

Search and Intersection

O'Rourke, Chapter 7

de Berg et al., Chapter 11

Announcements



- Assignment 3 web-page has been updated:
 - Additional extra credit
 - Hints for managing a dynamic half-edge representation

Outline



- Review
 - Duality
 - Linear Programming
- Half-Spaces and Convex Hulls (2D)

Duality



Definition:

Given a point $p = (\alpha, \beta)$ in the plane, define the dual line to be the (non-vertical) line with equation: $p^* = \{(x, y) | y = 2\alpha x - \beta\}$

Given a line $L = \{(x,y)|y = mx + b\}$, define the *dual* point to be the point with coordinates:

$$L^* = \left(\frac{m}{2}, -b\right)$$

Duality



Properties:

Given a point p and lines L, L_1 , and L_2 :

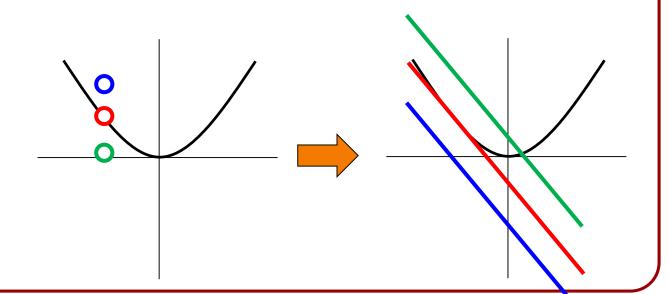
- $(p^*)^* = p$ and $(L^*)^* = L$.
- $p \in L$ iff. $L^* \in p^*$.
- $p \in L_1 \cap L_2$ iff. $L_1^*, L_2^* \in p^*$.
- L is below/above p iff. L^* is above/below p^* .
- p is on the parabola $y = x^2$ iff. p^* is tangent to the parabola at p.

Duality



Properties:

- Given a point $p = (\alpha, \beta)$:
 - The slope of p^* is the slope of the tangent to the parabola at (α, α^2) .
 - p^* passes through the point $(\alpha, \alpha^2 + (\alpha^2 \beta))$.



Linear Programming



Goal:

Given a set of linear constraints:

$$C_i = \{p | \langle p, n_i \rangle \ge d_i \}$$

and a linear energy function:

$$E(p) = \langle p, n \rangle + d$$

we would like to find the point p that satisfies the constraints and minimizes the energy.

Linear Programming



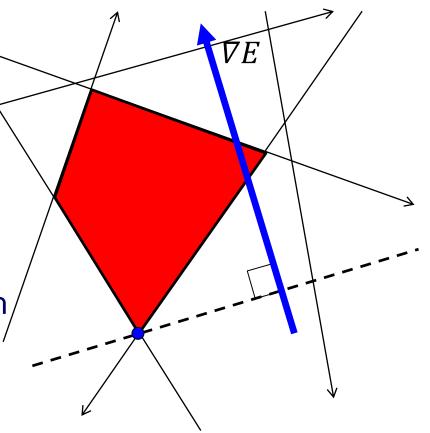
Approach:

 Since the constraints are linear, each one defines a half-space of valid solutions.

 The intersection of these half-spaces is convex.

 Since the energy is linear, it has a constant gradient \(\nabla E \) pointing away from the minimum.

 The minimizer is the point in the convex region which is extreme along direction n.



Linear Programming



Approach:

 Since the constraints are linear, each one defines a half-space of valid solutions.

 The intersection of these half-spaces is convex.

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How do we compute the convex hull corresponding to the intersection of half spaces?

Outline



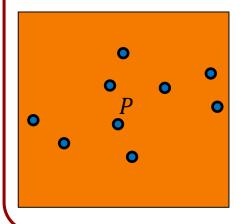
- Review
- Half-Spaces and Convex Hulls (2D)

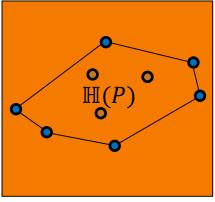


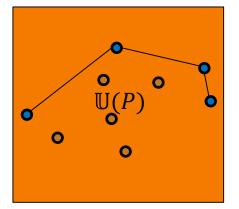
Notation:

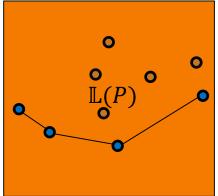
Given a set of points, $P = \{p_1, ..., p_n\}$:

- Denote the convex hull of the points as: $\mathbb{H}(P)$
- Denote the upper hull of the points as: $\mathbb{U}(P)$
- Denote the lower hull of the points as: $\mathbb{L}(P)$





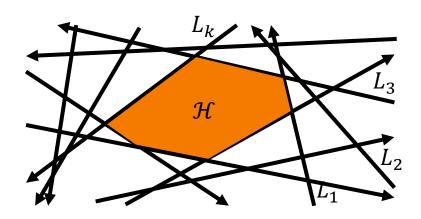






Goal:

Given a set of half-spaces, represented by directed lines $\{L_1, ..., L_n\}$ compute the convex hull, \mathcal{H} , corresponding to the boundary of their intersection.

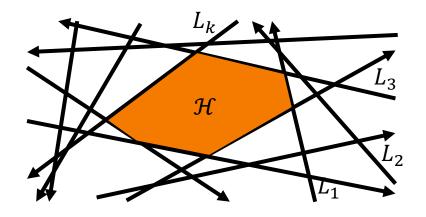


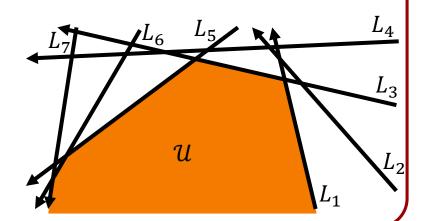


Approach:

Consider the upper/lower hulls, \mathcal{U}/\mathcal{L} , independently.

∘ For the upper (resp. lower) part, assume all line segments have $(0, -\infty)$ (resp. $(0, \infty)$) to their left.







Approach:

Consider the upper/lower hulls, \mathcal{U}/\mathcal{L} , independently.

An edge $e \subset \mathcal{U}$ is the set of points on a line that are below all the other lines:

$$e \subset L_i$$
 and Below $(v, L_k) \forall k \neq i, v \in e$

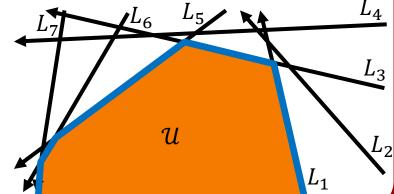
Dually:

$$L_i^* \in v^*$$
 and Below $(L_k^*, v^*) \ \forall k \neq i, v \in e$

 \Leftrightarrow Lines v^* , with $v \in e$, pass through L_i^* and have all other $\{L_1^*, \dots, L_n^*\}$ below.

 $\Leftrightarrow L_i^*$ is a vertex of:

$$\mathcal{U}^* = \mathbb{U}(L_1^*, \dots, L_n^*)$$





Approach:

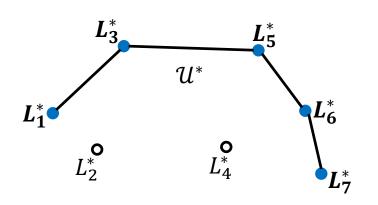
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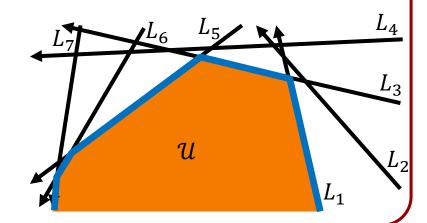
An edge $e \subset \mathcal{U}$ is the set of points on a line that are below all the other lines:

$$e \subset L_i$$
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Dually:

$$L_i^* \in v^*$$
 and Below $(L_k^*, v^*) \ \forall k \neq i, v \in e$







Approach:

Consider the upper/lower hulls, \mathcal{U}/\mathcal{L} , independently.

A vertex $v \in \mathcal{U}$ lies on the intersection of two lines and is below all the other lines:

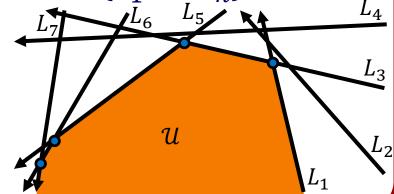
$$v \in L_i \cap L_j$$
 and Below $(v, L_k) \forall k \neq i, j$

<u>Dually</u>:

$$L_i^*, L_j^* \in v^*$$
 and $Below(L_k^*, v^*) \ \forall k \neq i, j$

 \Leftrightarrow The line segment $\overline{L_i^*L_i^*}$ has all other $\{L_1^*, ..., L_n^*\}$ below.

$$\Leftrightarrow \overline{L_i^* L_j^*}$$
 is an edge of $\mathcal{U}^* = \mathbb{U}(L_1^*, ..., L_n^*)$





Approach:

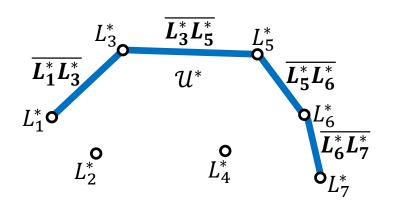
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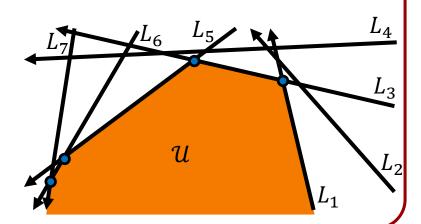
A vertex $v \in \mathcal{U}$ lies on the intersection of two lines and is below all the other lines:

$$v \in L_i \cap L_j$$
 and Below $(v, L_k) \forall k \neq i, j$

Dually:

$$L_i^*, L_i^* \in v^*$$
 and Below $(L_k^*, v^*) \forall k \neq i, j$





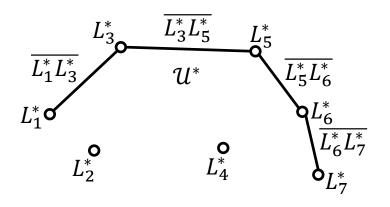


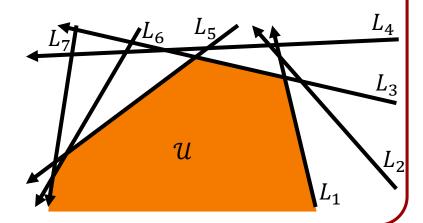
Implementation:

UpperHull($\{L_1, ..., L_n\}$)

- ∘ $\{v_1, ..., v_m\}$ ← UpperHull($\{L_1^*, ..., L_n^*\}$)
- return $\{v_m^*, ..., v_1^*\}$

This gives the lines in the (CCW) order in which they appear on the intersection of half-spaces.







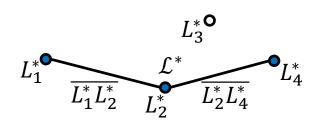
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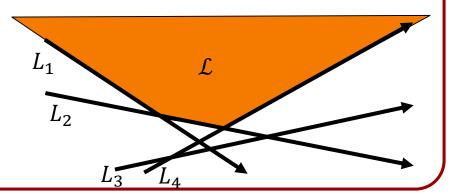
UpperHull($\{L_1, ..., L_n\}$)

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LowerHull($\{L_1, ..., L_n\}$)

- ∘ $\{v_1, ..., v_m\}$ ← LowerHull($\{L_1^*, ..., L_n^*\}$)
- return $\{v_1^*, ..., v_m^*\}$







Implementation:

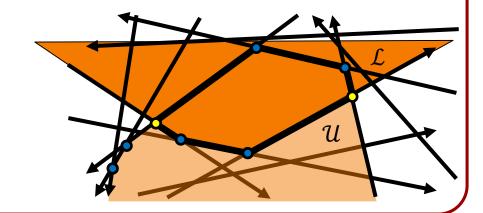
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UpperHull( \{L_1, ..., L_n\} )
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- ∘ $\{v_1, ..., v_m\}$ ← UpperHull($\{L_1^*, ..., L_n^*\}$)
- return $\{v_m^*, ..., v_1^*\}$

LowerHull($\{L_1, ..., L_n\}$)

- ∘ $\{v_1, ..., v_m\}$ ← LowerHull($\{L_1^*, ..., L_n^*\}$)
- \circ return $\{v_1^*, \dots, v_m^*\}$

Taking the intersection (in linear time), we get the convex hull.





<u>Implementation</u>:

We have to separately compute the upper and lower hulls because the dual map is undefined (discontinuous) as lines approach vertical.

Is there an alternate definition of duality that would allow us to compute the hulls simultaneously?

No! The convex hull of the dual cannot be empty, but the intersection of half-spaces can be.

(in linear time), we get the convex hull.



Writing $p \in \mathbb{R}^n$ as:

$$p = (\alpha_1, \dots, \alpha_{n-1}, \beta) = (\vec{\alpha}, \beta)$$

we can define duality in n-dimensional space.

Given a point $p = (\vec{\alpha}, \beta)$, define the *dual hyper-plane* to be the (non-vertical) plane:

$$p^* = \{(\vec{x}, y) | y = 2\langle \vec{x}, \vec{\alpha} \rangle - \beta\}$$

Given a plane $H = \{(\vec{x}, y) | y = \langle \vec{m}, \vec{x} \rangle + b\}$, define the *dual point* to be the point with coordinates:

$$H^* = \left(\frac{\overrightarrow{m}}{2}, -b\right)$$



Writing $p \in \mathbb{R}^n$ as:

$$n = (\alpha_1 \quad \alpha_{--} + R) = (\vec{\alpha} R)$$

The same properties hold in n-dimensions:

- $(p^*)^* = p$ and $(H^*)^* = H$.
- $p \in H \text{ iff. } H^* \in p^*.$
- $p \in H_1 \cap H_2 \text{ iff. } H_1^*, H_2^* \in p^*.$
- H is below/above p iff. H^* is above/below p^* .
- $p = (\vec{x}, y)$ is on the parabola $y = ||\vec{x}||^2$ iff. p^* is tangent to the parabola at p.

dual point to be the point with coordinates:

$$H^* = \left(\frac{\overline{m}}{2}, -b\right)$$



<u>In 3D</u>:

Let $\{H_1, ..., H_m\}$ be oriented hyper-planes in \mathbb{R}^3 with $(0, 0, -\infty)$ to the left and let \mathcal{U} be the intersection of the associate half-spaces.

$$\{H_i^*, H_j^*, H_k^*\}$$
 is a triangle of $\mathbb{U}(\{H_1^*, \dots, H_m^*\})$

 $H_i \cap H_j \cap H_k$ is a vertex of \mathcal{U} .



<u>In 3D</u>:

Let $\{H_1, ..., H_m\}$ be oriented hyper-planes in \mathbb{R}^3 with $(0, 0, -\infty)$ to the left and let \mathcal{U} be the intersection of the associate half-spaces.

$$\{H_i^*, H_j^*\}$$
 is an edge of $\mathbb{U}(\{H_1^*, \dots, H_m^*\})$

 $H_i \cap H_i$ contains one edge of \mathcal{U} .



<u>In 3D</u>:

Let $\{H_1, ..., H_m\}$ be oriented hyper-planes in \mathbb{R}^3 with $(0, 0, -\infty)$ to the left and let \mathcal{U} be the intersection of the associate half-spaces.

$$H_i^*$$
 is a vertex of $\mathbb{U}(\{H_1^*, ..., H_m^*\})$

 H_i contains one face of \mathcal{U} .



More Generally:

Let $\{H_1, ..., H_m\}$ be oriented hyper-planes in \mathbb{R}^n with $(0, ..., 0, -\infty)$ to the left and let \mathcal{U} be the intersection of the associate half-spaces.

$$\{H_{i_1}^*, \dots, H_{i_k}^*\}$$
 is a k -simplex of $\mathbb{U}(\{H_1^*, \dots, H_m^*\})$

 $\bigcap_{j=1}^k H_{i_j}$ contains one (n-d)-dimensional face of \mathcal{U} .