

1/30/25;

Exponential Mechanism:

Previous setup: utility was directly connected to noise added.

- not always the case! Sometimes utility very complicated fn of noise, or not even numeric in first place!

Ex from book: auctions.

want to set price, but keep bids private.
(max revenue)
Four bidders A, B, C, D .

$A, B, C: \$1$

$D: \$3.01$

utility (revenue): $x \leq 3.01: x$
 $> 3.01: 0$

Exponential mechanism: queries w/ arbitrary utilities,
possibly non-numeric range, DP.

Start from same setup: databases \mathcal{D} , want to output something from range R .

What is "something"?

$u: \mathcal{D} \times \mathcal{R} \rightarrow \mathbb{R}$ utility function

- If our database is D , and output $r \in \mathcal{R}$,
 $u(D, r)$ is how happy we are.

So want to output $\arg \max_{r \in \mathcal{R}} u(D, r)$

Sensitivity: only measured for database parameter

$$\Delta u = \max_{r \in \mathcal{R}} \max_{D \sim D'} |u(D, r) - u(D', r)|$$

Def (exponential mechanism): $M_E(D, u, \mathcal{R})$:

output $r \in \mathcal{R}$ with probability proportional to
 $\exp\left(\frac{\epsilon u(D, r)}{2\Delta u}\right)$

Thm: Exponential mechanism is ϵ -DP.

PF: Let $D \sim D' \in \mathcal{D}$, $r \in \mathcal{R}$.

$$\frac{\Pr[M_E(D, u, \mathcal{R}) = r]}{\Pr[M_E(D', u, \mathcal{R}) = r]} = \frac{\frac{\exp\left(\frac{\epsilon u(D, r)}{2\Delta u}\right)}{\sum_{r' \in \mathcal{R}} \exp\left(\frac{\epsilon u(D, r')}{2\Delta u}\right)}}{\frac{\exp\left(\frac{\epsilon u(D', r)}{2\Delta u}\right)}{\sum_{r' \in \mathcal{R}} \exp\left(\frac{\epsilon u(D', r')}{2\Delta u}\right)}}$$

$$= \frac{\exp\left(\frac{\varepsilon \psi(D, r)}{2D_n}\right)}{\exp\left(\frac{\varepsilon \psi(D', r')}{2D_n}\right)} \cdot \frac{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D', r')}{2D_n}\right)}{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D, r')}{2D_n}\right)}$$

$$= \exp\left(\frac{\varepsilon (\psi(D, r) - \psi(D', r'))}{2D_n}\right) \cdot \frac{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D', r')}{2D_n}\right)}{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D, r')}{2D_n}\right)}$$

$$\leq \exp\left(\frac{\varepsilon}{2}\right) \cdot \frac{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D', r')}{2D_n}\right)}{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D, r')}{2D_n}\right)}$$

$$\leq \exp\left(\frac{\varepsilon}{2}\right) \cdot \frac{\sum_{r' \in R} \exp\left(\frac{\varepsilon (\psi(D, r') + D_n)}{2D_n}\right)}{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D, r')}{2D_n}\right)}$$

$$= \exp\left(\frac{\varepsilon}{2}\right) \cdot \frac{\sum_{r' \in R} \exp\left(\frac{\varepsilon}{2}\right) \exp\left(\frac{\varepsilon \psi(D, r')}{2D_n}\right)}{\sum_{r' \in R} \exp\left(\frac{\varepsilon \psi(D, r')}{2D_n}\right)}$$

$$= \exp\left(\frac{\varepsilon}{2}\right) \cdot \exp\left(\frac{\varepsilon}{2}\right) = \exp(\varepsilon) \quad \checkmark$$

Quality: Is exponential mechanism "good"?

Intuition: probability drops exponentially w/ utility,
so might be pretty good quality!

Def: Let $OPT_u(D) = \max_{r \in R} u(D, r)$ be max

achievable utility.

Let $R_{opt}(D) = \{r \in R : u(D, r) = OPT_u(D)\}$ be outputs
with max utility

Thm: $Pr[u(M_F(D, u, R)) \leq OPT_u(D) - \frac{2\Delta u}{\epsilon} (\ln(\frac{|R|}{|R_{opt}(D)|}) + t)]$
 $\leq e^{-t}$

Pf: Fix some c . Let $S_c = \{r \in R : u(D, r) \leq c\}$

$$Pr[u(M_F(D, u, R)) \leq c] = Pr[M_F(D, u, r) \in S_c]$$

$$= \frac{\sum_{r \in S_c} \exp\left(\frac{\epsilon u(D, r)}{2\Delta u}\right)}{\sum_{r \in R} \exp\left(\frac{\epsilon u(D, r)}{2\Delta u}\right)} \leq \frac{\sum_{r \in S_c} \exp\left(\frac{\epsilon c}{2\Delta u}\right)}{\sum_{r \in R_{opt}(D)} \exp\left(\frac{\epsilon \cdot OPT_u(D)}{2\Delta u}\right)}$$

$$\leq \frac{|S_c| \exp\left(\frac{\xi c}{2D_n}\right)}{|R_{\text{opt}}(D)| \exp\left(\frac{\xi \text{OPT}_n(D)}{2D_n}\right)}$$

$$\leq \frac{|R|}{|R_{\text{opt}}(D)|} \cdot \exp\left(\frac{\xi(c - \text{OPT}_n(D))}{2D_n}\right)$$

$$S_c + c - \text{OPT}_n(D) - \frac{2D_n}{\xi} \left(\ln\left(\frac{|R|}{|R_{\text{opt}}(D)|}\right) + t \right) :$$

$$\leq \frac{|R|}{|R_{\text{opt}}(D)|} \cdot \exp\left(\frac{\xi \left(-\frac{2D_n}{\xi} \left(\ln\left(\frac{|R|}{|R_{\text{opt}}(D)|}\right) + t \right) \right)}{2D_n}\right)$$

$$= \frac{|R|}{|R_{\text{opt}}(D)|} \cdot \exp\left(-\ln\left(\frac{|R|}{|R_{\text{opt}}(D)|}\right) - t\right)$$

$$= \exp(-t) \quad \checkmark$$

(Note: other, more refined versions. See, e.g., Wikipedia (1.4).)

Corollary: $\Pr\left[\ln(M_F(D, n, R)) \leq \text{OPT}_n(D) - \frac{2D_n}{\xi} (\ln |R| + t)\right]$

$$\leq e^{-t}$$

PF: $|R_{\text{opt}}(D)| \geq 1$

✓

Ex: Best of 2.

Two medical conditions, A and B. which one
(common?)

Sp's true count of B = true count of A + C.

$u(B) = u(A) + C$ (utilities are true counts).

$$\Pr[\text{Count}_{\text{opt}} A] = \Pr[u(M_F(D, u))] \leq \text{OPT}_u(D) - C$$

$$C \leq \frac{2\Delta}{\epsilon} (\ln |R| + t) \leq \frac{2}{\epsilon} (t + 1)$$

$$\Rightarrow t \geq C \frac{\epsilon}{2} - 1$$

$$\Rightarrow \Pr[\text{Count}_{\text{opt}} A] \leq e^{-\frac{C\epsilon}{2}}$$

Ex: Laplace Mechanism!

Sp's $f: D \rightarrow \mathbb{R}$, sensitivity of

Laplace is add $Lap(\frac{\Delta f}{\epsilon})$ noise, i.e.

$$\Pr[M_L(f, D) = x] \propto \exp\left(-\frac{\epsilon |x - f(D)|}{\Delta f}\right)$$

Exponential: use $R = \mathbb{R}$, $u(D, r) = |r - f(D)|$

$$\Rightarrow \Pr[M_F(D, r) = x] \propto \exp\left(-\frac{\epsilon |x - f(D)|}{2\Delta u}\right)$$

Discussion: Exponential Mechanism super powerful!

Main downside: how do we actually implement?

Naive version: enumerate R , write down all probabilities, sample.

works, but requires enumerating R !

Often, don't want to enumerate full range: may be slow.

Ex: Synthetic database.

See given D , want to output synthetic database that looks "similar" to D : for some class of queries, answers similar.

Idea: $\sim(D, \hat{D})$ how close is \hat{D} to D in terms of query class?

\Rightarrow exponential mechanism works!

But need to enumerate all possible synthetic databases...