
MIPS Pseudo Instructions and Functions

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pseudo instruction

Assembler



2

- Assembler convert readable instructions into machine code
 - assembly language: `add $t0, $s1, $s2`
 - machine code: `00000010 00110010 01000000 00100000`
- Make life easier with address labels

Address	Instruction
loop	...
	...
	j loop

Pseudo Instructions



- Some instructions would be nice to have
- For instance: load 32 bit value into register

```
li $s0, 32648278h
```

- Requires 2 instructions

```
lui $s0, 3264h  
ori $s0, $s0, 8278h
```

- Pseudo instruction
 - available in assembly
 - gets compiled into 2 machine code instructions

Syntactic Sugar



- Move

```
move $t0, $t1
```

- Compiled into add instruction

```
add $t0, $zero, $t
```

Reserved Register



- Example: load word from arbitrary memory address

```
lw $s0, 32648278h
```

- Memory address 32648278h has to be stored in register

- Solution: use reserved register \$at

```
lui $at, 3264h  
ori $at, $s0, 8278h  
lw $s0, 0($at)
```

Another Example



- Branch if less than

```
blt $t0, $t1, address
```

- Compiled into add instruction

```
slt $at, $t0, $t1  
bne $at, $zero, address
```

(slt = set if less than)

code example

Factorial



- Compute $n! = n \times n - 1 \times n - 2 \times \dots \times 2 \times 1$
- Iterative loop
 - initialize sum with n
 - loop through n-1, n-2, ..., 1
 - multiple sum with loop variable

Implementation



- Registers

- \$a0: n (loop variable)
- \$v0: sum

- Initialize

```
move $v0, $a0           # initialize sum with n
```

- Loop setup

```
loop:
    addi $a0, $a0, -1     # decrement n
    beq $a0, $zero, exit  # = 0? then done
    ...
    j loop
```

- Multiplication

```
mul $v0, $v0, $a0        # sum = sum * n
```



```
.text
main:
    li $a0, 5           # compute 5!
    move $v0, $a0       # initialize sum with n

loop:
    addi $a0, $a0, -1   # decrement n
    beq $a0, $zero, exit # = 0? then done
    mul $v0, $v0, $a0   # sum = sum * n
    j loop

exit:
    jr $ra              # done
```



jumps and subroutines

Jump

- MIPS instruction

j address

- Only 26 bits available for address (6 bits of op-code)■

⇒ 32 bit address constructed by concatenating

- upper 4 bits from current program counter
- 26 bits as specified
- 2 bits with value "0"■

- Proper 32 bit addressing available with

jr \$register

Jump and Link: Subroutines

- MIPS instructions

```
jal address  
jalr $register
```

- Address handling as before
- Stores return address in register \$ra (31st register)
- Return from subroutine

```
jr $ra
```

Register Conventions

- Arguments to subroutine: registers \$a0, \$a1, \$a2, \$a3
- Return values from subroutine: registers \$v0, \$v1, \$v2, \$v3
- Conceptually

$$(\$v0, \$v1, \$v2, \$v3) = f(\$a0, \$a1, \$a2, \$a3)$$

Example

15



- Subroutine to add three numbers

main:

```
li $a0, 10
li $a1, 21
li $a2, 33
jal add3
```

add3:

```
add $v0, $a0, $a1
add $v0, $v0, $a2
jr $ra
```


Another Example

- Subroutine for $a + b - c$

main:

```
li $a0, 10
li $a1, 21
li $a2, 33
jal add-and-sub
```

add-and-sub:

```
add $a0, $a0, $a1
move $a1, $a2
jal my-sub
jr $ra
```

my-sub:

```
sub $v0, $a0, $a1
jr $ra
```

- What could go wrong?

Safekeeping

- Recursive calls: must keep return address `$ra` in safe place
- May also want to preserve other registers
- Temporary registers `$t0-$t9` may be overwritten by subroutine
- Saved registers `$s0-$s7` must be preserved by subroutine
- Note
 - all this is by convention
 - you have to do this yourself



stack

- Recall: 6502
 - JSR stored return address on stack
 - RTS retrieved return address from stack
 - special instructions to store accumulator, status register
- MIPS: software stack
- By convention: stack pointer register \$sp (29th register)
- Why not always use the stack? It's slow

Alternate Idea

- Store return address in saved register \$s0
- But: now have to preserve \$s0 on stack

Store Return Address on Stack

- Decrease stack pointer

```
addi $sp, $sp, -4
```

32-bit address has 4 bytes■

- Store return address

```
sw $ra 0($sp)
```

sw = store word

- Stack pointer points to last used address

Retrieve Return Address from Stack

22



- Load return address

```
lw $ra 0($sp)
```

lw = store word

- Increase stack pointer

```
addi $sp, $sp, 4
```

Multiple Registers

- Store multiple registers

```
addi $sp, $sp, -12
sw $ra 0($sp)
sw $s0 4($sp)
sw $s1 8($sp)
```

- Load

```
lw $ra 0($sp)
lw $s0 4($sp)
lw $s1 8($sp)
addi $sp, $sp, 12
```


Frame Pointer

- What if we want to consult values stored on the stack?
- Example
 - subroutine stores return address and some save registers on stack
 - some code does something
(maybe even store more things on stack)
 - subroutine wants to consult stored return address■
- Stack pointer has changed
 - may be difficult to track down■
- Solution
 - store entry stack pointer in frame pointer \$fp (30th register)
 `move $fp, $sp`
 - retrieve return address using frame pointer
 `lw $s0, 0($fp)`



example

Recall: Factorial

.text

```
li $a0, 5           # compute 5!
move $v0, $a0        # initialize sum with n
```

loop:

```
addi $a0, $a0, -1    # decrement n
beq $a0, $zero, exit  # = 0? then done
mul $v0, $v0, $a0     # sum = sum * n
j loop
```

exit:

```
jr $ra              # done
```

Implemented as a Function

- Subroutine call (function argument in \$a0)

```
main:
    li $a0, 5           # compute 5!
    jal fact            # call function
```

- Return from subroutine (return value is in \$v0)

```
exit:
    jr $ra              # done
```

Scaffolding

```
.text
main:
    li $a0, 5           # compute 5!
    jal fact            # call function
    jr $ra              # done

fact:
    (old code)

exit:
    jr $ra              # done
```

Complete Code

```
.text
main:
    li $a0, 5           # compute 5!
    jal fact            # call function
    jr $ra              # done

fact:
    move $v0, $a0       # initialize sum with n

loop:
    addi $a0, $a0, -1    # decrement n
    beq $a0, $zero, exit # = 0? then done
    mul $v0, $v0, $a0    # sum = sum * n
    j loop

exit:
    jr $ra              # done
```

Recursive Implementation

- Idea: $f(n) = f(n-1) * n$
- Recursive call needs to preserve
 - return address
 - argument (n)

Termination Condition

- Check if argument is 0

fact:

```
    beq $a0, $zero, final    # = 0? then done
```

(common case)

final:

```
    li $v0, 1
    jr $ra                # done
```

- Note: no need to preserve registers

- Recursive call $f(n-1)$

```
addi $a0, $a0, -1      # decrement n
jal fact               # recursive call -> $v0 is f(n-1)
```

- Multiply with argument

```
mul $v0, $v0, $a0      # f(n-1) * n
```

Save and Restore Registers

- Save registers

```
addi $sp, $sp, -8
sw $ra 0($sp)      # return address on stack
sw $a0 4($sp)      # argument on stack
```

- Restore registers

```
lw $ra 0($sp)      # return address from stack
lw $a0 4($sp)      # argument from stack
addi $sp, $sp, 8
```

Complete Code

fact:

```
    beq $a0, $zero, final    # = 0? then done

    addi $sp, $sp, -8
    sw $ra 0($sp)            # return address on stack
    sw $a0 4($sp)            # argument on stack

    addi $a0, $a0, -1        # decrement n
    jal fact                  # recursive call -> $v0 is f(n-1)

    lw $ra 0($sp)            # return address from stack
    lw $a0 4($sp)            # argument from stack
    addi $sp, $sp, 8

    mul $v0, $v0, $a0        # f(n-1) * n
    jr $ra                   # done
```

final:

```
    li $v0, 1
    jr $ra                   # done
```