



# **FFTs in Graphics and Vision**

The Fast Fourier Transform

# Outline

The FFT Algorithm

Multi-Dimensional FFTs

More Applications

Real FFTs





# Computational Complexity

To compute the correlation of two periodic,  $n$ -dimensional arrays  $\mathbf{f}, \mathbf{g} \in \mathbb{C}^n$ :

1. Express  $\mathbf{f}$  and  $\mathbf{g}$  in the basis  $\{\mathbf{z}^0, \dots, \mathbf{z}^{n-1}\} \subset \mathbb{C}^n$ :

$$\mathbf{f} = \sum_{k=0}^{n-1} \hat{f}_k \cdot \mathbf{z}^k \quad \text{and} \quad \mathbf{g} = \sum_{k=0}^{n-1} \hat{g}_k \cdot \mathbf{z}^k$$

2. Multiply (and scale) the coefficients:

$$(\mathbf{f} \star \mathbf{g}) = \sqrt{2\pi} \cdot \sum_{k=0}^{n-1} \hat{f}_k \cdot \bar{\hat{g}}_k \cdot \mathbf{z}^{-k}$$

3. Evaluate at every index  $j$ :

$$(\mathbf{f} \star \mathbf{g})_j = \sqrt{2\pi} \cdot \sum_{k=0}^{n-1} \hat{f}_k \cdot \bar{\hat{g}}_k \cdot z_j^{-k}$$



# Goal

Given  $\mathbf{g} \in \mathbb{C}^n$ , we would like to compute the Fourier coefficients  $\hat{\mathbf{g}}$ :

$$\mathbf{g} = \sum_{k=0}^{n-1} \hat{g}_k \cdot \mathbf{z}^k$$

with  $\mathbf{z}^k$  the normalized discrete samples of the complex exponentials at  $n$  regular positions:

$$\mathbf{z}^k = \left( \frac{e^{ik\theta_0}}{\sqrt{2\pi}}, \dots, \frac{e^{ik\theta_{n-1}}}{\sqrt{2\pi}} \right), \quad \theta_j = \frac{2j\pi}{n}$$

# Challenge



How can we compute all  $n$  Fourier coefficients efficiently?

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Since the vectors  $\{\mathbf{z}^0, \dots, \mathbf{z}^{n-1}\}$  are orthonormal, we can compute the  $k$ -th Fourier coefficient of  $\mathbf{g}$  by computing the dot-product:

$$\begin{aligned}\hat{g}_k &= \langle \mathbf{g}, \mathbf{z}^k \rangle_{[0, 2\pi)} \\ &= \frac{2\pi}{n} \cdot \sum_{j=0}^{n-1} g_j \cdot \bar{z}_j^k \\ &= \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{n-1} g_j \cdot e^{-\frac{ik2\pi j}{n}}\end{aligned}$$



# Challenge

How can we compute all  $n$  Fourier coefficients?

Since the vectors  $\{\mathbf{z}^0, \dots, \mathbf{z}^{n-1}\}$  are orthonormal, we can compute the  $k$ -th Fourier coefficient of  $\mathbf{g}$  by computing the dot-product:

$$\hat{g}_k = \langle \mathbf{g}, \mathbf{z}^k \rangle_{[0, 2\pi)}$$

Computing one coefficient has complexity  $O(n)$

Computing all of them has complexity  $O(n^2)$

$$\sum_{j=0}^{n-1} g_j e^{-\pi j}$$



# Challenge



How can we compute all  $n$  Fourier coefficients efficiently?

# Approach (Divide and Conquer)



## Key Idea:

If we decompose the array into the even and odd halves, we can solve smaller problems and then combine.



# Approach (Divide and Conquer)

Key Idea:

Take the 8-dimensional array:

$$\mathbf{g} = (g_0, g_1, g_2, g_3, g_4, g_5, g_6, g_7)$$

And consider its even/odd decomposition:

$$\mathbf{g}^e = (g_0, g_2, g_4, g_6)$$

$$\mathbf{g}^o = (g_1, g_3, g_5, g_7)$$

How are the coefficients of  $\mathbf{g}$  are weighted when computing the  $k$ -th Fourier coefficient, relative to how the coefficients of  $\mathbf{g}^e$  and  $\mathbf{g}^o$  are weighted?



# Approach (Divide and Conquer)

Key Idea:

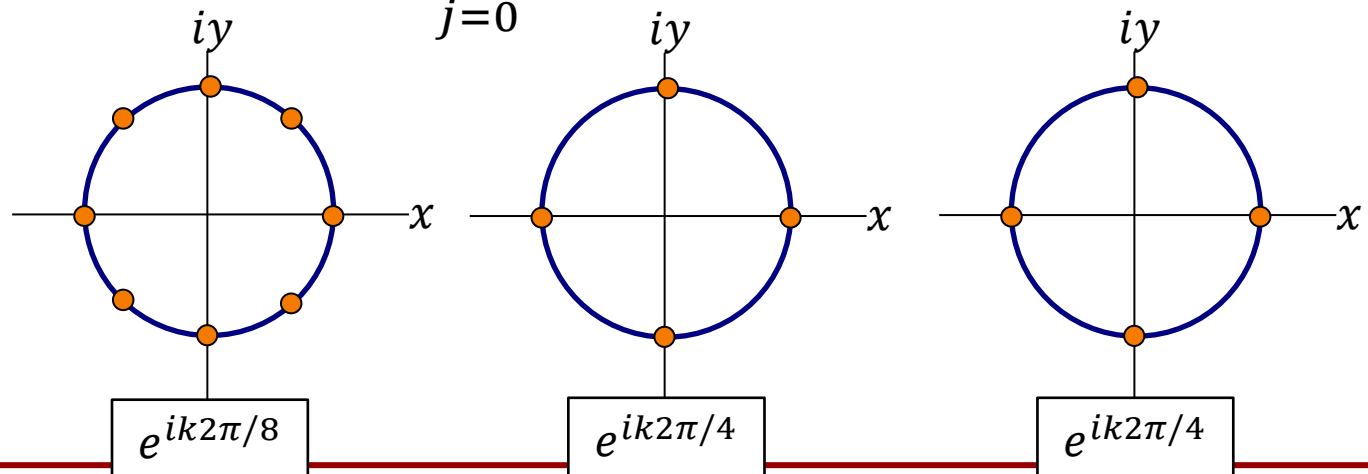
$$\mathbf{g} = (g_0, g_1, g_2, g_3, g_4, g_5, g_6, g_7)$$

$$\mathbf{g}^e = (g_0, g_2, g_4, g_6)$$

$$\mathbf{g}^o = (g_1, g_3, g_5, g_7)$$

Consider the Fourier coefficients:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{n-1} g_j \cdot e^{-\frac{ik2\pi j}{n}}$$





# Approach (Divide and Conquer)

Key Idea:

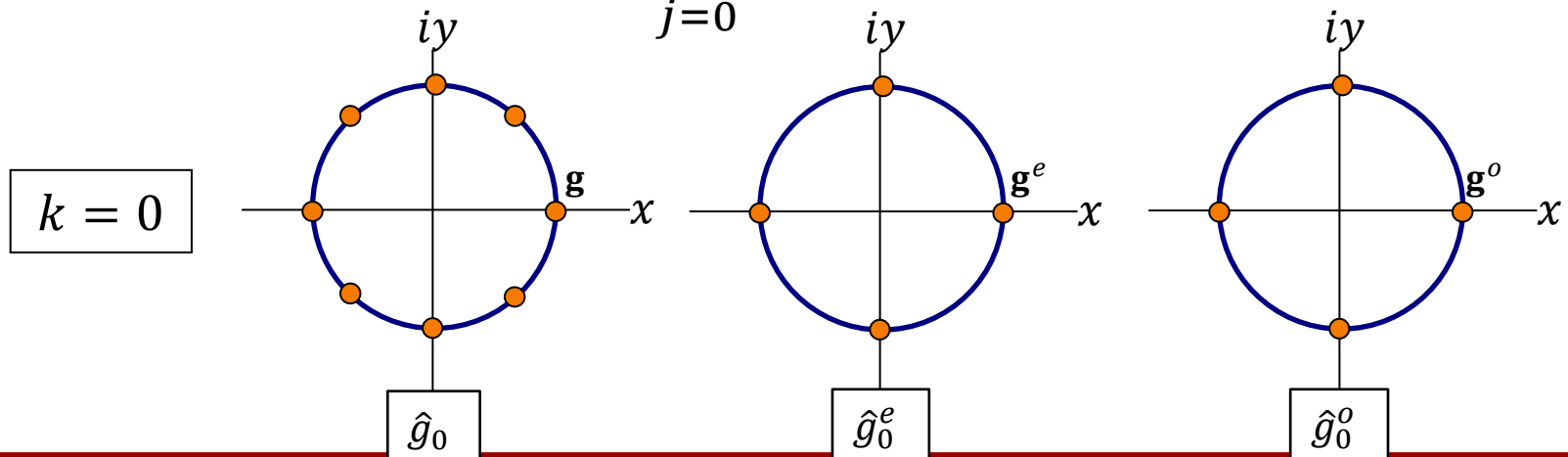
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Consider the Fourier coefficients:

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# Approach (Divide and Conquer)

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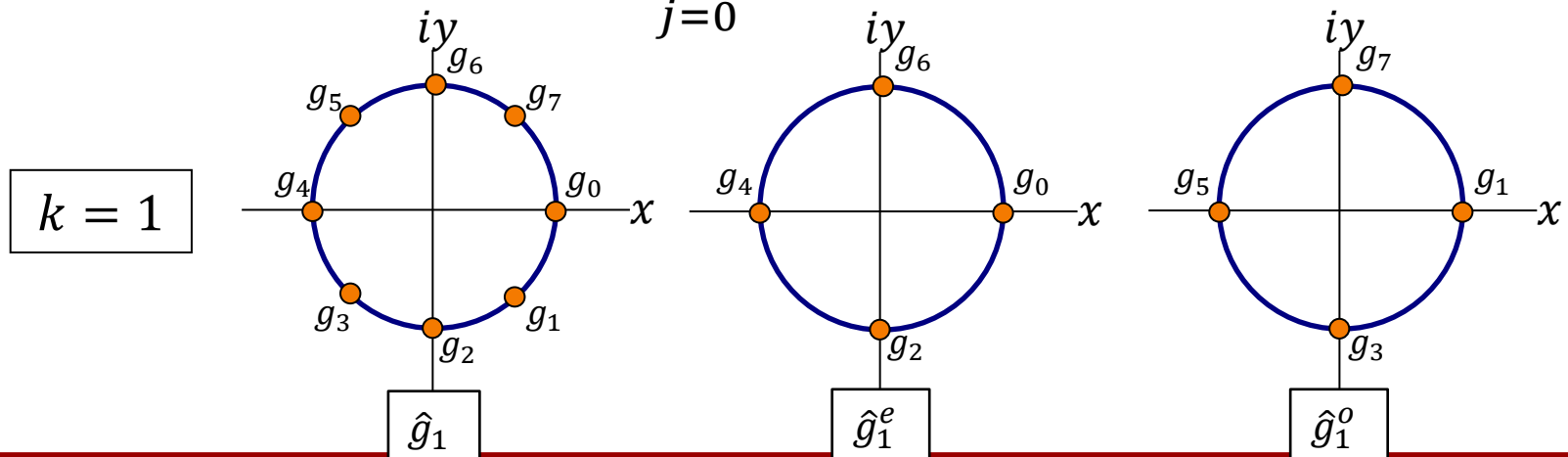
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Consider the Fourier coefficients:

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# Approach (Divide and Conquer)

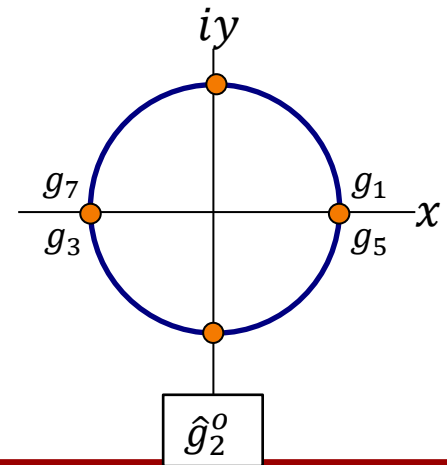
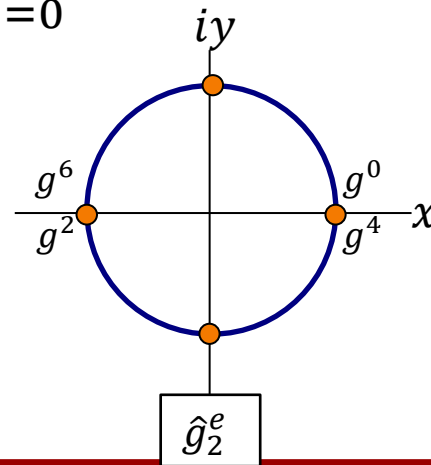
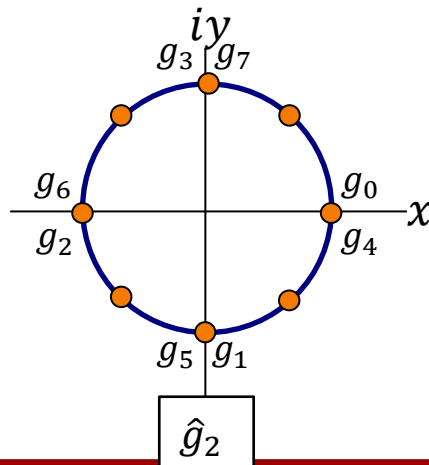
Key Idea:

$$\mathbf{g} = (g_0, g_1, g_2, g_3, g_4, g_5, g_6, g_7)$$
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Consider the Fourier coefficients:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{n-1} g_j \cdot e^{-\frac{ik2\pi j}{n}}$$

$$k = 2$$





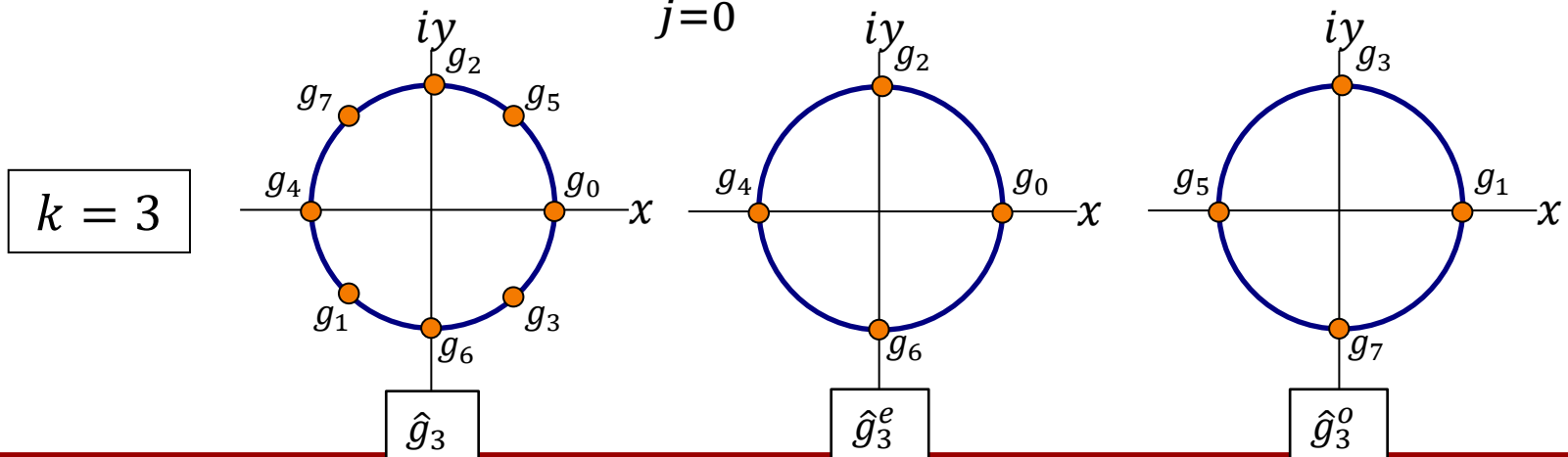
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Key Idea:

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Consider the Fourier coefficients:

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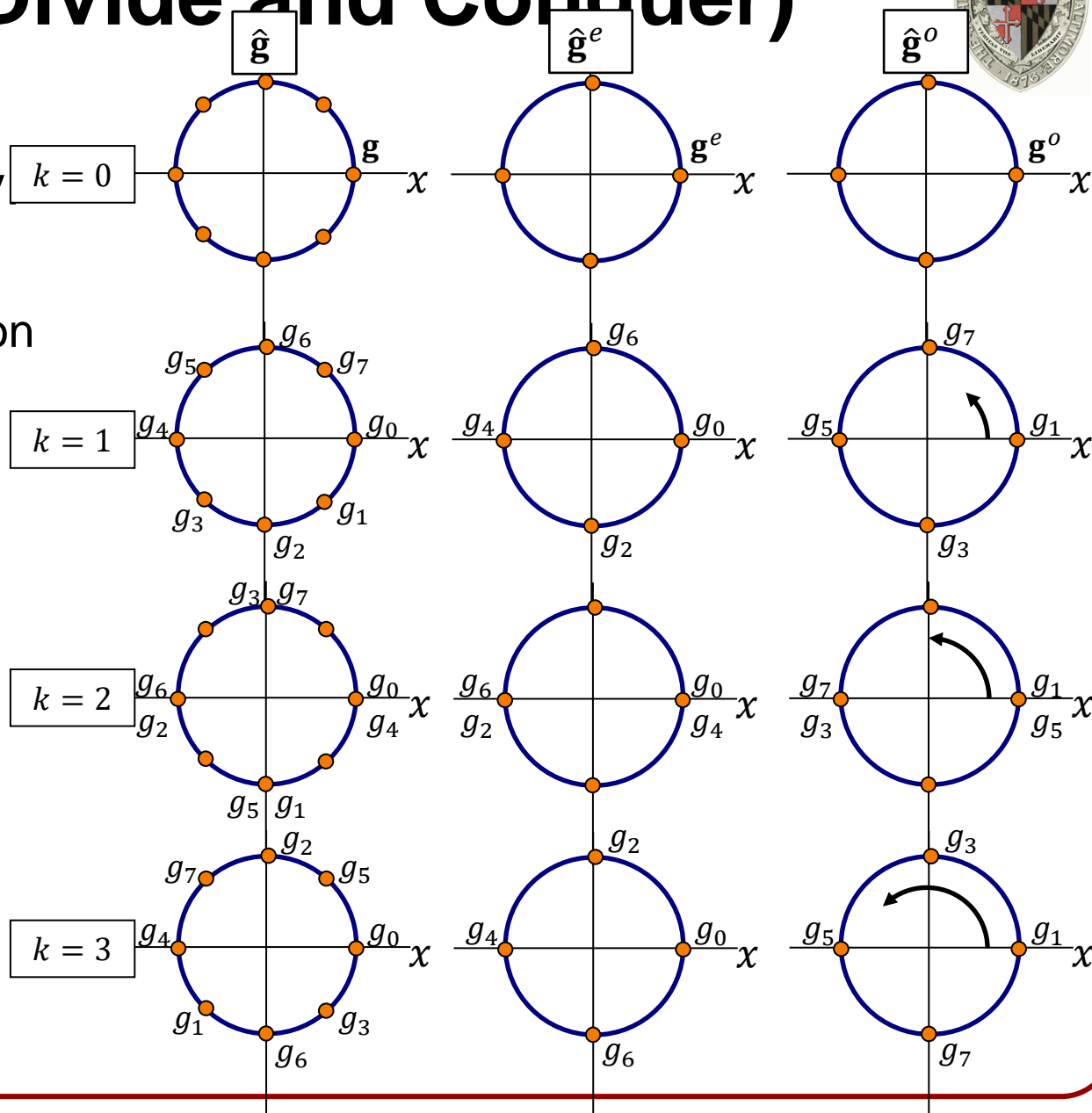


# Approach (Divide and Conquer)



For every frequency

- The even components are in the same position as the original
- The odd ones are off by a fixed angle.





# Approach (Divide and Conquer)

Assume that  $n$  is even ( $n = 2m$ ) and consider the even and odd entries separately.

That is, let  $\mathbf{g}^e, \mathbf{g}^o \in \mathbb{C}^m$  be the (periodic) arrays:

$$g_k^e = g_{2k} \quad g_k^o = g_{2k+1}$$



# Approach (Divide and Conquer)

$$g_k^e = g_{2k} \quad g_k^o = g_{2k+1}$$

The  $k$ -th Fourier coefficient of  $\mathbf{g}$  is:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{n-1} g_j \cdot e^{-\frac{ik2\pi j}{n}}$$

Splitting the summation into even/odd terms:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{m-1} \left( g_{2j} \cdot e^{-\frac{ik2\pi(2j)}{2m}} + g_{2j+1} \cdot e^{-\frac{ik2\pi(2j+1)}{2m}} \right)$$



# Approach (Divide and Conquer)

$$g_k^e = g_{2k} \quad g_k^o = g_{2k+1}$$

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Re-writing the exponent, we get:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{m-1} \left( g_{2j} \cdot e^{-\frac{ik2\pi j}{m}} + g_{2j+1} \cdot e^{-\frac{ik2\pi j}{m}} \cdot e^{-\frac{ik2\pi}{n}} \right)$$



# Approach (Divide and Conquer)

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Plugging in our even/odd expression gives:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{m-1} \left( g_j^e \cdot e^{-\frac{ik2\pi j}{m}} + g_j^o \cdot e^{-\frac{ik2\pi j}{m}} \cdot e^{-\frac{ik2\pi}{n}} \right)$$



# Approach (Divide and Conquer)

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Plugging in our even/odd expression gives:

$$\hat{g}_k = \frac{\sqrt{2\pi}}{n} \cdot \sum_{j=0}^{m-1} \left( g_j^e \cdot e^{-\frac{ik2\pi j}{m}} + g_j^o \cdot e^{-\frac{ik2\pi j}{m}} \cdot e^{-\frac{ik2\pi}{n}} \right)$$

Re-writing this equation gives:

$$\hat{g}_k = \frac{1}{2} \left( \frac{\sqrt{2\pi}}{m} \cdot \sum_{j=0}^{m-1} g_j^e \cdot e^{-\frac{ik2\pi j}{m}} \right) + \frac{1}{2} \left( e^{-\frac{ik2\pi}{n}} \cdot \frac{\sqrt{2\pi}}{m} \cdot \sum_{j=0}^{m-1} g_j^o \cdot e^{-\frac{ik2\pi j}{m}} \right)$$



# Approach (Divide and Conquer)

$$g_k^e = g_{2k} \quad g_k^o = g_{2k+1}$$

Re-writing this equation gives:

$$\hat{g}_k = \frac{1}{2} \left( \frac{\sqrt{2\pi}}{m} \cdot \sum_{j=0}^{m-1} g_j^e \cdot e^{-\frac{ik2\pi j}{m}} \right) + \frac{1}{2} \left( e^{-\frac{ik2\pi}{n}} \cdot \frac{\sqrt{2\pi}}{m} \cdot \sum_{j=0}^{m-1} g_j^o \cdot e^{-\frac{ik2\pi j}{m}} \right)$$

This is a linear combination of Fourier coefficients:

$$\hat{g}_k = \frac{1}{2} \cdot \left( \hat{g}_k^e + e^{-\frac{ik2\pi}{n}} \cdot \hat{g}_k^o \right)$$



# Approach (Divide and Conquer)

$$g_k^e = g_{2k} \quad g_k^o = g_{2k+1}$$
$$\hat{g}_k = \frac{1}{2} \cdot \left( \hat{g}_k^e + e^{-\frac{ik2\pi}{n}} \cdot \hat{g}_k^o \right)$$

If we get the Fourier coefficients of  $\mathbf{g}^e$  and  $\mathbf{g}^o$ , we can combine them to get the Fourier coefficients of  $\mathbf{g}$ .

Assuming that  $m$  is also even, we can repeat for  $\mathbf{g}^e$  and  $\mathbf{g}^o$  to get their Fourier coefficients.

If  $n$  is a power of 2, we can keep splitting, to get an  $O(n \log n)$  algorithm.



# The Inverse Fourier Transform



It turns out that computing the inverse Fourier transform is (almost) equivalent to computing the forward transform...



# The Inverse Fourier Transform

The Fourier Transform is a change of basis:

Complex Exponential Basis

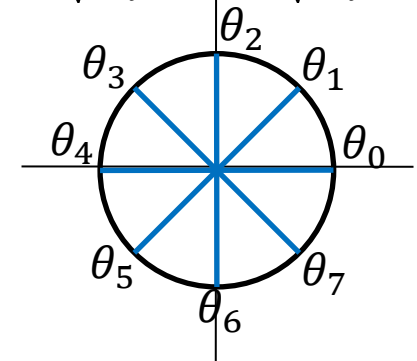
Evaluation Basis

$(1, 0, \dots, 0, 0)$   
 $(0, 1, \dots, 0, 0)$   
 $\vdots$   
 $(0, 0, \dots, 1, 0)$   
 $(0, 0, \dots, 0, 1)$

Fourier  
 Transform  $\rightarrow$

$$\begin{pmatrix} \frac{e^{i0\theta_0}}{\sqrt{2\pi}}, \frac{e^{i0\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i0\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i0\theta_{n-1}}}{\sqrt{2\pi}} \\
 \frac{e^{i1\theta_0}}{\sqrt{2\pi}}, \frac{e^{i1\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i1\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i1\theta_{n-1}}}{\sqrt{2\pi}} \\
 \vdots
 \end{pmatrix}$$

$$\begin{pmatrix} \frac{e^{i(n-2)\theta_0}}{\sqrt{2\pi}}, \frac{e^{i(n-2)\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i(n-2)\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i(n-2)\theta_{n-1}}}{\sqrt{2\pi}} \\
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 \end{pmatrix}$$





# The Inverse Fourier Transform

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Complex Exponential Basis

Evaluation Basis

$(1, 0, \dots, 0, 0)$   
 $(0, 1, \dots, 0, 0)$   
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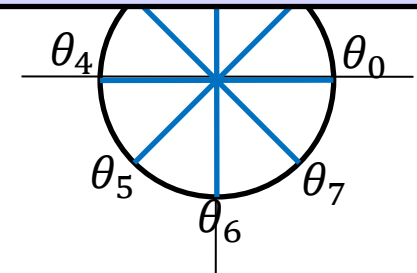
Fourier  
 Transform  $\rightarrow$

$$\begin{pmatrix} \frac{e^{i0\theta_0}}{\sqrt{2\pi}}, \frac{e^{i0\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i0\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i0\theta_{n-1}}}{\sqrt{2\pi}} \\ \frac{e^{i1\theta_0}}{\sqrt{2\pi}}, \frac{e^{i1\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i1\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i1\theta_{n-1}}}{\sqrt{2\pi}} \\ \vdots \\ \frac{e^{i(n-2)\theta_0}}{\sqrt{2\pi}}, \frac{e^{i(n-2)\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i(n-2)\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i(n-2)\theta_{n-1}}}{\sqrt{2\pi}} \\ \frac{e^{i(n-1)\theta_0}}{\sqrt{2\pi}}, \frac{e^{i(n-1)\theta_1}}{\sqrt{2\pi}}, \dots, \frac{e^{i(n-1)\theta_{n-2}}}{\sqrt{2\pi}}, \frac{e^{i(n-1)\theta_{n-1}}}{\sqrt{2\pi}} \end{pmatrix}$$

Note:

The evaluation basis is not orthonormal.

It can be made orthonormal by scaling all the basis vectors by the same scale factor,  $\sqrt{n/2\pi}$ .





# The Inverse Fourier Transform

Since the Fourier basis  $\{\mathbf{z}^0, \dots, \mathbf{z}^{n-1}\}$  is orthonormal, the  $k$ -th Fourier coefficient of a vector  $\mathbf{a} \in \mathbb{C}^n$  is the (Hermitian) dot-product:

$$\begin{aligned}\hat{a}_k &= \langle \mathbf{a}, \mathbf{z}^k \rangle_{[0, 2\pi)} \\ &= \frac{2\pi}{n} (\bar{z}_0^k, \dots, \bar{z}_{n-1}^k) \cdot \begin{pmatrix} a_0 \\ \vdots \\ a_{n-1} \end{pmatrix} \\ &= \frac{2\pi}{n} \left( \frac{e^{-ik \cdot 0\theta}}{\sqrt{2\pi}}, \dots, \frac{e^{-ik \cdot (n-1)\theta}}{\sqrt{2\pi}} \right) \cdot \begin{pmatrix} a_0 \\ \vdots \\ a_{n-1} \end{pmatrix}\end{aligned}$$

where  $\theta$  is the angle  $\theta = \frac{2\pi}{n}$ .



# The Inverse Fourier Transform

$$\hat{a}_k = \frac{\sqrt{2\pi}}{n} (e^{-ik \cdot 0\theta}, \dots, e^{-ik \cdot (n-1)\theta}) \cdot \begin{pmatrix} a_0 \\ \vdots \\ a_{n-1} \end{pmatrix}, \quad \omega / \theta = \frac{2\pi}{n}$$

Stacking the rows we can represent the Fourier transform by the matrix:

$$\mathbf{F} = \frac{\sqrt{2\pi}}{n} \begin{pmatrix} 1 & 1 & \dots & 1 & 1 \\ 1 & e^{-i\theta} & \dots & e^{-i(n-2)\theta} & e^{-i(n-1)\theta} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & e^{-i(n-2)\theta} & \dots & e^{-i(n-2)(n-2)\theta} & e^{-i(n-2)(n-1)\theta} \\ 1 & e^{-i(n-1)\theta} & \dots & e^{-i(n-1)(n-2)\theta} & e^{-i(n-1)(n-1)\theta} \end{pmatrix}.$$



# The Inverse Fourier Transform

$$\mathbf{F} = \frac{\sqrt{2\pi}}{n} \begin{pmatrix} 1 & 1 & \dots & 1 & 1 \\ 1 & e^{-i\theta} & \dots & e^{-i(n-2)\theta} & e^{-i(n-1)\theta} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & e^{-i(n-2)\theta} & \dots & e^{-i(n-2)(n-2)\theta} & e^{-i(n-2)(n-1)\theta} \\ 1 & e^{-i(n-1)\theta} & \dots & e^{-i(n-1)(n-2)\theta} & e^{-i(n-1)(n-1)\theta} \end{pmatrix}$$

Since the evaluation and Fourier bases are (essentially) orthonormal\*, the matrix is unitary, and the inverse Fourier transform is (up to a constant) the transpose conjugate of the forward transform.

\*Up to a scale factor of  $\sqrt{n/2\pi}$ , the evaluation basis is orthonormal.



# The Inverse Fourier Transform

$$\mathbf{F} = \frac{\sqrt{2\pi}}{n} \begin{pmatrix} 1 & 1 & \dots & 1 & 1 \\ 1 & e^{-i\theta} & \dots & e^{-i(n-2)\theta} & e^{-i(n-1)\theta} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & e^{-i(n-2)\theta} & \dots & e^{-i(n-2)(n-2)\theta} & e^{-i(n-2)(n-1)\theta} \\ 1 & e^{-i(n-1)\theta} & \dots & e^{-i(n-1)(n-2)\theta} & e^{-i(n-1)(n-1)\theta} \end{pmatrix}$$

In particular, given the Fourier coefficients:

$$\hat{\mathbf{a}} = (\hat{a}_0, \dots, \hat{a}_{n-1})$$

the inverse Fourier transform is:

$$\mathbf{F}^{-1} \hat{\mathbf{a}} = c \cdot \bar{\mathbf{F}}^t \hat{\mathbf{a}}$$

(for some real constant  $c$  accounting for the non-orthonormality of the evaluation basis).



# The Inverse Fourier Transform

$$\mathbf{F} = \frac{\sqrt{2\pi}}{n} \begin{pmatrix} 1 & 1 & \dots & 1 & 1 \\ 1 & e^{-i\theta} & \dots & e^{-i(n-2)\theta} & e^{-i(n-1)\theta} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & e^{-i(n-2)\theta} & \dots & e^{-i(n-2)(n-2)\theta} & e^{-i(n-2)(n-1)\theta} \\ 1 & e^{-i(n-1)\theta} & \dots & e^{-i(n-1)(n-2)\theta} & e^{-i(n-1)(n-1)\theta} \end{pmatrix}$$

Taking the double conjugate, we get:

$$\begin{aligned} \mathbf{F}^{-1} \hat{\mathbf{a}} &= c \cdot \bar{\mathbf{F}}^t \hat{\mathbf{a}} \\ &= c \cdot \overline{\overline{\mathbf{F}}^t \hat{\mathbf{a}}} \\ &= c \cdot \mathbf{F}^t \bar{\hat{\mathbf{a}}} \end{aligned}$$





# The Inverse Fourier Transform

$$\mathbf{F} = \frac{\sqrt{2\pi}}{n} \begin{pmatrix} 1 & 1 & \dots & 1 & 1 \\ 1 & e^{-i\theta} & \dots & e^{-i(n-2)\theta} & e^{-i(n-1)\theta} \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 1 & e^{-i(n-2)\theta} & \dots & e^{-i(n-2)(n-2)\theta} & e^{-i(n-2)(n-1)\theta} \\ 1 & e^{-i(n-1)\theta} & \dots & e^{-i(n-1)(n-2)\theta} & e^{-i(n-1)(n-1)\theta} \end{pmatrix}$$

Since  $\mathbf{F} = \mathbf{F}^t$ :

$$\begin{aligned} \mathbf{F}^{-1} \hat{\mathbf{a}} &= c \cdot \overline{\mathbf{F}^t \hat{\mathbf{a}}} \\ &= c \cdot \overline{\mathbf{F} \hat{\mathbf{a}}} \end{aligned}$$

⇒ Up to a constant scale factor, the inverse Fourier transform can be obtained by taking the conjugate of the coefficients, computing the forward transform, and conjugating the result.



# The Inverse Fourier Transform

$$\mathbf{F}^{-1}\hat{\mathbf{a}} = c \cdot \overline{\mathbf{F}\hat{\mathbf{a}}}$$

Q: What is the constant  $c$ ?

A: The matrix  $\mathbf{F}$ , gives the Fourier coefficients of a vector represented in the evaluation basis.

⇒ The evaluation basis becomes orthonormal if we scale by  $\sqrt{n/2\pi}$ .

⇒ The matrix  $\sqrt{n/2\pi} \cdot \mathbf{F}$  is a change of basis between two orthonormal bases.

⇒ The matrix  $\sqrt{n/2\pi} \cdot \mathbf{F}$  is unitary.



# The Inverse Fourier Transform

$$\mathbf{F}^{-1}\hat{\mathbf{a}} = c \cdot \overline{\mathbf{F}\hat{\mathbf{a}}}$$

⇒ The matrix  $\sqrt{n/2\pi} \cdot \mathbf{F}$  is unitary.

⇒ Its inverse is its conjugate transpose:

$$\left(\sqrt{n/2\pi} \cdot \mathbf{F}\right)^{-1} = \sqrt{n/2\pi} \cdot \overline{\mathbf{F}}^t$$



$$\mathbf{F}^{-1} = n/2\pi \cdot \overline{\mathbf{F}}^t$$

⇒ The inverse Fourier transform is defined by:

$$\mathbf{F}^{-1}\hat{\mathbf{a}} = n/2\pi \cdot \overline{\mathbf{F}\hat{\mathbf{a}}}$$



# The Inverse Fourier Transform

$$\mathbf{F}^{-1}\hat{\mathbf{a}} = n/2\pi \cdot \overline{\mathbf{F}\hat{\mathbf{a}}}$$

Normalized Inverse Fourier Transform:

1. Take the conjugate of the Fourier coefficients
2. Compute the forward Fourier transform
3. Take the conjugate of the resultant coefficients
4. Scale by  $n/2\pi$



# The Inverse Fourier Transform

$$\mathbf{F}^{-1} \hat{\mathbf{a}} = n/2\pi \cdot \overline{\mathbf{F} \hat{\mathbf{a}}}$$

Note:

To avoid an expensive square root, most FFT implementations compute the un-normalized Fourier coefficients:

$$\tilde{a}_k = \langle \tilde{\mathbf{z}}^k, \mathbf{a} \rangle = \frac{n}{2\pi} \sqrt{2\pi} \langle \mathbf{z}^k, \mathbf{a} \rangle_{[0, 2\pi)} = n/\sqrt{2\pi} \hat{a}_k$$

⇒ The two Fourier transforms are related by:

$$\tilde{\mathbf{F}} = n/\sqrt{2\pi} \mathbf{F}$$



# The Inverse Fourier Transform

$$\mathbf{F}^{-1} \hat{\mathbf{a}} = n/2\pi \cdot \overline{\mathbf{F} \hat{\tilde{\mathbf{a}}}}$$
$$\tilde{\mathbf{F}} = n/\sqrt{2\pi} \mathbf{F}$$

Note:

Taking the inverse, we get:

$$\begin{aligned} \tilde{\mathbf{F}}^{-1} \tilde{\hat{\mathbf{a}}} &= \sqrt{2\pi}/n \cdot \mathbf{F}^{-1} \tilde{\hat{\mathbf{a}}} \\ &= \sqrt{2\pi}/n \cdot n/2\pi \cdot \overline{\mathbf{F} \tilde{\hat{\mathbf{a}}}} \\ &= \sqrt{2\pi}/n \cdot n/2\pi \cdot \sqrt{2\pi}/n \cdot \overline{\tilde{\mathbf{F}} \tilde{\hat{\mathbf{a}}}} \\ &= 1/n \cdot \overline{\tilde{\mathbf{F}} \tilde{\hat{\mathbf{a}}}} \end{aligned}$$



# The Inverse Fourier Transform

$$\tilde{\mathbf{F}}^{-1} \tilde{\mathbf{a}} = 1/n \cdot \overline{\tilde{\mathbf{F}} \tilde{\mathbf{a}}}$$

Unnormalized Inverse Fourier Transform:

1. Take the conjugate of the Fourier coefficients
2. Compute the forward Fourier transform
3. Take the conjugate of the resultant coefficients
4. Scale by  $1/n$

Implementations like the FFTW do not  
adjust for the scaling term.



Running the forward Fourier Transform and then  
the inverse scales the input by a factor of  $1/n$ .



# Outline

The FFT Algorithm

**Multi-Dimensional FFTs**

More Applications

Real FFTs





# Multi-Dimensional FFTs

How do we compute the Fourier transform of a multi-dimensional signal?

## Examples:

- Images
- Voxel Grids
- Etc.



# 2D FFTs

For regularly sampled,  $n \times n$  grids, the irreducible representations are spanned by the orthonormal basis  $\{\mathbf{z}^{lm}\}$  where:

$$z_{jk}^{lm} = \frac{1}{2\pi} \cdot e^{\frac{il2\pi j}{n}} \cdot e^{\frac{im2\pi k}{n}}$$

To see this, consider the shift by  $(\alpha, \beta) \in \mathbb{Z}^2$  :

$$\begin{aligned} \left( \rho_{\alpha, \beta}(\mathbf{z}^{lm}) \right)_{jk} &= z_{j-\alpha, k-\beta}^{lm} \\ &= \frac{1}{2\pi} \cdot e^{\frac{il2\pi(j-\alpha)}{n}} \cdot e^{\frac{im2\pi(k-\beta)}{n}} \\ &= \left( e^{-\frac{il2\pi\alpha}{n}} \cdot e^{-\frac{im2\pi\beta}{n}} \right) \cdot z_{jk}^{lm} \end{aligned}$$



# 2D FFTs

How can we compute the 2D Fourier coefficients efficiently?

To do this, we will leverage the fact that the 2D basis vectors can be expressed as the product of 1D basis vectors.

Setting  $\mathbf{z}^l$  to be the  $n$ -dimensional array:

$$z_j^l = \frac{1}{\sqrt{2\pi}} \cdot e^{\frac{il2\pi j}{n}}$$

we get:

$$z_{jk}^{lm} = z_j^l \cdot z_k^m$$



# 2D FFTs

To compute the  $(l, m)$ -th Fourier coefficient of an  $(n \times n)$ -dimensional grid  $\mathbf{g}$ , we compute the dot-product of  $\mathbf{g}$  with  $\mathbf{z}^{lm}$ :

$$\begin{aligned}\hat{g}_{lm} &= \langle \mathbf{g}, \mathbf{z}^{lm} \rangle_{[0, 2\pi) \times [0, 2\pi)} \\ &= \frac{(2\pi)^2}{n^2} \cdot \sum_{j=0}^{n-1} \sum_{k=0}^{n-1} g_{jk} \cdot \bar{z}_{jk}^{lm} \\ &= \frac{(2\pi)^2}{n^2} \cdot \sum_{j=0}^{n-1} \sum_{k=0}^{n-1} g_{jk} \cdot \bar{z}_j^l \cdot \bar{z}_k^m \\ &= \frac{2\pi}{n} \cdot \sum_{j=0}^{n-1} \left( \frac{2\pi}{n} \cdot \sum_{k=0}^{n-1} g_{jk} \cdot \bar{z}_k^m \right) \cdot \bar{z}_j^l\end{aligned}$$



# 2D FFTs

$$\hat{\mathbf{g}}_{lm} = \frac{2\pi}{n} \cdot \sum_{j=0}^{n-1} \left( \frac{2\pi}{n} \cdot \sum_{k=0}^{n-1} g_{jk} \cdot \bar{z}_k^m \right) \cdot \bar{z}_j^l$$

The interior sum is a 1D Fourier transform!

In particular, setting  $\mathbf{g}^j$  to be the 1D array:

$$g_k^j = g_{jk}$$

we get:

$$\begin{aligned} \hat{g}_{lm} &= \frac{2\pi}{n} \cdot \sum_{j=0}^{n-1} \left( \frac{2\pi}{n} \cdot \sum_{k=0}^{n-1} g_k^j \cdot \bar{z}_k^m \right) \cdot \bar{z}_j^l \\ &= \frac{2\pi}{n} \cdot \sum_{j=0}^{n-1} \hat{g}_m^j \cdot \bar{z}_j^l \end{aligned}$$



# 2D FFTs

$$\hat{g}_{lm} = \frac{2\pi}{n} \cdot \sum_{j=0}^{n-1} \hat{g}_m^j \cdot \bar{z}_j^l$$

But this sum is also a 1D Fourier transform!

Thus, to compute the Fourier coefficients, we:

1. Compute the Fourier coefficients of each row
2. Compute the Fourier coefficients of each column

And the total complexity of this operation is:

$$O(n^2 \log n)$$



# Outline

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More Applications

- Gaussian Smoothing
- Up-Sampling
- Differentiation
- Boundary Detection
- Gaussian Sharpening

Real FFTs



# Gaussian Smoothing

Given an  $n \times n$  grid of values, we would like to smooth the grid.







# Gaussian Smoothing

Given an  $n \times n$  grid of values, we would like to smooth the grid.

To do this we need:

- **f**: The smoothing filter, usually a Gaussian:

$$\tilde{f}_{jk} = e^{-(j^2+k^2)/2\sigma^2} \quad \text{with} \quad -\frac{n}{2} < j, k \leq \frac{n}{2}$$

normalized to have unit mass:

$$f_{jk} = \frac{\tilde{f}_{jk}}{\langle \tilde{\mathbf{f}}, 1 \rangle_{[0,2\pi) \times [0,2\pi)}}$$

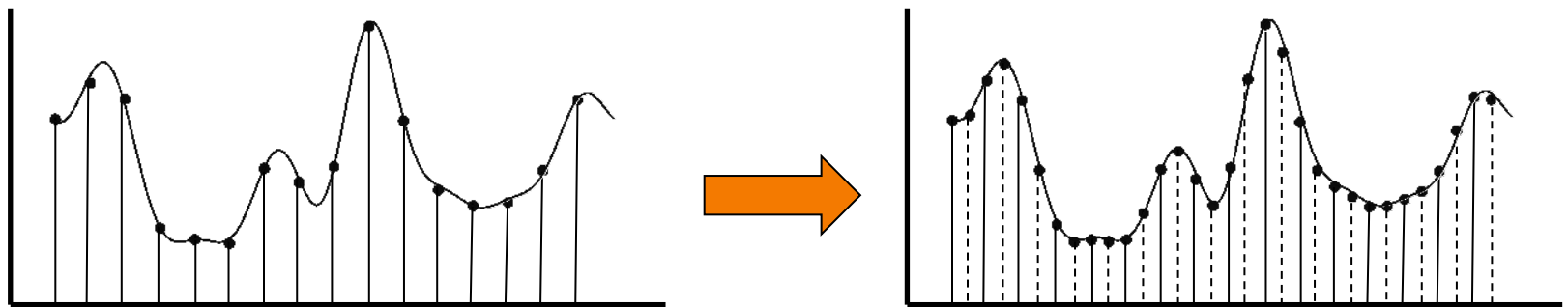
- **g**: The initial grid
- **f \* g**: The Gaussian-smoothed grid



# Up-Sampling

Given an  $n$ -dimensional array, we would like to extrapolate the array to a  $2n$ -dimensional array.

One way to do this is to fit a continuous function to the original array and then sample at  $2n$  regular samples.





# Up-Sampling

How do we generate a continuous function from a set of  $n$  samples?

Recall that the Fourier decomposition of  $\mathbf{g}$  is:

$$\mathbf{g}_j = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{\frac{ik2\pi j}{n}}$$



# Up-Sampling

How do we generate a continuous function from a set of  $n$  samples?

We can fit a continuous function to the data by associating the discrete index  $0 \leq j \leq n$  with a continuous index  $0 \leq s \leq 2\pi$ :

$$g_j = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{\frac{ik2\pi j}{n}} \Rightarrow g(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{iks}$$

Since at integer values of  $j$  we have:

$$g_j = g(2\pi j/n)$$

the continuous function interpolates the  $n$  discrete samples.



# Up-Sampling

## Word of Warning:

Recall that for integer values of  $j$ , the complex exponential satisfies the condition:

$$e^{\frac{ik2\pi j}{n}} = e^{\frac{i(k+n)2\pi j}{n}}$$

Thus, if we were to fit a function:

$$g_j = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{\frac{ik2\pi j}{n}} \Rightarrow g(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{i(k+2\pi)s}$$

we would also get a continuous function that interpolates the  $n$  discrete samples.

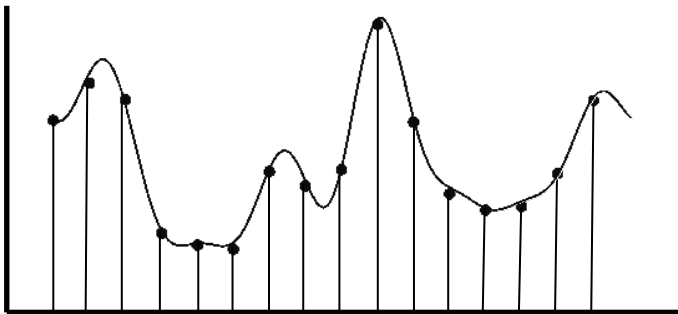


# Up-Sampling

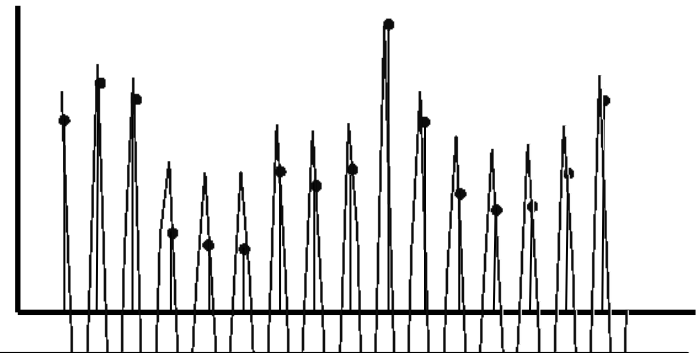
## Word of Warning:

The difference is in how the array is interpolated:

$$g(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{iks}$$



$$g(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_k \hat{g}_k \cdot e^{i(k+2\pi)s}$$



In fact, if not done carefully, a real-valued input array will not necessarily generate a real-valued up-sampled array.

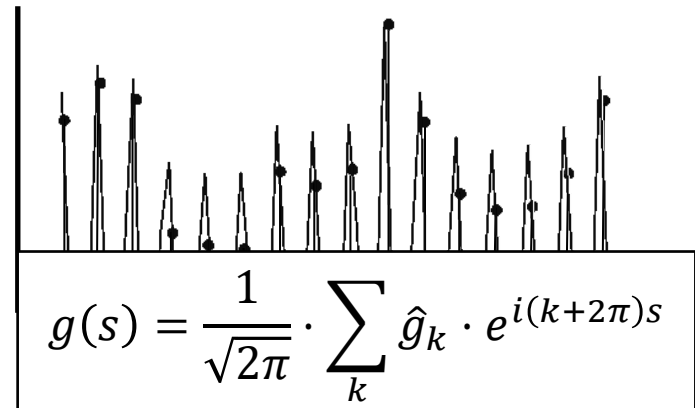
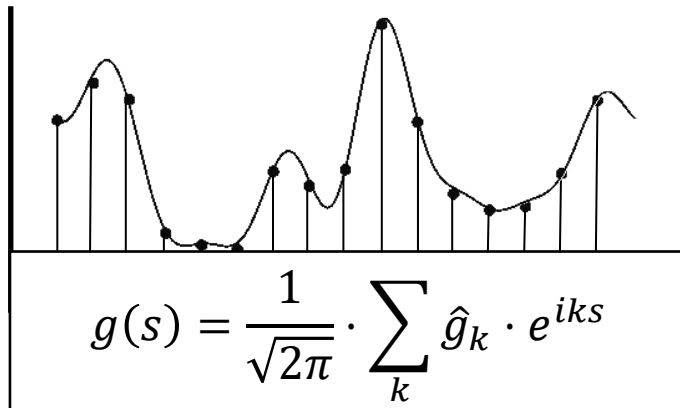


# Up-Sampling

## Word of Warning:

Extrapolating the discrete samples is an under-constrained problem, and so there are many different solutions.

In practice, we would like the “smoothest” possible fit, so we would like to minimize the contribution of high frequency terms





# Up-Sampling

## Word of Warning:

The “best” fit is obtained using the function:

$$g(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_{k=-n/2+1}^{n/2} \hat{g}_k \cdot e^{iks}$$

A simple algorithm for implementing this is:

1. Compute the  $n$  Fourier coefficients of  $\mathbf{g}$
2. Generate an array of  $2n$  Fourier coefficients:

$$\hat{h}_k = \begin{cases} \hat{g}_k & -n/2 \leq k \leq n/2 \\ 0 & \text{otherwise} \end{cases}$$

3. Compute the inverse Fourier transform to get back the  $2n$ -dimensional array  $\mathbf{h}$





# Differentiation

Given an  $n$ -dimensional array  $\mathbf{g}$ , how do we compute the derivative of  $\mathbf{g}$ ?

Finite (Central) Differences:

Define the derivative at some index  $j$  as the average of the discrete left and right derivatives:

$$\begin{aligned} g'_j &= \frac{(g_{j+1} - g_j) + (g_j - g_{j-1}))}{2} \\ &= \frac{g_{j+1} - g_{j-1}}{2} \end{aligned}$$



# Differentiation

Given an  $n$ -dimensional array  $\mathbf{g}$ , how do we compute the derivative of  $\mathbf{g}$ ?

Finite (Central) Differences:

Define the derivative at some index  $j$  as the average of the discrete left and right derivatives:

$$g'_j = \frac{g_{j+1} - g_{j-1}}{2}$$

Continuous Differentiation:

Fit a continuous function to the samples and take the derivative of the continuous function.



# Differentiation

## Continuous Differentiation:

Fit a continuous function to the samples and take the derivative of the continuous function.

Fit a continuous function to **g** by setting:

$$g(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_{k=-n/2+1}^{n/2} \hat{g}_k \cdot e^{iks}$$

Since  $\frac{\partial}{\partial s} e^{\lambda s} = \lambda e^{\lambda s}$  we get:

$$g'(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_{k=-n/2+1}^{n/2} \hat{g}_k \cdot ik \cdot e^{iks}$$



# Differentiation

Continuous Differentiation:

$$g'(s) = \frac{1}{\sqrt{2\pi}} \cdot \sum_{k=-n/2+1}^{n/2} \hat{g}_k \cdot ik \cdot e^{iks}$$

⇒ Get the values of the continuous derivative by multiplying the  $k$ -th Fourier coefficient by  $ik$ :

$$\hat{g}'_k = ik \cdot \hat{g}_k$$

Note:

As with up-sampling, there are many continuous functions we could fit to the discrete samples.

We use the smoothest one.



# Edge Detection

To compute the edges in a grid  $g$ , we would like to measure how much the grid  $g$  is changing at every index.

Specifically, we would like to define a grid  $h$  such that  $h_{jk}$  measures the “extent of change” of  $g$  at the index  $(j, k)$ .





# Edge Detection

## Gradient Method:

One way to measure the extent change is by computing the gradient of  $g$  at every point and setting:

$$h_{jk} = \|\nabla g_{jk}\|$$

To compute the gradient, we need to compute the partial derivatives of  $g$  at every point.



# Edge Detection

## Gradient Method:

This can be done with the FFT:

1. Compute the Fourier coefficients of  $\mathbf{g}$
2. Generate the two grids corresponding to the Fourier coefficients of the partial derivatives:

$$\hat{g}_{jk}^x = \hat{g}_{jk} \cdot ij$$

$$\hat{g}_{jk}^y = \hat{g}_{jk} \cdot ik$$

3. Compute the inverse Fourier transforms to get the grids of partial derivatives  $\mathbf{g}^x$  and  $\mathbf{g}^y$ .
4. Set  $\mathbf{h}$  to be the grid of gradient lengths:

$$h_{jk} = \sqrt{\left(g_{jk}^x\right)^2 + \left(g_{jk}^y\right)^2}$$



# Edge Detection

## Laplacian Method:

An alternate way to measure the rate of change is to compute the difference between the original grid and a smoothed version of the grid.

A measure of this difference can be obtained by computing the Laplacian (the sum of unmixed) partial derivatives:

$$\Delta f = f_{xx} + f_{yy}$$





# Edge Detection

## Laplacian Method:

This can be done with the FFT:

1. Compute the Fourier coefficients of  $\mathbf{g}$
2. Generate the grid corresponding to the Fourier coefficients of the Laplacian:

$$\begin{aligned}\hat{h}_{jk} &= (\hat{\mathbf{g}}^{xx} + \hat{\mathbf{g}}^{yy})_{jk} = \hat{g}_{jk} \cdot ((ij)^2 + (ik)^2) \\ &= -\hat{g}_{jk} \cdot (j^2 + k^2)\end{aligned}$$

3. Set  $\mathbf{h}$  to be the inverse Fourier transform of the Laplacian Fourier coefficients.



# Gaussian Sharpening

How do we undo the effects of Gaussian-smoothing a grid?

We perform Gaussian smoothing of  $g$  by:

1. Computing the Fourier transforms of  $g$  and the Gaussian filter  $f$
2. Multiplying the Fourier coefficients of  $g$  by the Fourier coefficients of  $f$
3. Computing the inverse Fourier transform



# Gaussian Sharpening

How do we undo the effects of Gaussian-smoothing a grid?

We un-perform Gaussian smoothing of  $g$  by:

1. Computing the Fourier transforms of  $g$  and the Gaussian grid  $f$
2. Multiplying the Fourier coefficients of  $g$  by the reciprocals of the Fourier coefficients of  $f$
3. Computing the inverse Fourier transform

This is doable as long as the Fourier coefficients of  $f$  are non-zero.

This is numerically stable as long as the Fourier coefficients of  $f$  are not too small.



# Outline

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**Real FFTs**



# Real FFTs

So far, we have considered the Fourier transform of complex-valued functions. What happens when the values of the function are all real?

$$g(\theta) - \overline{g(\theta)} = 0$$



# Real FFTs

If we write out the function  $g$  in terms of its Fourier decomposition, we get:

$$g(\theta) = \frac{1}{\sqrt{2\pi}} \cdot \sum_{k=-\infty}^{\infty} \hat{g}_k \cdot e^{ik\theta}$$

Using the fact that  $g(\theta) - \bar{g}(\theta) = 0$  we get:

$$0 = \frac{1}{\sqrt{2\pi}} \cdot \left( \sum_{k=-\infty}^{\infty} \hat{g}_k \cdot e^{ik\theta} - \sum_{k=-\infty}^{\infty} \bar{\hat{g}}_k \cdot e^{-ik\theta} \right)$$

This can be re-written as:

$$0 = \frac{1}{\sqrt{2\pi}} \cdot \left( \sum_{k=-\infty}^{\infty} \hat{g}_k \cdot e^{ik\theta} - \sum_{k=-\infty}^{\infty} \bar{\hat{g}}_{-k} \cdot e^{ik\theta} \right)$$



# Real FFTs

$$0 = \frac{1}{\sqrt{2\pi}} \cdot \left( \sum_{k=-\infty}^{\infty} \hat{g}_k \cdot e^{ik\theta} - \sum_{k=-\infty}^{\infty} \bar{\hat{g}}_{-k} \cdot e^{ik\theta} \right)$$

Simplifying, we get:

$$0 = \sum_{k=-\infty}^{\infty} (\hat{g}_k - \bar{\hat{g}}_{-k}) \cdot e^{ik\theta}$$

But since complex exponentials are linearly independent, this implies that all the Fourier coefficients are equal to zero:

$$\begin{aligned} 0 &= \hat{g}_k - \bar{\hat{g}}_{-k} \\ &\quad \Updownarrow \\ \hat{g}_k &= \bar{\hat{g}}_{-k} \end{aligned}$$



# Real FFTs

$$\hat{g}_k = \bar{\hat{g}}_{-k}$$

⇒ When the function  $g$  is real-valued, the  $k$ -th and  $(-k)$ -th Fourier coefficient are conjugate.





# Real FFTs

Although this discussion is true for real functions, a similar argument shows that for real-valued arrays  $\mathbf{g} \in \mathbb{R}^n$ , the Fourier coefficients have the property that:

$$\hat{g}_k = \bar{\hat{g}_{n-k}}$$

In particular:

1. The dc term,  $\hat{g}_0$  must be real, and
2. If  $n$  is even,  $\hat{g}_{n/2}$  must be real as well.



# Real FFTs

For real-valued arrays:

- In 1D we have:

$$\hat{g}_j = \bar{\hat{g}}_{n-j}$$

- In 2D we have:

$$\hat{g}_{jk} = \bar{\hat{g}}_{n-j, n-k}$$

- In 3D we have:

$$\hat{g}_{jkl} = \bar{\hat{g}}_{n-j, n-k, n-l}$$

⇒ When the array is real-valued, we only have to compute and store half of the coefficients.