Final Exam

600.233 Computer Systems Fundamentals

Fall 2016 Johns Hopkins University

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Complete all questions.

Use additional paper if needed.

Time: 60 minutes.

Name of student:

Answer the following questions either with YES or NO.

- 1. Does loading a value from a hard-coded memory address takes two instructions in MIPS?
- 2. Does MIPS use a stack for the return address of a jump to subroutine?
- 3. Can a MIPS instruction carry out both a memory lookup and a subtraction?
- 4. To reduce cache misses with twice the cache memory is it better to double block size than double the number of blocks?
- 5. Is a joint component for instruction fetch and data fetch an example for a structural hazard?
- 6. Do x86 processors follow the Complex Instruction Set Computer (CISC) design?
- 7. Does the linker change the addresses of machine code instructions?
- 8. Do dynamically linked shared libraries avoid adjustment of addresses because they are loaded always into the same memory location?
- 9. Can you always avoid using the stack and use registers instead for return addresses in function calls?
- 10. Is it possible to run a process as fast in a virtual machine as in real machine?
- 11. Do you want 5 points?

Q2. Caching	25 points
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Consider use of a 2-way associative cache that addresses blocks of 4 bytes, with 4 sets in a 8-bit address space.

(a) How are the 8 bits of the address used as tag, index, and offset for the cache?

(b) Consider a following sequence of requests to the cache.

Enter the **tag** for each cache slot after each request in the table below. Assume FIFO as caching strategy (do not worry about internal bookkeeping of timestamps). Note: use " to indicate that the value in the slot is identical to the previous value.

Request	Set 0		Set 1		Set 2		Set 3	
	Slot 0	Slot 1						
	empty							
00110101								
01101000								
01101001								
10010111								
10010110								
10110001								
10110101								

Q3. Pipelining 25 points

Consider the following sequence of MIPS commands.

```
0000 load $s0, 20($t0)
0004 load $s1, 24($t0)
0008 sub $s2, $s0, $s1
000c bne $s2, $zero, 0018 ; assume that the test fails
0010 add $s4, $s2, $s1
0014 add $s2, $s2, $s1
0018 jr $ra ; return from subroutine
```

(a) Identify a data hazard in the code.

(b) Identify a control hazard in the code.

(c) Complete the table on the next page that shows which command is currently in which stage of the CPU.

Identify each command by its address and operation name. End with instruction fetch of the jr command. Add more lines / use less lines as needed.

State your assumptions about handling data hazards and control hazard.

Time	Instruction Fetch	Instr. Decode	Execute / ALU	Memory Access	Write-Back
		/ Register Read			
1	0000 load	-	-	-	-
2					
3					
4					
5					
6					
7					
8					
9					
10					
11					
12					
13					
14					
15					

Q4. Virtual Memory

25 points

(a) Process memory.

Draw a diagram that shows:

- virtual memory layout of a process,
- its page table,
- the system RAM,
- and the system disk

The diagram should include the following memory areas of a process:

- code
- shared libraries
- stack, and
- allocated memory for data structures.

Note: assume that the page table is just a flat lookup table and the RAM does not use any caches.

(b) Change in memory use.

Indicate how entries in page table and memory use in RAM change ...

- 1. ... when process is launched
- 2. ... when process executed some instructions
- 3. ... when process allocates memory for a data structure

(c) Multiple processes.

Now a second process using the same code is launched. Draw its page table and memory use and how it relates to the first process.