



# FFTs in Graphics and Vision

Fast String Matching  
and  
Math Review

*Fast Pattern Matching in Strings*  
Knuth *et al.*, 1977



# Outline

Fast Substring Matching

Math Review

- Complex Numbers
- Vector Spaces
- Linear Operators



# Fast Substring Matching

Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .



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Example:

$S = A \boxed{CDB} EF \boxed{CDB} E$

$T = CDB$



# Fast Substring Matching

## Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .

## Brute Force:

- For each position in  $S$ :
  - Test if the next  $|T|$  letters in  $S$  match those in  $T$

$S = \text{ACDBEFCDDBE}$        $T = \text{CDB}$   
~~COBDBBDBBDBB~~  
 $O(|S| * |T|)$



# Fast Substring Matching

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Can we do this more efficiently?



# Fast Substring Matching

## Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .

## Observation:

On a failed match, we don't have to compare all  $|T|$  letters in  $T$ :

$S = \text{ACDBEFCDDBE}$   
 $\text{COBBCCOBBB}$

$T = \text{CDB}$

**Comparisons: 3**



# Fast Substring Matching

## Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .

## Observation:

What if the situation is more complex?

$S = \text{AAAAAA}\boxed{\text{AAAB}}$   
 $\text{AAABABABABAB}$

$T = \text{AAAB}$   
Comparisons: 4





# Fast Substring Matching

## Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .

## Knuth *et al.* (1977):

On a failed match, we don't have to re-start the matching.

Diagram illustrating the naive string matching algorithm:

- String  $S = \text{AAAAAAAAAAB}$  (length 11)
- String  $T = \text{AAAB}$  (length 4)
- The algorithm compares  $T$  with every substring of  $S$  of length  $|T|$ .
- The complexity is  $O(|S| \cdot |T|)$ .



# Fast Substring Matching

## Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .

## Knuth *et al.* (1977):

On a failed match, we don't have to re-start the matching.

The key is to know where in  $T$  we have to start comparing.



# Fast Substring Matching

## Challenge:

Given strings  $S$  and  $T$ , find all occurrences of  $T$  as a substring of  $S$ .

## Knuth *et al.* (1977):

The size of the shift on a mismatch is determined by the repetitions in  $T$ , is independent of  $S$ , and can be computed in  $O(|T|)$  time.

For more details, see:

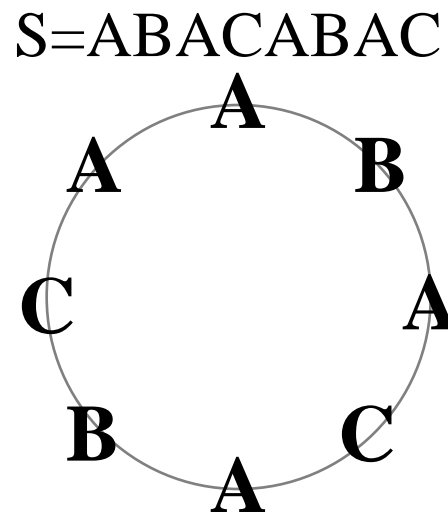
*Fast Pattern Matching in Strings.*



# Fast Substring Matching

## Applications:

If we think of a string as a signal on a circle:





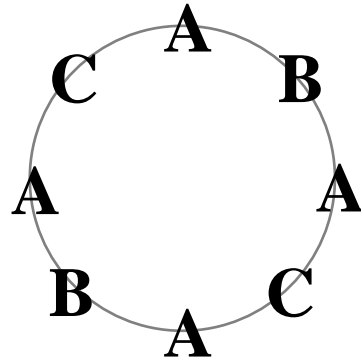
# Fast Substring Matching

## Applications:

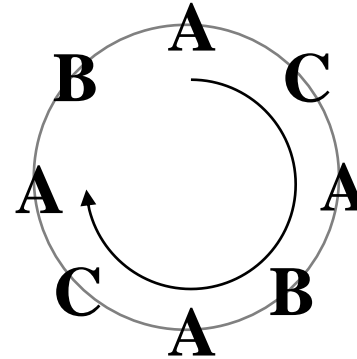
If we think of a string as a signal on a circle:

- We can test if signal  $T$  is a rotation of  $S$  by testing if  $T$  is a substring of  $SS$

$S=ABACABAC$



$T=ACABACAB$



$SS=ABACABACABACABAC$   
 $T=$   $ACABACAB$



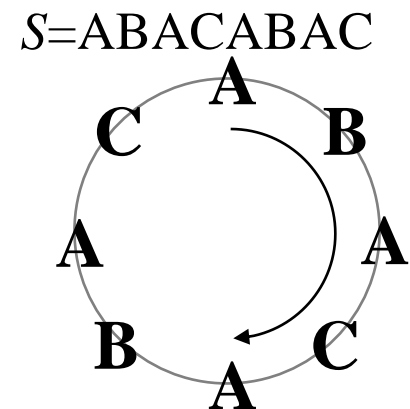
# Fast Substring Matching

## Applications:

If we think of a string as a signal on a circle:

- We can test if signal  $T$  is a rotation of  $S$  by testing if  $T$  is a substring of  $SS$
- We can test if  $S$  has rotational symmetry by testing if  $S$  is a substring of  $SS$

$SS = \text{ABACABACABACABAC}$   
 $T = \text{ABACABAC}$





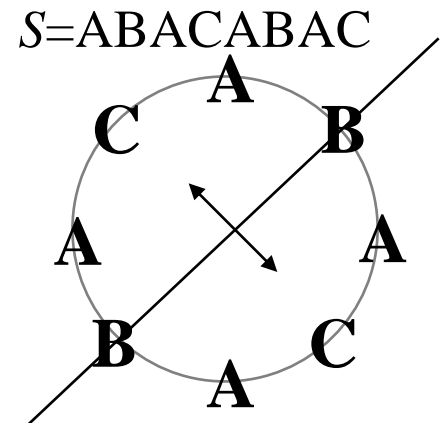
# Fast Substring Matching

## Applications:

If we think of a string as a signal on a circle:

- We can test if signal  $T$  is a rotation of  $S$  by testing if  $T$  is a substring of  $SS$ :
- We can test if  $S$  has rotational symmetry by testing if  $S$  is a substring of  $SS$ .
- We can test if  $S$  has reflective symmetry by testing if  $S$  is a substring of  $(SS)^t$

$(SS)^t = \text{CABACABACABACABA}$   
 $T = \text{ABACABAC}$





# Fast Substring Matching

## Advantages:

- A fast (linear time) algorithm for performing pattern detection on discrete signals.

## Disadvantages:

- Can only tell us if there is a perfect match
  - We need a continuous measure of similarity for real-world data
- Only works for signals on a circle (or a line)
  - Hard to generalize to signals on more complex / interesting domains





# Outline

Fast Substring Matching

Math Review

- Complex Numbers
- Vector Spaces
- Linear Operators



# Complex Numbers

- A complex number  $c$  is any number that can be written as:

$$c = a + ib$$

where  $a$  and  $b$  are real numbers and  $i$  is the square root of -1:

$$i^2 = -1$$



# Complex Numbers

Given two complex numbers,  $c_1 = a_1 + ib_1$  and  $c_2 = a_2 + ib_2$ :

- The sum of the numbers is:

$$c_1 + c_2 = (a_1 + a_2) + i(b_1 + b_2)$$



# Complex Numbers

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- The sum of the numbers is:

$$c_1 + c_2 = (a_1 + a_2) + i(b_1 + b_2)$$

- The product of the numbers is:

$$\begin{aligned} c_1 c_2 &= (a_1 + ib_1)(a_2 + ib_2) \\ &= a_1 a_2 + ib_1 ib_2 + a_1 ib_2 + ib_1 a_2 \\ &= (a_1 a_2 - b_1 b_2) + i(a_1 b_2 + b_1 a_2) \end{aligned}$$



# Complex Numbers

Given a complex numbers,  $c=a+ib$ :

- The negation of the number is:

$$-c = -a - ib$$



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# Complex Numbers

Given a complex numbers,  $c=a+ib$ :

- The negation of the number is:

$$-c = -a - ib$$

- The conjugate of the number is:

$$\bar{c} = a - ib$$

- The reciprocal of the number is:

$$\frac{1}{c} = \frac{1}{c} \frac{\bar{c}}{\bar{c}} = \left( \frac{a}{a^2 + b^2} \right) - i \left( \frac{b}{a^2 + b^2} \right)$$

# Complex Numbers

Why do we care?







# Complex Numbers

Why do we care?

## Fundamental Theorem of Algebra

Given any polynomial:

$$p(x) = a_0 + a_1x + a_2x^2 + \cdots + a_nx^n$$

there always exists a complex number  $c_0$  s.t.:

$$p(c_0) = 0$$



# Vector Spaces

A (real/complex) vector space  $V$  is a set of elements  $v \in V$ , with:

- An addition operator “+”, and
- A scaling operator “.”

(i.e. we can add any two vectors together to get a vector and if we scale a vector by a number we also get a vector.)



# Vector Spaces (Formal Properties 1)

For all  $u$ ,  $v$ , and  $w$  in  $V$ :

- Associative Addition:

$$(u+v)+w=u+(v+w)$$

- Commutative Addition:

$$u+v=v+u$$

- Additive Identity:

There exists a unique vector  $0$  in  $V$  such that:

$$u+0=u$$

- Additive Inverse:

There exists a vector  $(-u)$  in  $V$  such that:

$$u+(-u)=0$$



# Vector Spaces (Formal Properties 2)

For all  $u$ , and  $v$  in  $V$ , and all (real / complex) scalars  $a$  and  $b$ :

- Distributive over vector addition:

$$a(u+v) = (au) + (av)$$

- Distributive over scalar addition:

$$(a+b)u = (au) + (bu)$$

- Compatible scalar multiplication:

$$a(bu) = (ab)u$$

- Scalar Identity:

$$1u = u$$



# Vector Spaces: Examples

## Real Vector Spaces:

- The real / complex numbers
- The space of  $n$ -dimensional arrays with real / complex entries
- The space of  $m \times n$  matrices with real / complex entries
- The space of real / complex valued functions on a circle / line / plane / sphere / etc.

## Complex Vector Spaces:

- The complex numbers
- The space of  $n$ -dimensional arrays with complex entries
- The space of  $m \times n$  matrices with complex entries
- The space of complex valued functions on a circle / line / plane / sphere / etc.



# Vector Space Basis

A basis of  $V$  is a finite set  $\{v_1, \dots, v_n\}$  of vectors such that:

1. Any vector  $v$  in  $V$  can be expressed as:

$$v = a_1 v_1 + \dots + a_n v_n$$

where the  $a_i$  are (real / complex) scalars.

2. No basis vector  $v_i$  can be expressed as the linear sum of the other basis vectors.

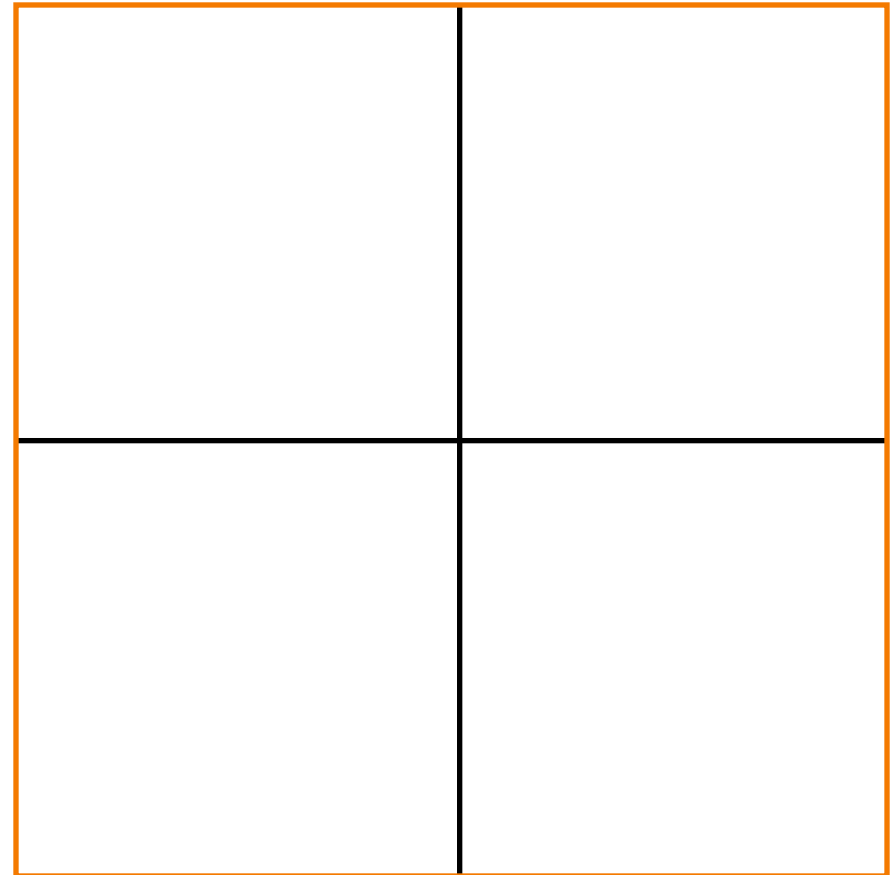


# Vector Space Basis

Many different bases can be used to represent the same vector space.

Example:

Consider the set of points in 2D Euclidean space.





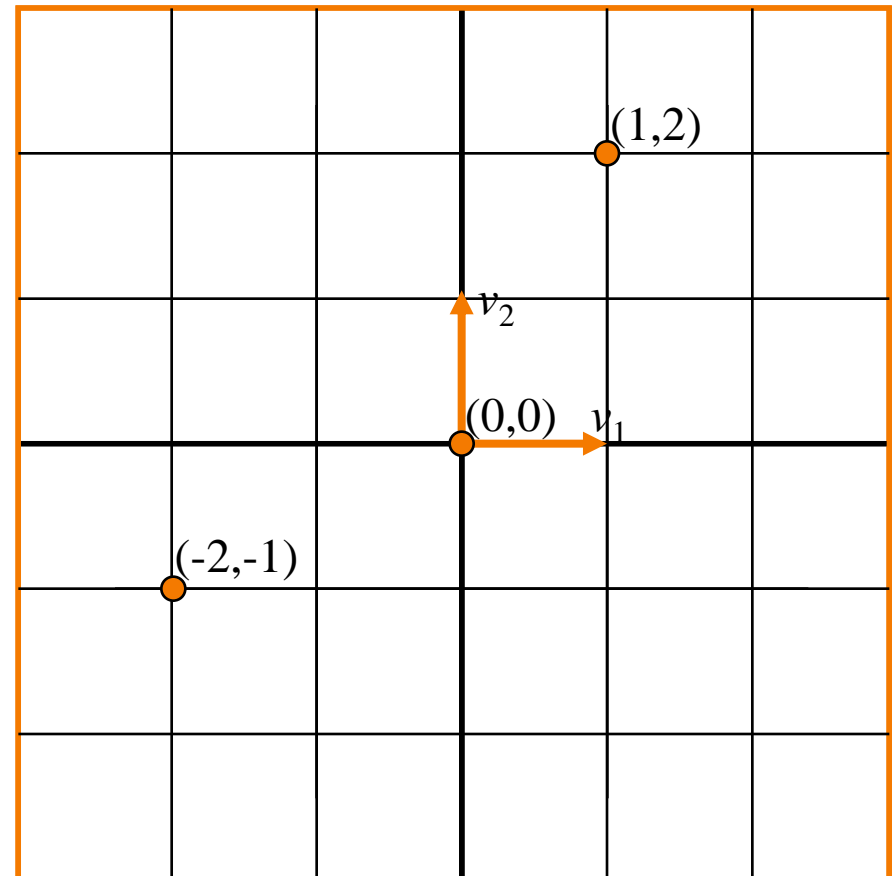
# Vector Space Basis

Many different bases can be used to represent the same vector space.

## Example:

Consider the set of points in 2D Euclidean space.

We can represent each vector in terms of its  $(x,y)$ -coordinates.







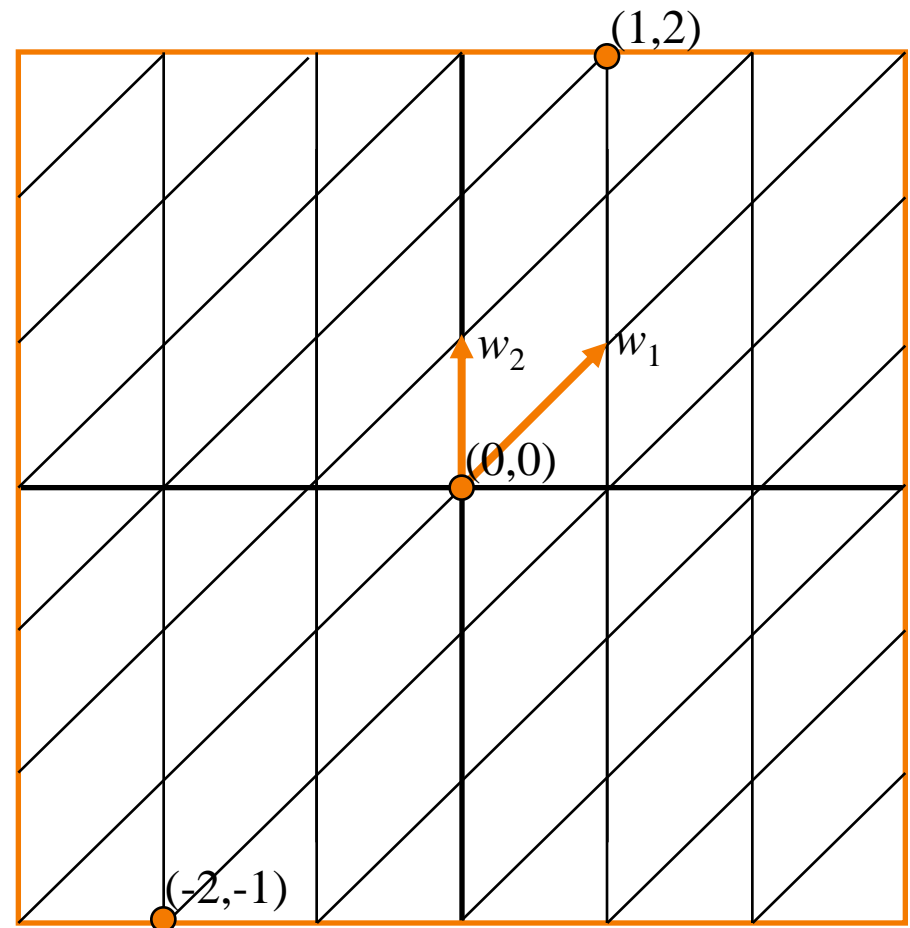
# Vector Space Basis

Many different bases can be used to represent the same vector space.

Example:

Consider the set of points in 2D Euclidean space.

Or we could use a different basis...





# Linear Maps

A function  $L: V \rightarrow W$ , is a linear map if for all  $v_1$  and  $v_2$  in  $V$  and all (real / complex) scalars  $a$  and  $b$ :

$$L(av_1 + bv_2) = a \cdot L(v_1) + b \cdot L(v_2)$$



# Linear Maps

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$$L(av_1 + bv_2) = a \cdot L(v_1) + b \cdot L(v_2)$$

If it exists, the inverse of a linear map  $L$  is the map  $L^{-1}$  with the property that:

$$L^{-1}(L(v)) = v$$



# Linear Maps

If  $L: V \rightarrow W$ , is a linear map:

The set of vectors:

$$K = \{v \in V \mid L(v) = 0\}$$

is a vector subspace called the kernel.

The set of vectors:

$$I = \{w \in W \mid \exists v \in V \text{ s.t. } L(v) = w\}$$

is a vector subspace called the image.



# Matrices

Given a vector space  $V$ , with basis  $\{v_1, \dots, v_n\}$ , a linear map can be expressed as an  $n \times n$  matrix  $M$  such that for any vector  $v = a_1 v_1 + \dots + a_n v_n$  in  $V$ :

$$L(v) = b_1 v_1 + \dots + b_n v_n$$

with:

$$\begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix} = \begin{pmatrix} M_{11} & \cdots & M_{1n} \\ \vdots & \ddots & \vdots \\ M_{n1} & \cdots & M_{nn} \end{pmatrix} \begin{bmatrix} a_1 \\ \vdots \\ a_n \end{bmatrix}$$



# Change of Basis

Given a vector space  $V$ , and given two bases  $\{v_1, \dots, v_n\}$  and  $\{w_1, \dots, w_n\}$ , then since  $\{v_1, \dots, v_n\}$  is a basis, there exist values  $B_{ij}$  such that:

$$\begin{aligned} w_1 &= B_{11}v_1 + \cdots + B_{1n}v_n \\ &\vdots \\ w_n &= B_{n1}v_1 + \cdots + B_{nn}v_n \end{aligned}$$



# Change of Basis

Given a vector space  $V$ , and given two bases  $\{v_1, \dots, v_n\}$  and  $\{w_1, \dots, w_n\}$ , the matrix  $B$  is the change of basis matrix.

If  $v$  is any vector in  $V$ , we can write out  $v$  in terms of the basis  $\{v_1, \dots, v_n\}$  as  $v = a_1 v_1 + \dots + a_n v_n$ .

We can also write out  $v$  in terms of the basis  $\{w_1, \dots, w_n\}$  as  $v = b_1 w_1 + \dots + b_n w_n$ .

The coefficients are related by:

$$\begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix} = \begin{pmatrix} B_{11} & \cdots & B_{1n} \\ \vdots & \ddots & \vdots \\ B_{n1} & \cdots & B_{nn} \end{pmatrix} \begin{bmatrix} a_1 \\ \vdots \\ a_n \end{bmatrix}$$



# Change of Basis

Given:

- A vector space  $V$ ,
- Two bases  $\{v_1, \dots, v_n\}$  and  $\{w_1, \dots, w_n\}$ ,
- A linear operator  $L$  represented by the matrix  $M$  in terms of the basis  $\{v_1, \dots, v_n\}$ .

The matrix representation for  $L$  in terms of the basis  $\{w_1, \dots, w_n\}$  is given by:

$$BMB^{-1}$$



# Change of Basis

Why do we care?





# Change of Basis

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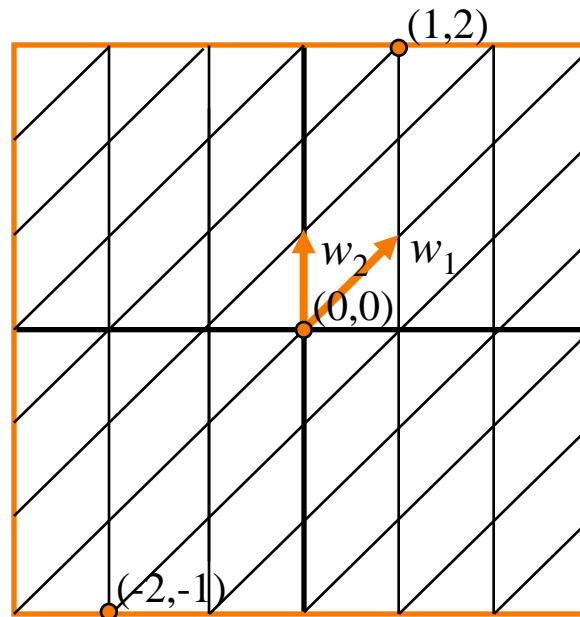
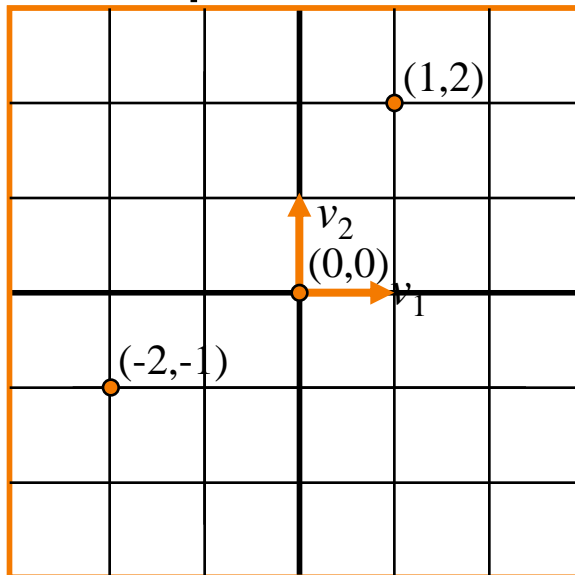
Choosing the appropriate basis can make it much easier to understand a linear operator.



# Change of Basis

Why do we care?

Example:



$$B = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$$

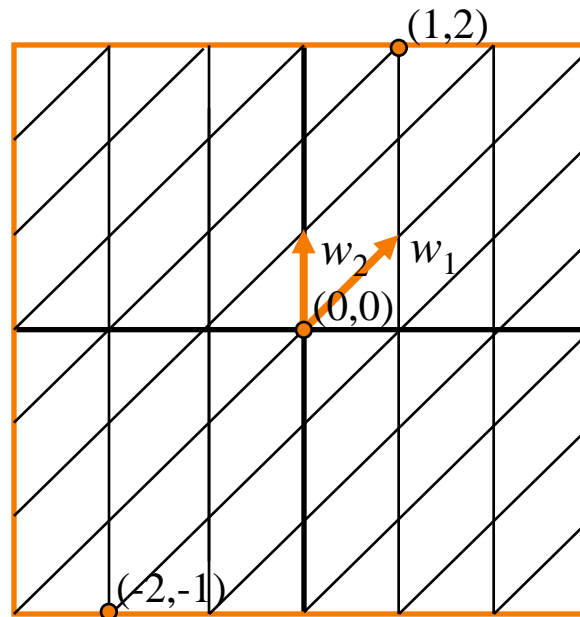
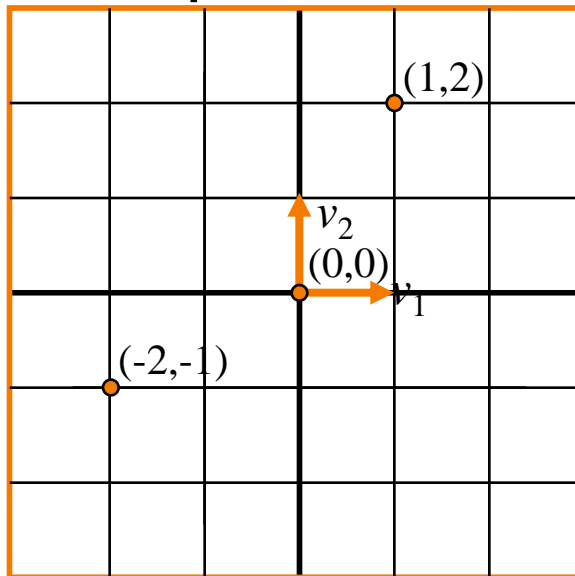
$$B^{-1} = \begin{pmatrix} 1 & 0 \\ -1 & 1 \end{pmatrix}$$



# Change of Basis

Why do we care?

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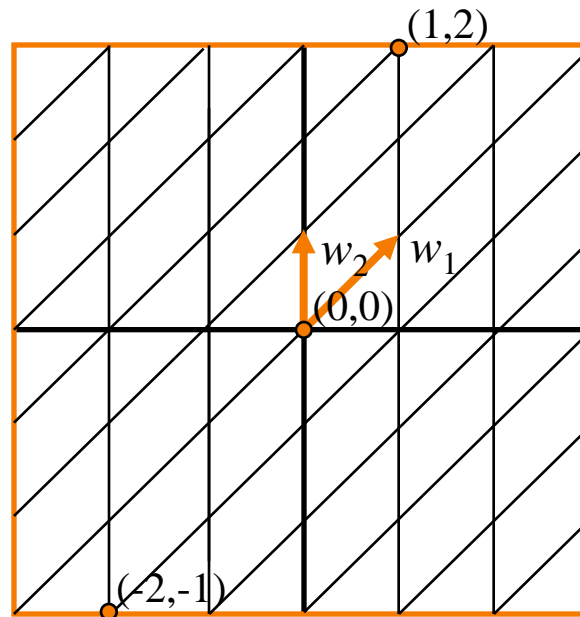
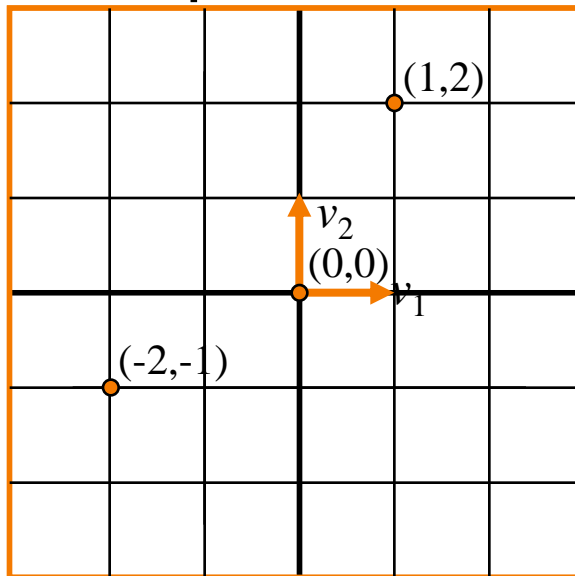
$$M = \begin{pmatrix} 2 & 0 \\ -1 & 1 \end{pmatrix} \longrightarrow BMB^{-1}$$



# Change of Basis

Why do we care?

Example:



$$B = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$$

$$B^{-1} = \begin{pmatrix} 1 & 0 \\ -1 & 1 \end{pmatrix}$$

$$M = \begin{pmatrix} 2 & 0 \\ -1 & 1 \end{pmatrix} \longrightarrow BMB^{-1} = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 2 & 0 \\ -1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ -1 & 1 \end{pmatrix} = \begin{pmatrix} 2 & 0 \\ 0 & 1 \end{pmatrix}$$



# Determinants

The determinant is a function that associates a scalar value to every square ( $n \times n$ ) matrix.

One way to think about this is to write out the matrix as a set of column vectors:

$$M = [w_1 | \cdots | w_n]$$

Then the determinant of  $M$  is the (signed) volume of the parallelepiped with sides  $\{w_1, \dots, w_n\}$



# Determinants

The determinant of a matrix  $M$  is equal to zero if and only if there exists a vector  $v$  in  $V$ , with  $v \neq 0$ , such that  $M(v) = 0$ .



# Eigenvalues and Eigenvectors

The scalar  $\lambda$  is an eigenvalue of a matrix  $M$  if there exists a vector  $v$  in  $V$  such that:

$$\lambda v = M(v)$$

In this case,  $v$  is an eigenvector of  $M$ .





# Eigenvalues and Eigenvectors

If  $M$  has an eigenvector  $v$  with eigenvalue  $\lambda$ , this must mean that:

$$0 = (M - \lambda)(v)$$

Thus, the matrix:

$$M - \lambda \cdot \text{Id} = \begin{pmatrix} M_{11} - \lambda & \cdots & M_{1n} \\ \vdots & \ddots & \vdots \\ M_{n1} & \cdots & M_{nn} - \lambda \end{pmatrix}$$

must have zero determinant.



# Characteristic Polynomials

If we treat  $\lambda$  as a variable, then the determinant:

$$\det(M - \lambda \cdot \text{Id})$$

is a polynomial of degree  $n$  in  $\lambda$ . This polynomial is the characteristic polynomial of  $M$ .



# Characteristic Polynomials

The roots of the characteristic polynomial of  $M$ :

$$\det(M - \lambda \cdot \text{Id})$$

are precisely the eigenvalues of the matrix  $M$ .

Thus, if we are considering  $M$  as a matrix acting on a complex vector space,  $M$  must always have at least one eigenvalue.