Arrangements

O'Rourke, Chapter 6

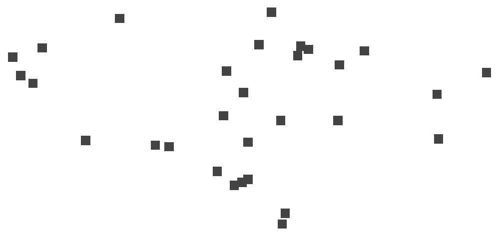
Outline

- Voronoi Diagrams
- Arrangements

Voronoi Diagrams

Recall:

We can compute the Delaunay Triangulation by raising the points to a paraboloid and computing the projection of the lower hull.



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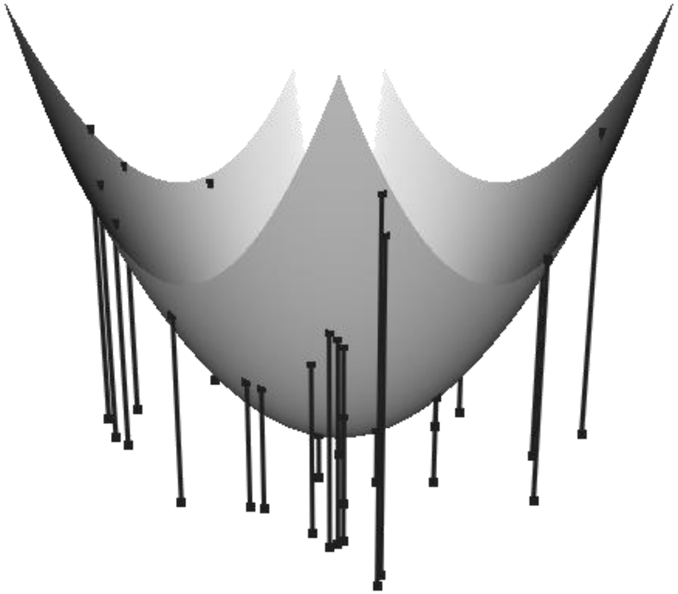

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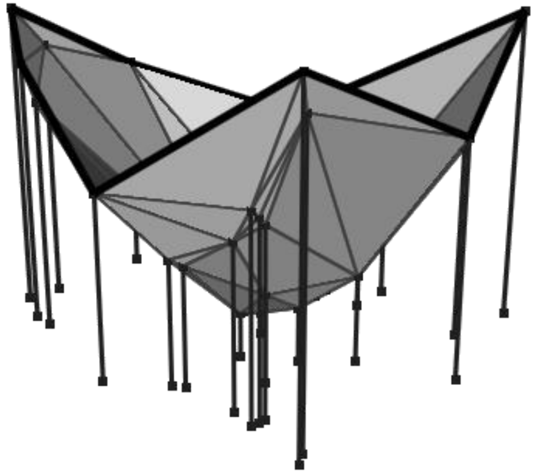

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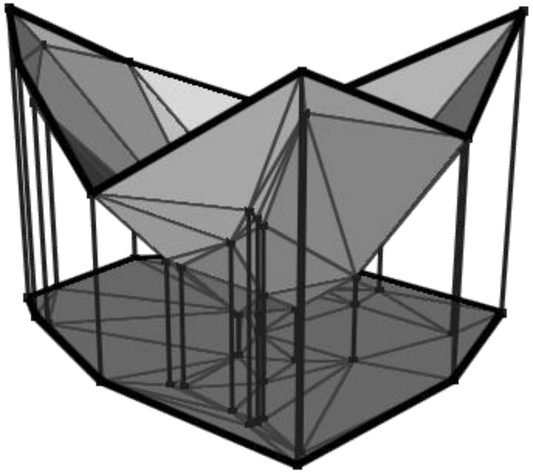

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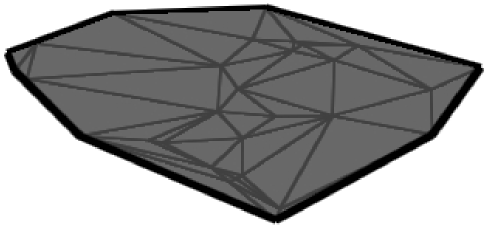


Voronoi Diagrams

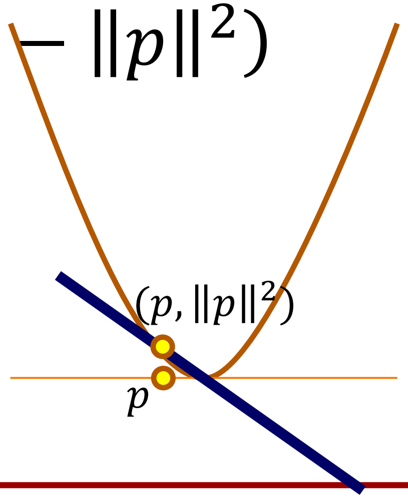
<u>Recall</u>:

Given a point $P(p) = (p, \|p\|^2)$ on the paraboloid, the tangent plane is given by:

$$z_p(r) = 2\langle p, r\rangle - \|p\|^2$$

For any point q the (vertical) distance between the points on the parabola and the tangent plane are:

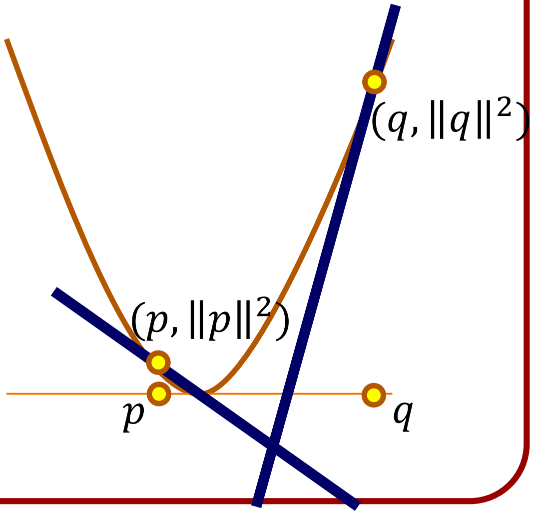
$$\begin{aligned} P(q) - z_p(q) &= \|q\|^2 - (2\langle q, p\rangle - \|p\|^2) \\ &= \|p - q\|^2 \end{aligned}$$



Voronoi Diagrams

⇒ Given points p and q, wherever the tangent plane at p is higher than the tangent plane at q, we are closer to p than to q.

$$z_p(r) \geq z_q(r)$$
$$\Updownarrow$$
$$P(r) - z_p(r) \leq P(r) - z_q(r)$$
$$\Updownarrow$$
$$\|p - r\|^2 \leq \|q - r\|^2$$



Voronoi Diagrams

⇒ Given points p and q, wherever the tangent plane at p is higher than the tangent plane at q, we are closer to p than to q.

⇒ We can visualize the Voronoi diagram by drawing the tangent planes at the sites and looking down the z-axis.

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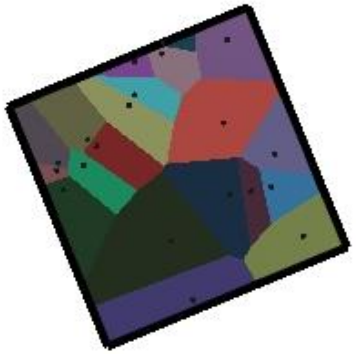

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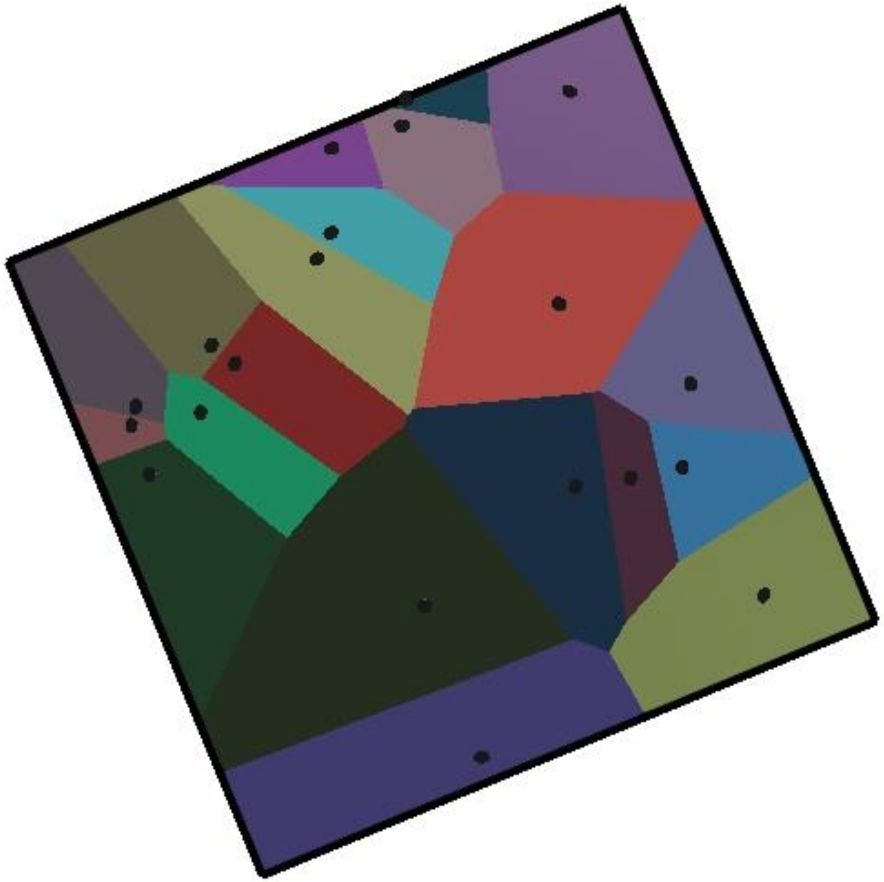


Outline

- Voronoi Diagrams
- Arrangements

Arrangements

Definition:

An *arrangement of lines* is a set of lines in the plane, inducing a partition of the domain into (convex) faces, edges, and vertices.

An arrangement is *simple* if all pairs of lines intersect, and no three lines intersect at the same point.

Combinatorics

<u>Claim</u>:

A simple arrangement of *n* lines has

- $\binom{n}{2}$ vertices,
- n^2 edges, and
- $\binom{n}{2} + n + 1$ faces.

Combinatorics

Proof (Vertices):

Since each pair of lines intersects exactly once, the total number of vertices is the number of distinct line pairs, $\binom{n}{2}$.

Combinatorics

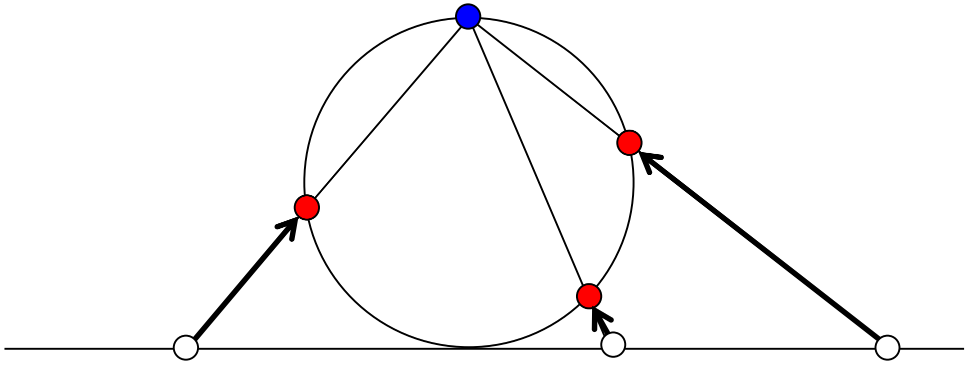
Proof (Edges):

Since each line is intersected by $n-1$ other lines, partitioning the lines into n edges, the total number of edges is n^2.

Combinatorics

Proof (Faces):

Using stereographic mapping, arrangements of lines in the plane can be thought of as polygonizations of the sphere.



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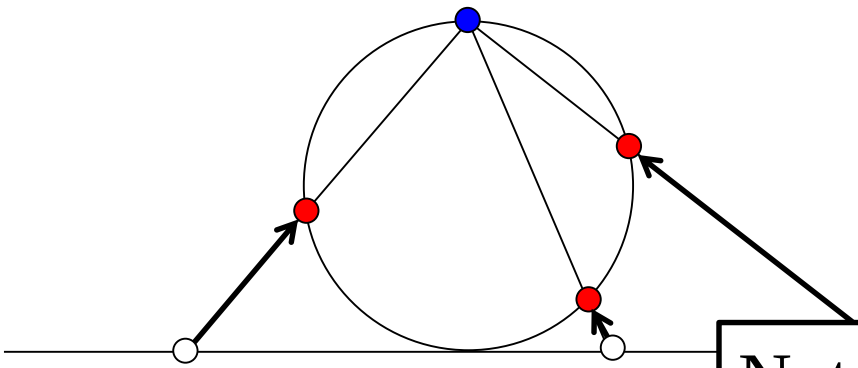
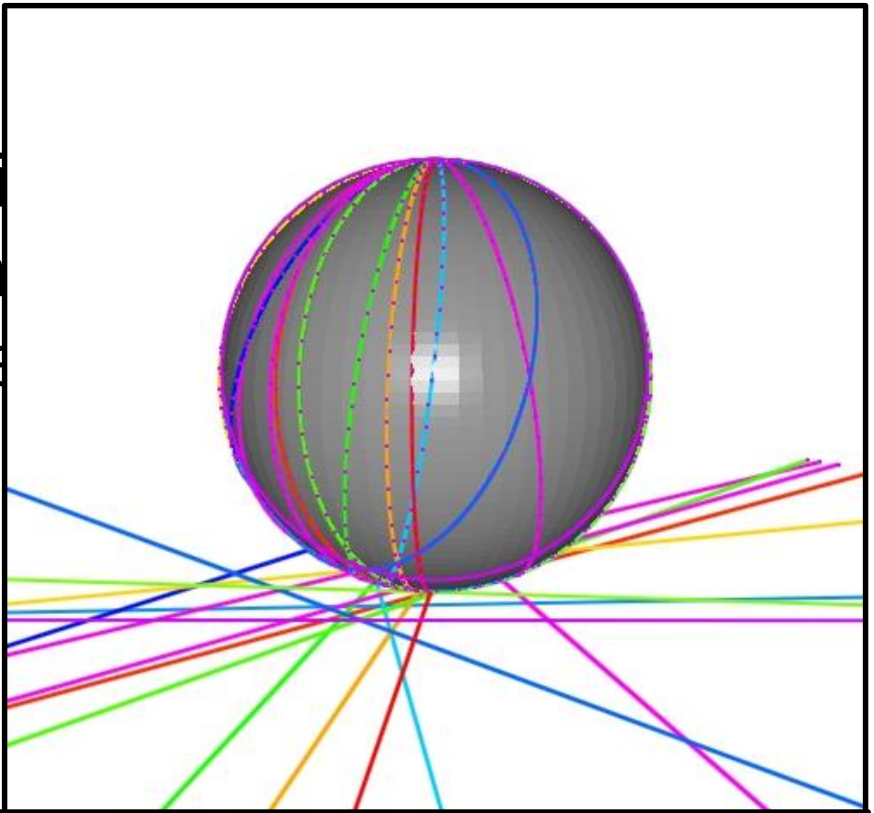




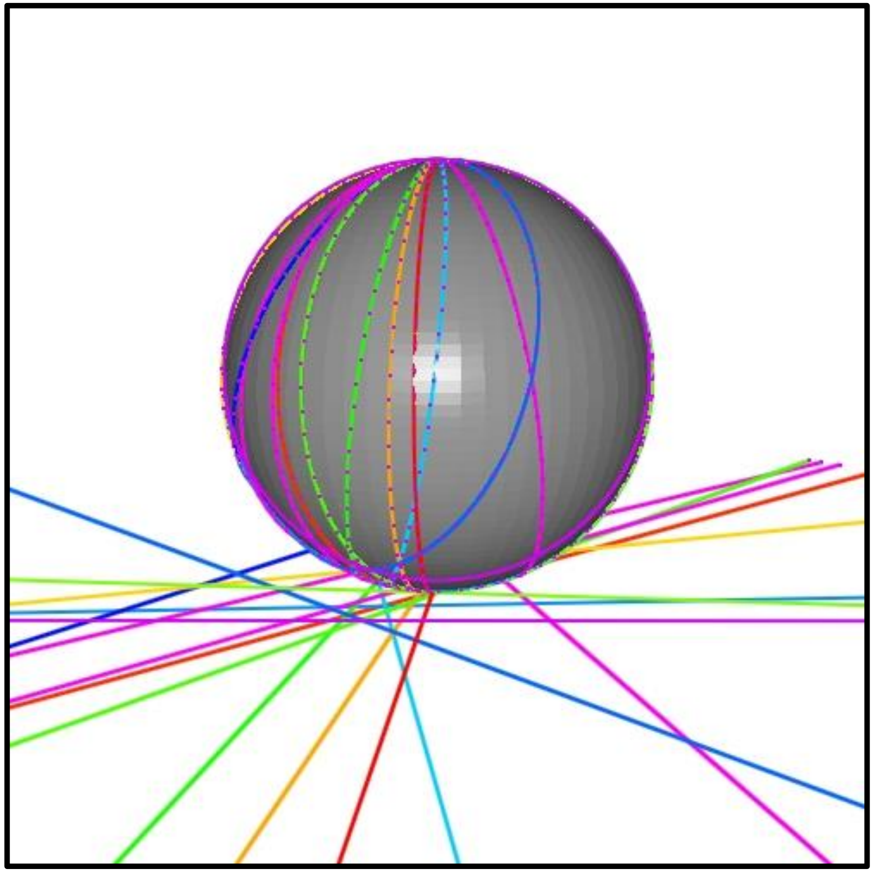
Note:
The stereographic mapping of the lines intersect at the North Pole.

Combinatorics

Proof (Faces):

By Euler's theorem the number of faces is:

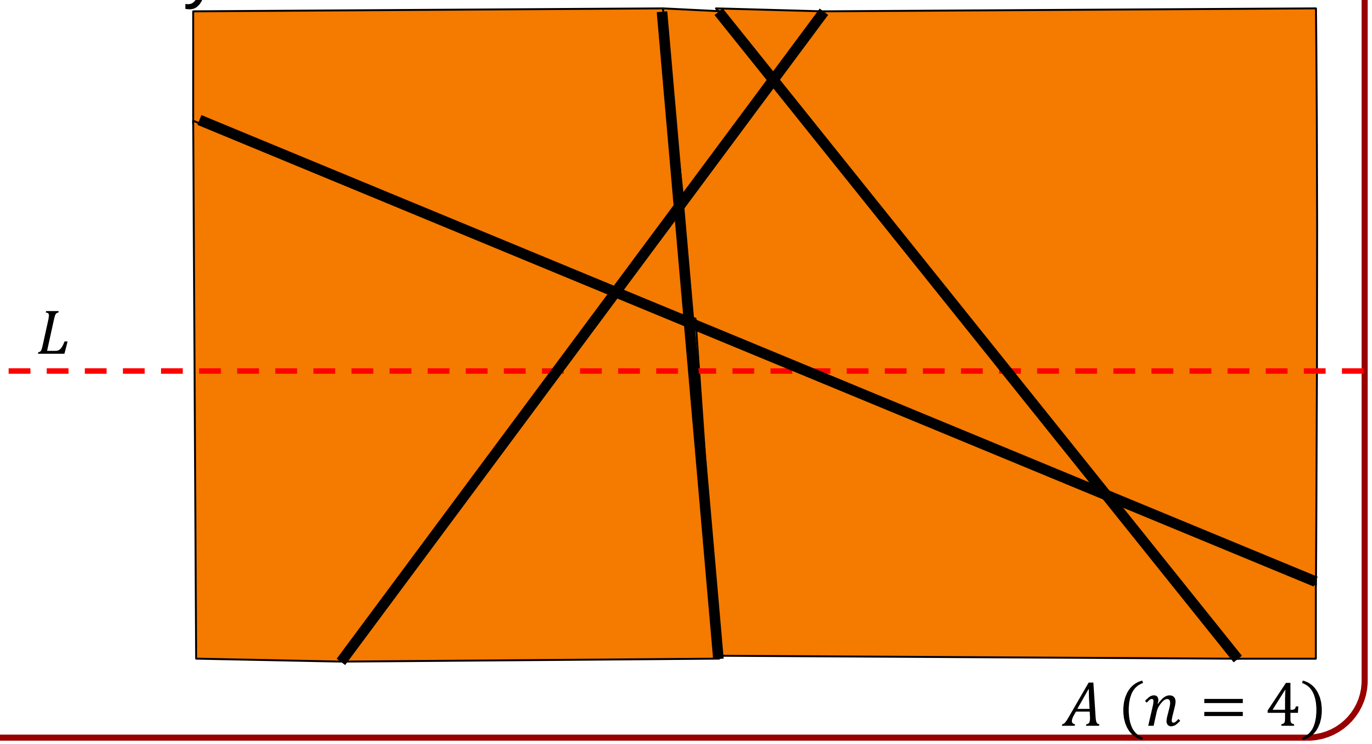
$$F = 2 - (V + 1) + E$$
$$= 2 - \binom{n}{2} - 1 + n^2$$
$$= \binom{n}{2} + n + 1$$



Zone Theorem

Definition:

Given an arrangement A and a line L (s.t. $A \cup \{L\}$ is simple) the *zone* of L in A, $Z(L)$, is the set of faces of A intersected by L.

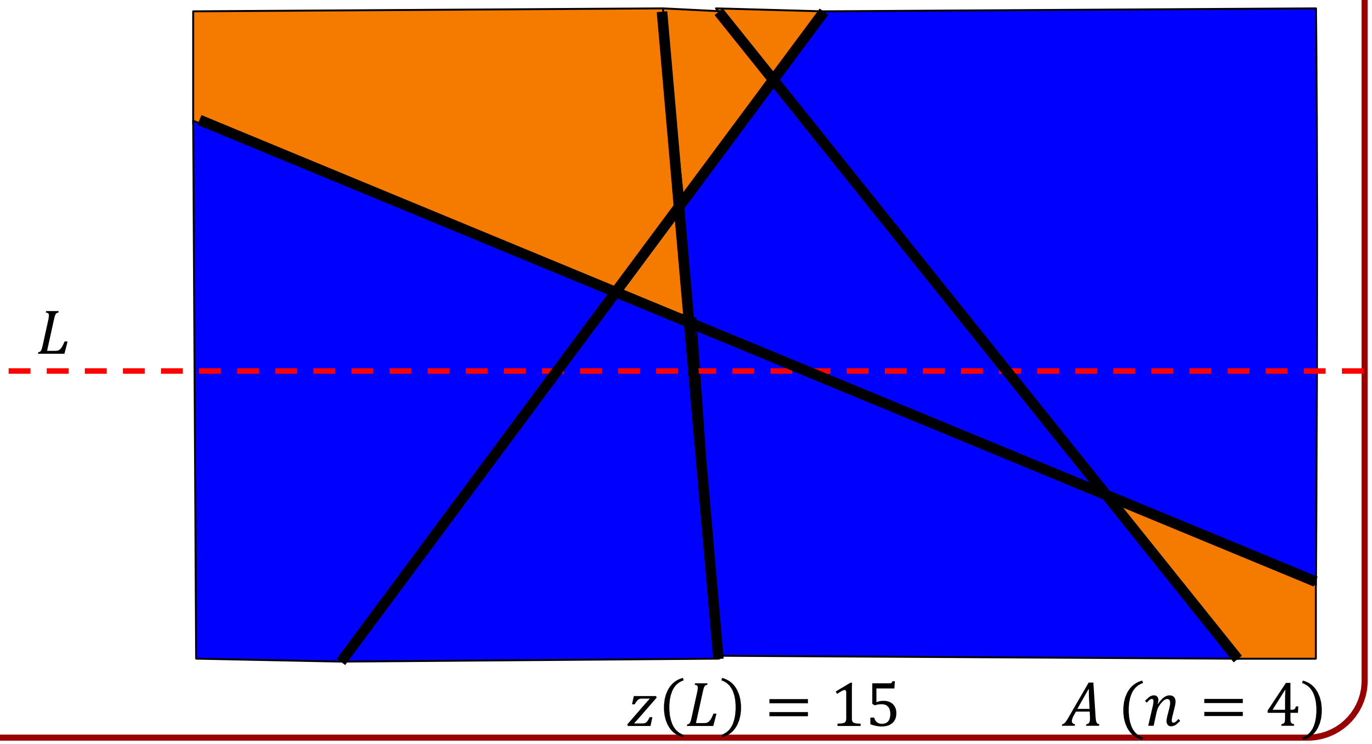


Zone Theorem

Notation:

The number of edges in $Z(L)$ is denoted $z(L)$.

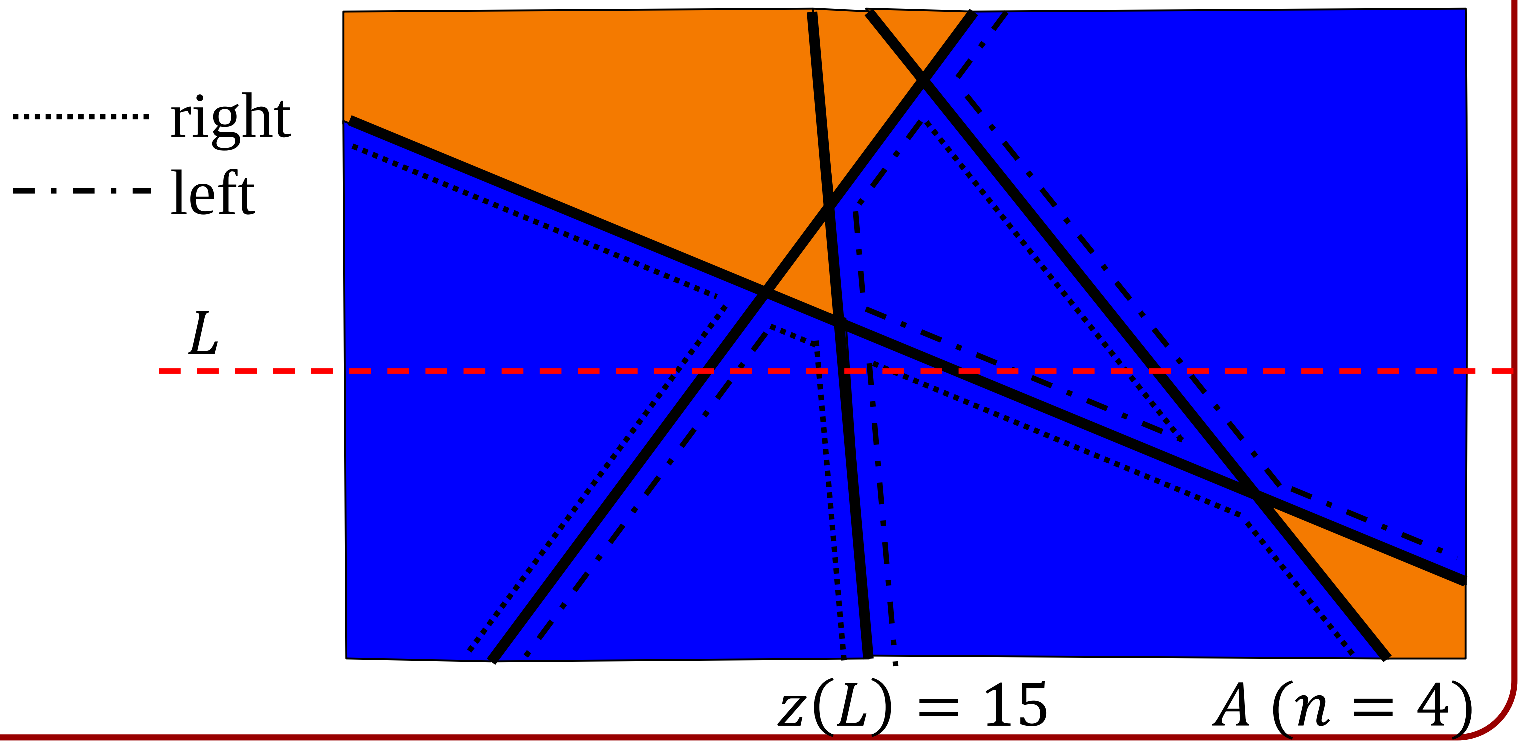
The max size of $z(L)$ over all lines is denoted z_n.



Zone Theorem

Note:

Assuming that no line in the arrangement is horizontal, we can mark each edge as left/right.

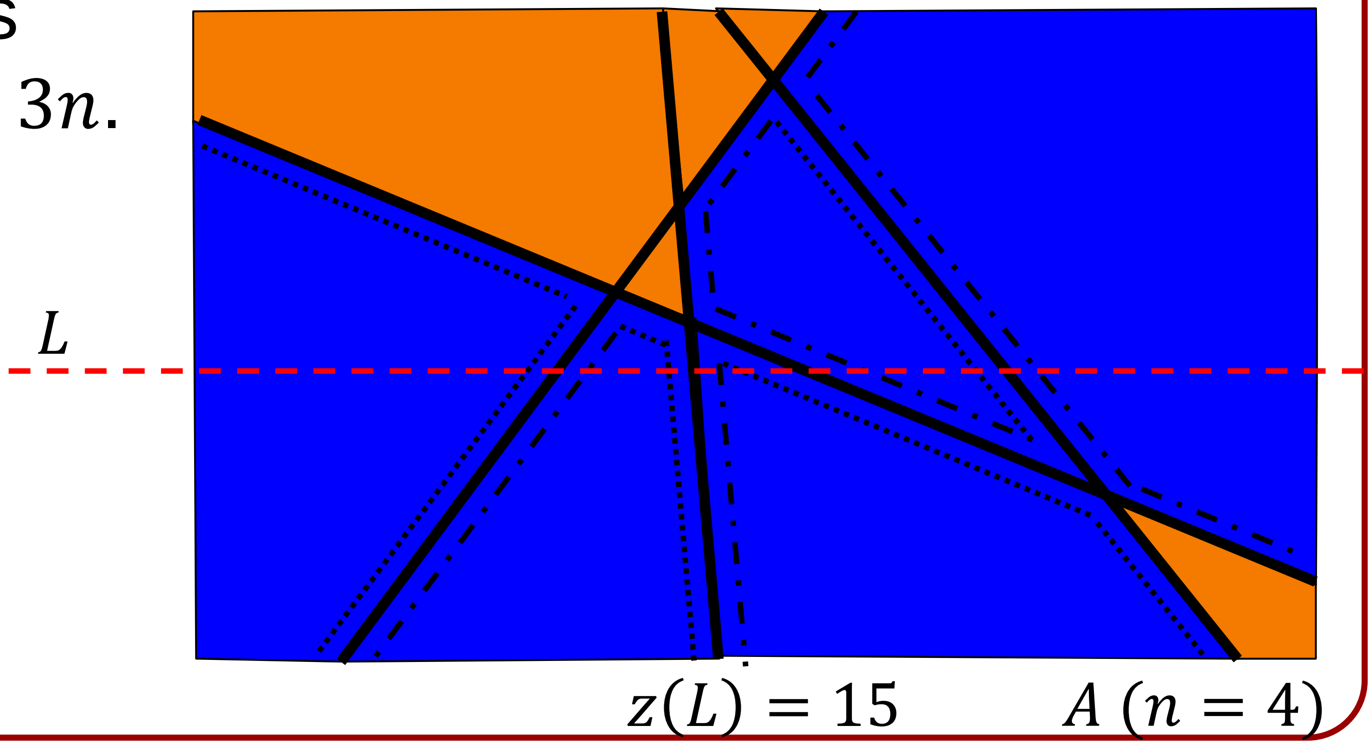


Zone Theorem

Claim:

For an arrangement of n lines, $2n \leq z_n \leq 6n$.

Specifically, the number of left/right edges crossed by a line L is bounded by $3n$.



Zone Theorem

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Note:

The expected complexity is $\sqrt{|E|} = O(n)$.

This means that we can't do worse than average.

Fo

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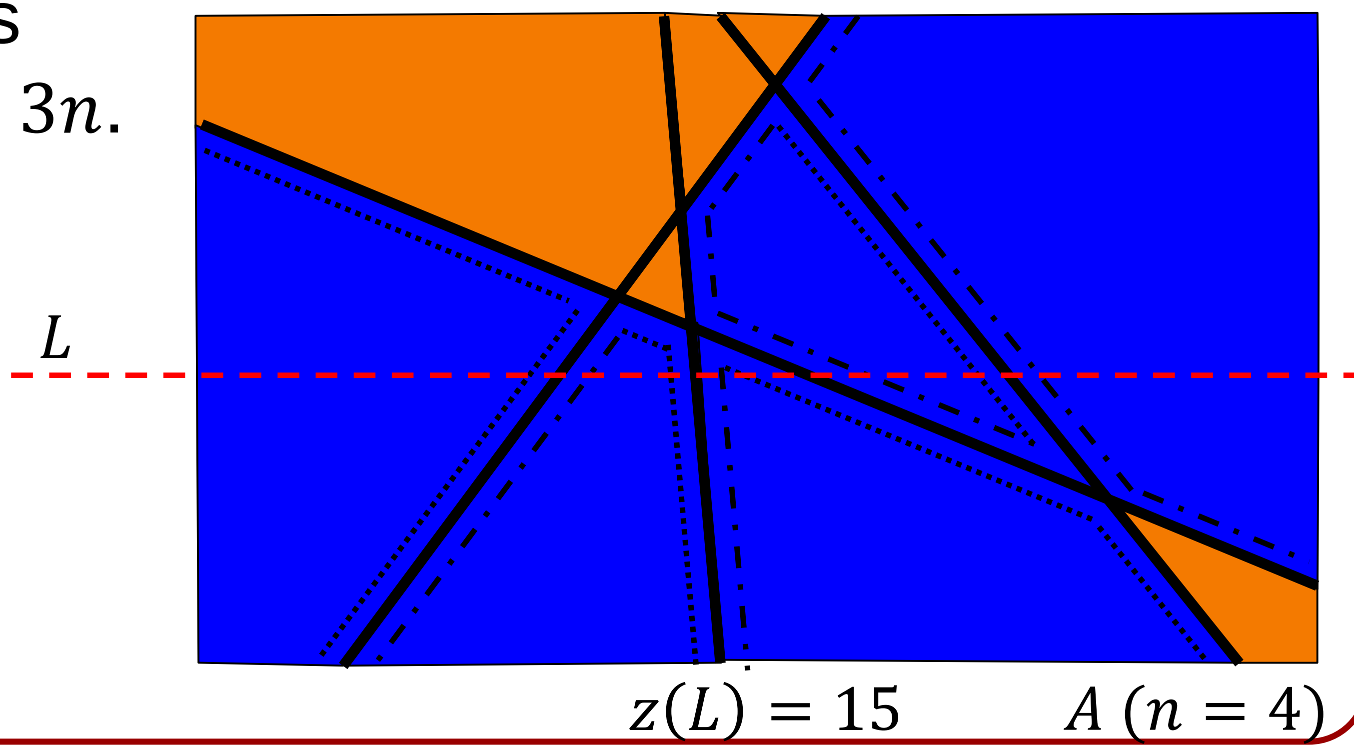


Zone Theorem

Proof:

Without loss of generality, assume that the line L maximizing the number of edges is horizontal.

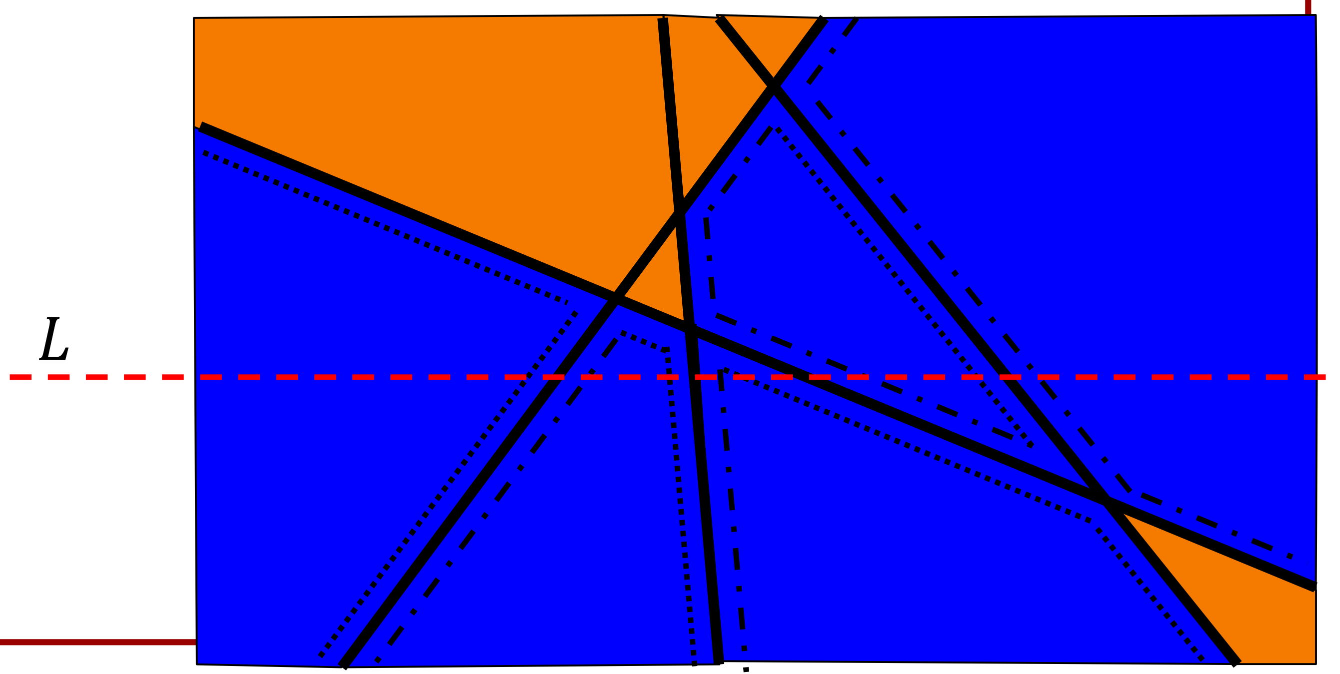
With a slight loss of generality, assume that none of the lines are vertical.

Proceed by induction.

Zone Theorem

Proof (base case):

Trivially true when $n = 0$.



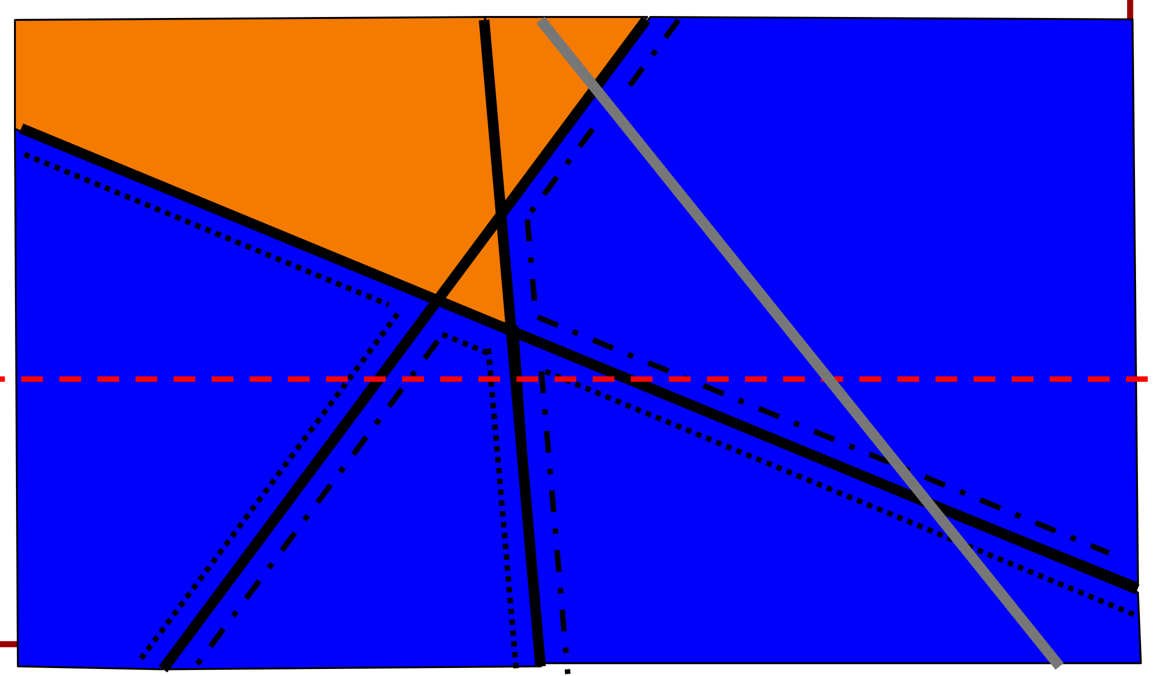
Zone Theorem

Proof (inductive case):

Remove the right-most line.

By induction, the number of left edges crossed is less than or equal to $3(n-1)$.

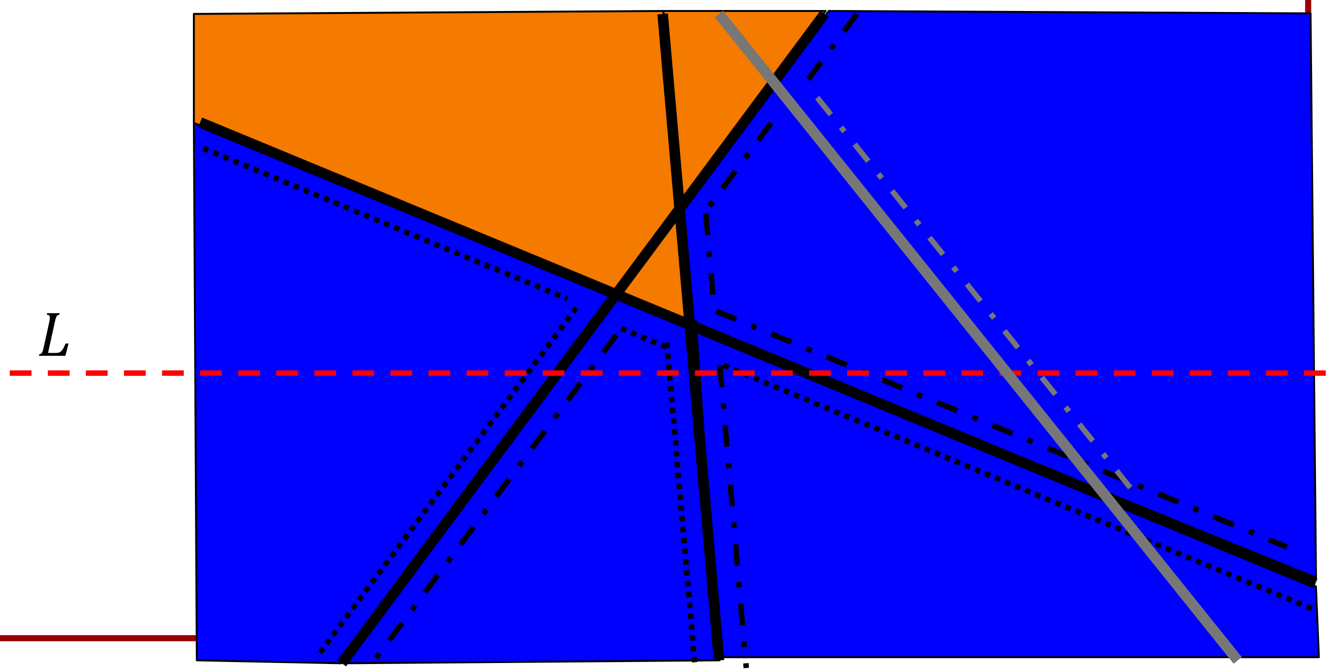
Need to show that adding the line back contributes at most 3 new L edges.



Zone Theorem

Claim:

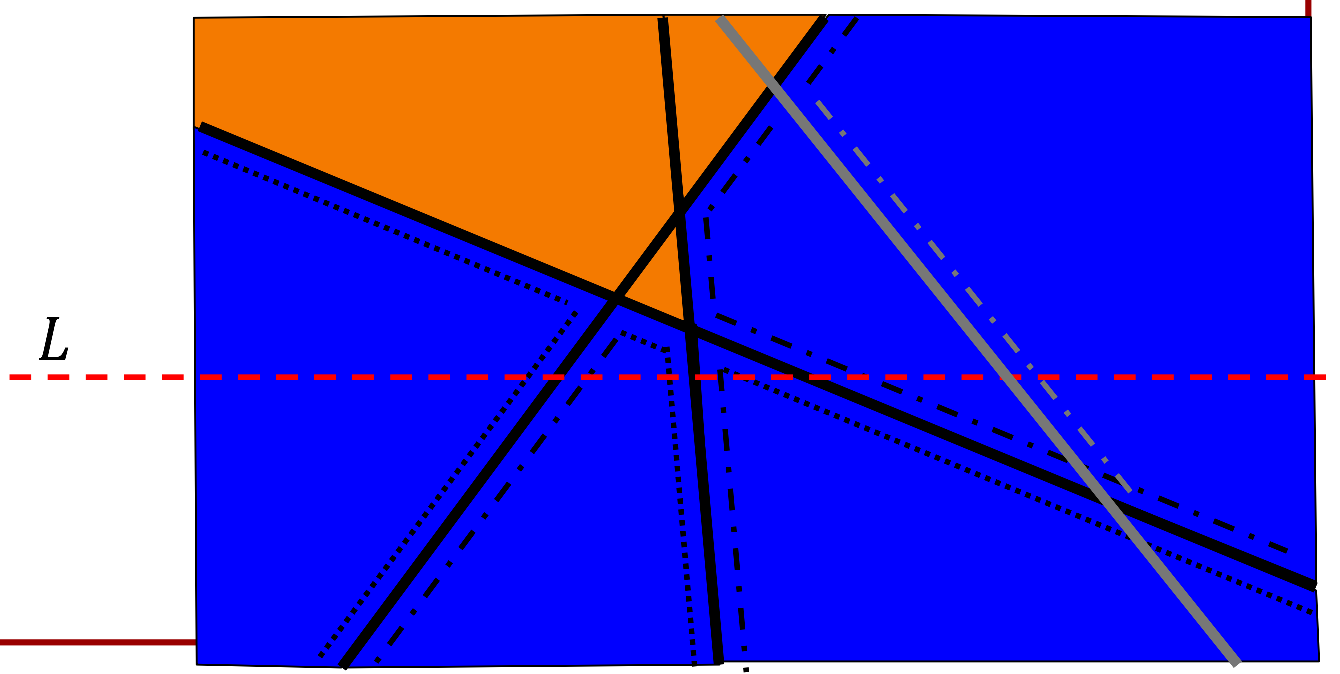
Adding the right-most line introduces exactly one new left edge.



Zone Theorem

Proof of Claim:

It introduces one because this will be a left edge of the right-most face.

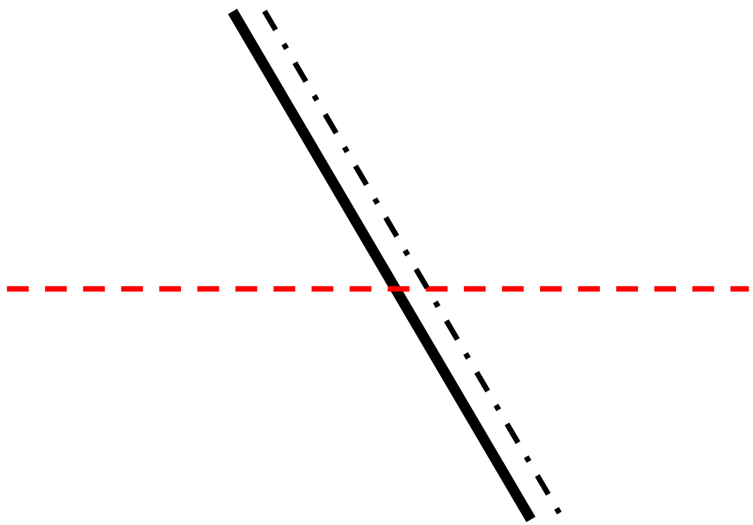


Zone Theorem

Proof of Claim:

It introduces one because this will be a left edge of the right-most face.

Exactly one because a right-most line cannot contribute more than one edge.



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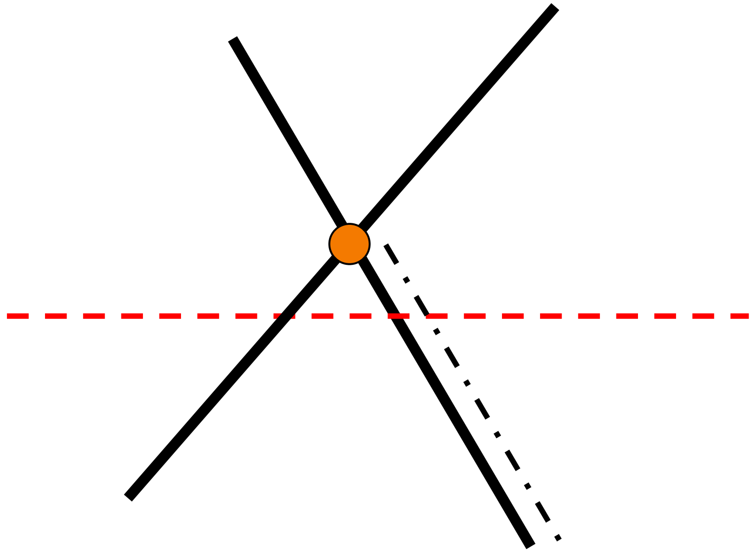

If it is split by a line from the left, only one of the two segments will be in the zone, (the one containing *L*.)

Zone Theorem

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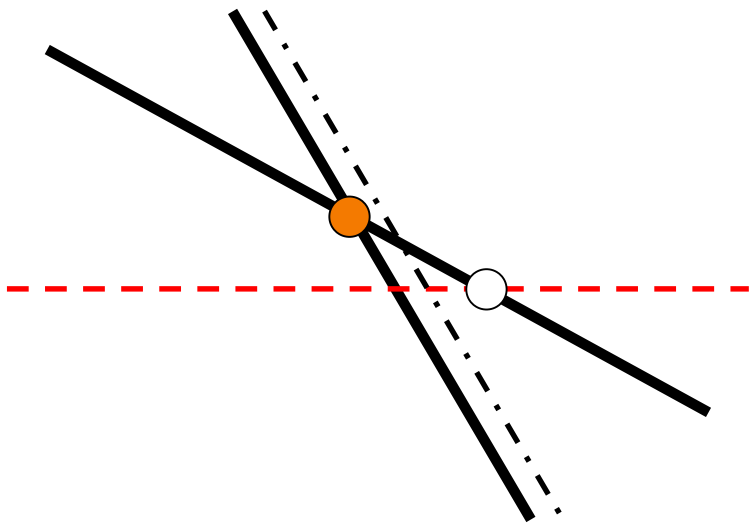
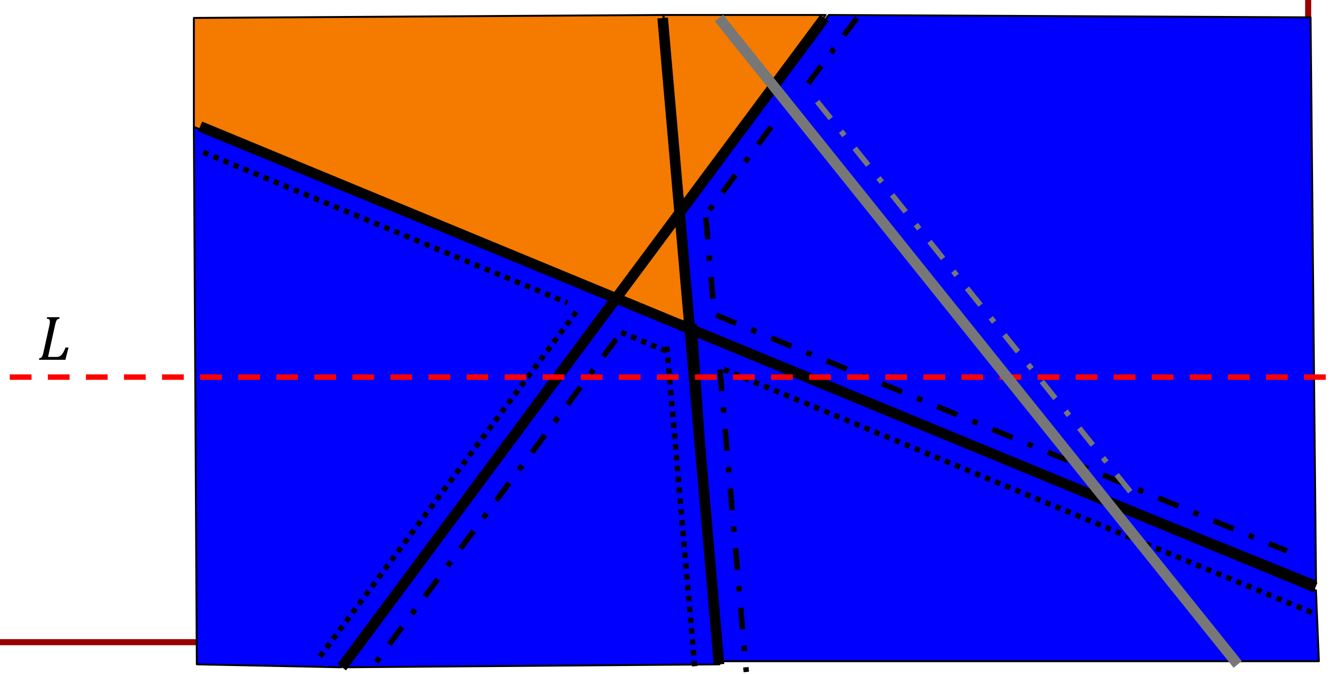


If it is split by a line from the right, then it wasn't right-most.

Zone Theorem

Claim:

Adding the right-most line splits at most two existing left edges.

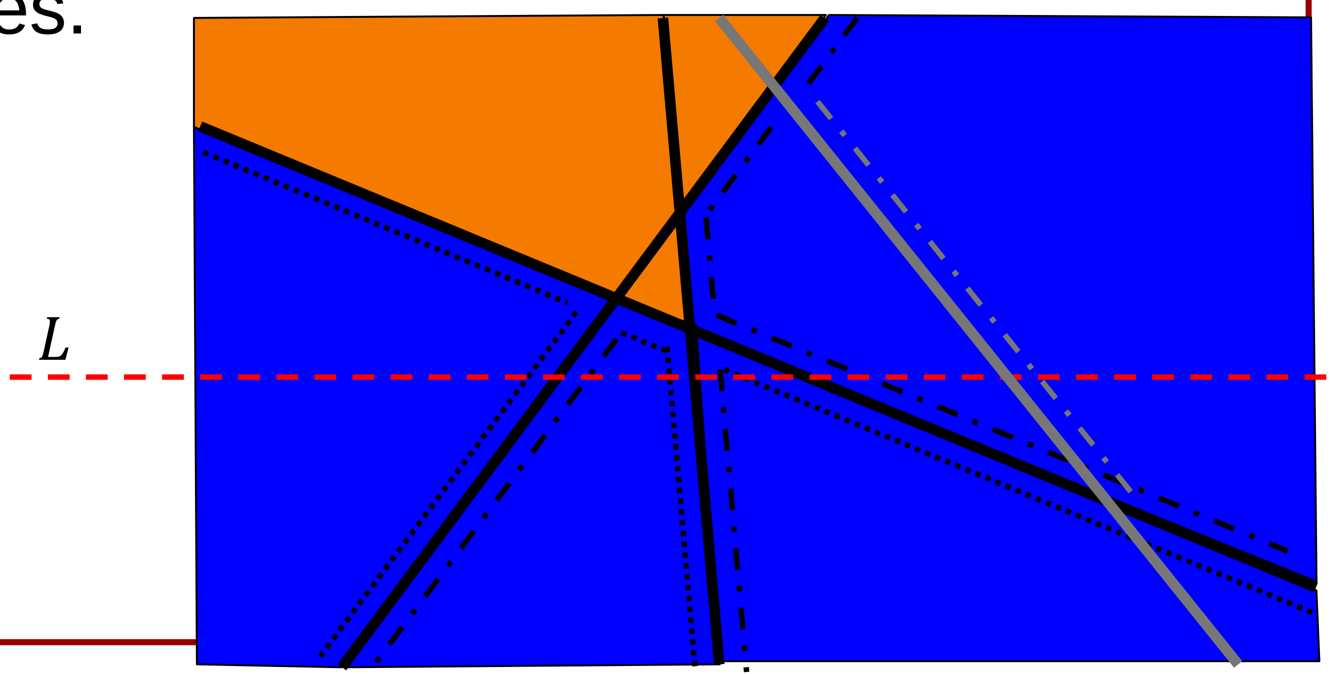


Zone Theorem

Proof of Claim:

As above, if the right-most line splits a left edge in two, the edge has to be on the right-most face.

Since faces are convex, the line can split at most two such edges.



Zone Theorem

Corollary:

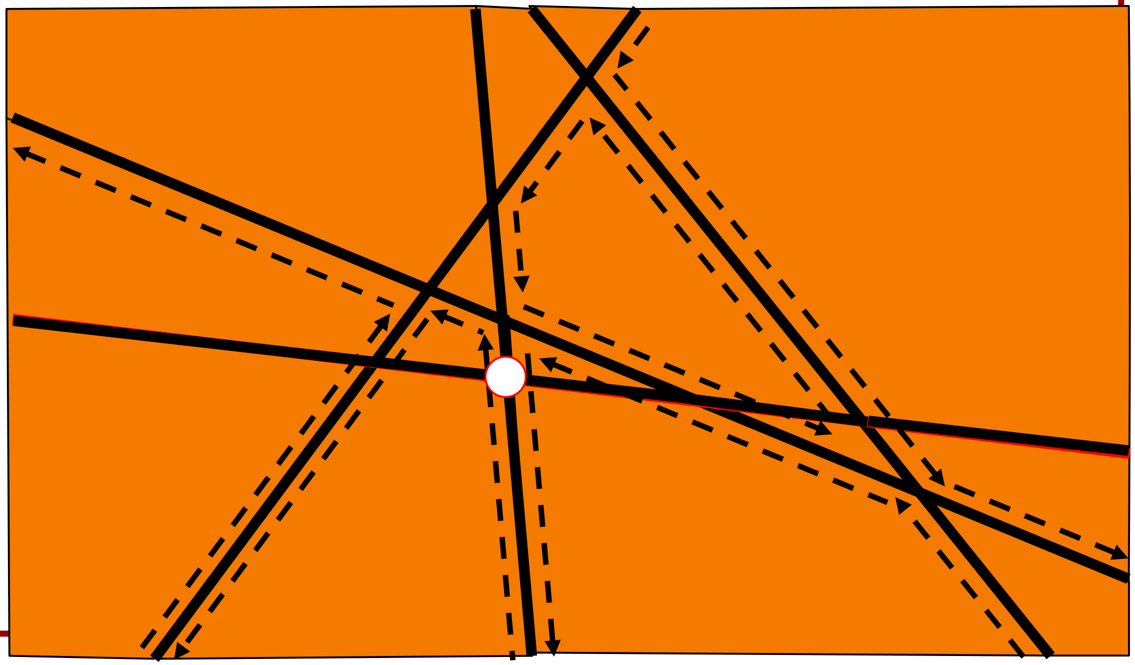
We can construct a (simple) arrangement of n lines in $O(n^2)$ time.

Zone Theorem

Proof:

Iteratively add lines.

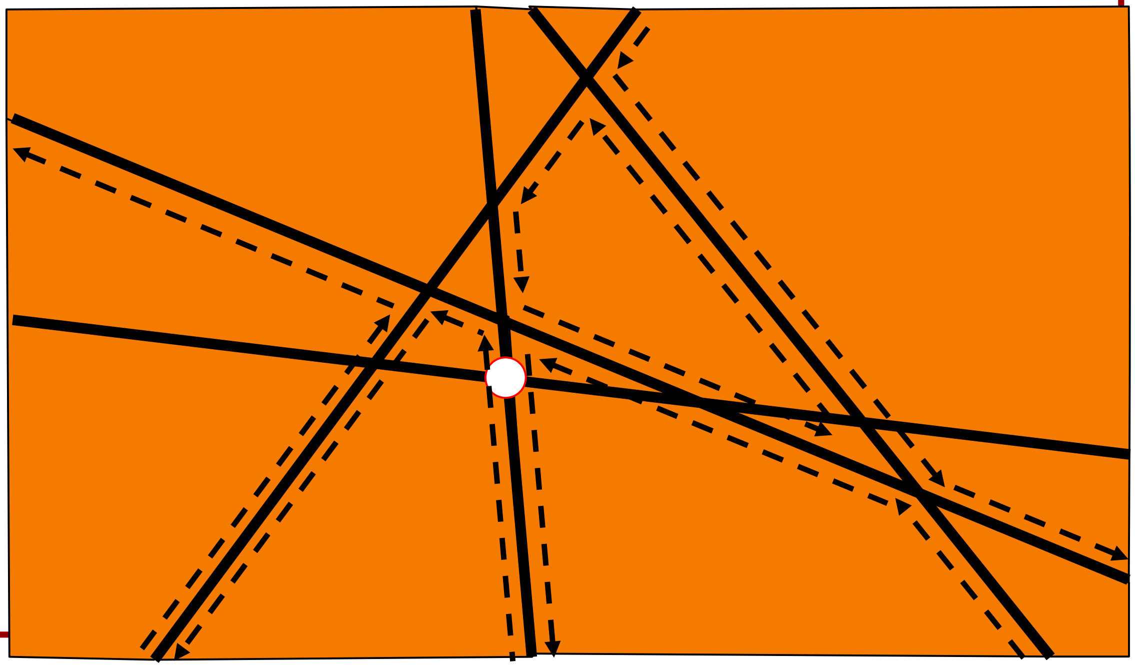
- Find an intersection with any existing line.
- Cycle around faces to the left
- Cycle around faces to the right



Zone Theorem

Proof:

Since the number of edges traversed at each iteration is $O(n)$, the total complexity is $O(n^2)$.



Zone Theorem

Generalizations:

In d-dimensional space:

- The number of faces of any dimension of an arrangement is $O(n^d)$.
- The number of faces in the zone of a hyper-plane is bounded by $O(n^{d-1})$.
- The arrangement can be computed in $O(n^d)$ time.