

## 1.1 The Topic

Citing a reference [CW87].

A description list.

**Part a:** The first part.

**Part b:** The second part.

**Part c:** The last part.

An itemized list.

- Item #1.
- Item #2.

A numbered list.

1. First
2. Second

### 1.1.1 Statements of Results

**Definition 1.1.1** *Define your problem here.*

**Theorem 1.1.2** *A theorem.*

**Lemma 1.1.3** *A helpful lemma.*

**Proof:** Proof of the lemma goes here. ■

Now we can prove Theorem 1.1.2.

**Proof of Theorem 1.1.2:** Follows from Lemma 1.1.3 ■

## 1.2 Some Formulas and Algorithms

### 1.2.1 A Linear Program

Consider linear program (**TSP-LP**) below.

$$\text{minimize: } \sum_e c(e) \cdot x_e \quad (\mathbf{TSP-LP})$$

$$\text{subject to: } x(\delta(S)) \geq 2 \quad \text{for each cut } \emptyset \subsetneq S \subsetneq V \quad (1.2.1)$$

$$x(\delta(v)) = 2 \quad \text{for each vertex } v \in V \quad (1.2.2)$$

$$x \geq 0$$

Constraints (1.2.1) are the *cut constraints* and Constraints (1.2.2) are the *degree constraints*.

### 1.2.2 Tips

Use  $\log n$ , not  $\log n$ .

$$V = \{v_1, v_2, \dots, v_n\}.$$

Check out  $\sum_{i=1}^n i$  vs.  $\sum_{i=1}^n i$ .

A displayed equation:

$$H_n = \sum_{k=1}^n \frac{1}{k} = \int_1^n \frac{dx}{x} + O(1) = \ln n + O(1)$$

A matrix:

$$\begin{pmatrix} x_1 & y_1 & z_1 \\ x_2 & y_2 & z_2 \\ 0 & 0 & 1 \end{pmatrix}$$

Problem names should look like this: SET COVER.

### 1.2.3 An Algorithm

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**Algorithm 1** Kruskal's MINIMUM SPANNING TREE Algorithm

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**Input:** Undirected graph  $G = (V, E)$  with edge costs  $c(e) \geq 0, e \in E$ .

**Output:** A minimum spanning tree of  $G$ .

$T \leftarrow \emptyset$

**for** each edge  $e \in E$  in increasing order of cost  $c(e)$  **do**

**if**  $T \cup \{e\}$  does not contain a cycle **then**

$T \leftarrow T \cup \{e\}$

**end if**

**end for**

**return**  $T$

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## References

- CW87 D. COPPERSMITH and S. WINOGRAD, Matrix multiplication via arithmetic progressions, *Proceedings of the 19th ACM Symposium on Theory of Computing*, 1987, pp. 1–6.
- S69 V. STRASSEN, Gaussian Elimination Is Not Optimal, *Numerische Mathematik* **13**, 1969, pp. 354–356.
- P84 V. PAN, *How To Multiply Matrices Faster*, Springer-Verlag, Lecture Notes in Computer Science Vol. 179, 1984.