

**Note:** For problems 2 and 3 you need to have a *formal* proof. Simply giving a state diagram or a description of a DFA/NFA is not enough.

1. (25 points) Give an NFA (both a state diagram and a formal description) recognizing the language  $0^*1^*0^+$  with three states. The alphabet is  $\{0, 1\}$ . You don't need to have a formal proof.
2. (25 points) This question studies the number of states in a DFA equivalent to an NFA. Recall that in class we showed an NFA with 4 states that recognizes the language which consists of all binary strings that have a 1 in the third position from the end. For any integer  $k$ , it is easy to generalize this construction to an NFA with  $k + 1$  states that recognizes the language which consists of all binary strings that have a 1 in the  $k$ 'th position from the end. The general transformation from an NFA to a DFA will give us a DFA with at most  $2^{k+1}$  states recognizing the same language.

Show that, any DFA that recognizes the same language must have at least  $2^k$  states.

**Hint:** start by looking at the following two strings:  $10^{k-1}$  and  $0^k$ . Observe that when a DFA takes them as inputs, it must end up at different states, since one string is accepted and the other is rejected.

3. (20 points) Say that a string  $x$  is a *prefix* of string  $y$  if a string  $z$  exists where  $xz = y$  and that  $x$  is a proper prefix of  $y$  if in addition  $x \neq y$ . Define the operation NOPREFIX on a language  $A$  to be

$$\text{NOPREFIX}(A) = \{w \in A \mid \text{no proper prefix of } w \text{ is a member of } A.\}$$

Show that the class of regular languages is closed under this operation.

**Hint:** Think about when a string  $w \in A$  can have a proper prefix which is in  $A$ , then modify the states/transitions of the machine to avoid this.

4. (25 points) Let  $\Sigma = \{0, 1\}$ .
  - (a) Write a regular expression for the language  $L$  consisting of all strings in  $\Sigma^*$  with exactly one occurrence of the substring 111.
  - (b) Write a regular expression for the language  $L$  consisting of all strings in  $\Sigma^*$  that do not end with 00.

You don't need to have a formal proof.