

Suffix Tries: size

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Suffix trie

Build a **trie** containing all **suffixes** of a text T

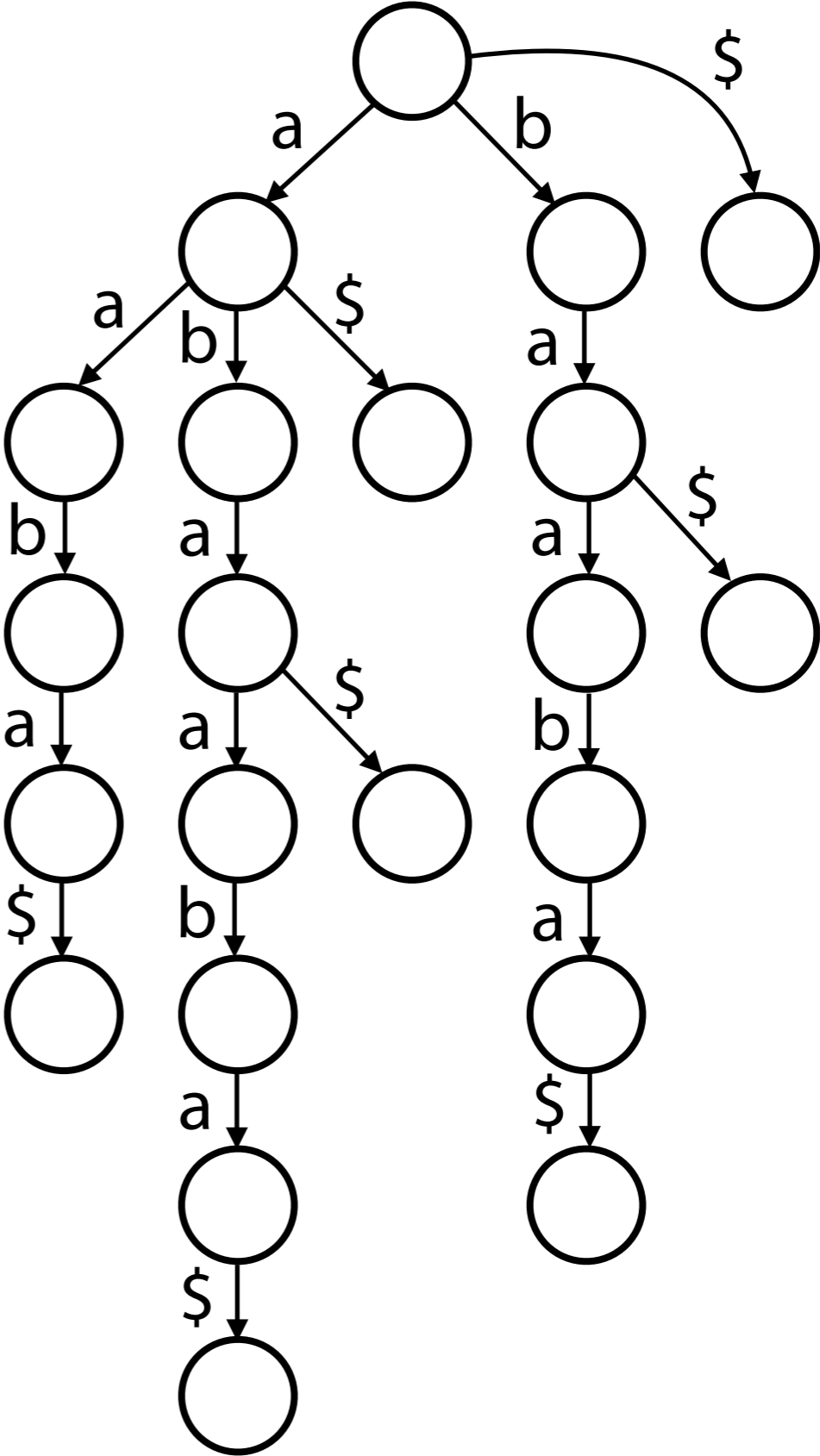
T : GTTATAGCTGATCGCGGGCGTAGCGG\$
GTTATAGCTGATCGCGGGCGTAGCGG\$
TTATAGCTGATCGCGGGCGTAGCGG\$
TATAGCTGATCGCGGGCGTAGCGG\$
ATAGCTGATCGCGGGCGTAGCGG\$
TAGCTGATCGCGGGCGTAGCGG\$
AGCTGATCGCGGGCGTAGCGG\$
GCTGATCGCGGGCGTAGCGG\$
CTGATCGCGGGCGTAGCGG\$
TGATCGCGGGCGTAGCGG\$
GATCGCGGGCGTAGCGG\$
ATCGCGGGCGTAGCGG\$
TCGCGGGCGTAGCGG\$
CGCGGGCGTAGCGG\$
GCGGGCGTAGCGG\$
CGGGCGTAGCGG\$
GGCGTAGCGG\$
GCGTAGCGG\$
CGTAGCGG\$
GTAGCGG\$
TAGCGG\$
AGCGG\$
GCGG\$
CGG\$
GG\$
G\$
\$

$m(m+1)/2$
chars

Suffix trie

How does the suffix trie grow with $|T| = m$?

T : **a b a a b a \$**
b a a b a \$
a a b a \$
a b a \$
b a \$
a \$
\$



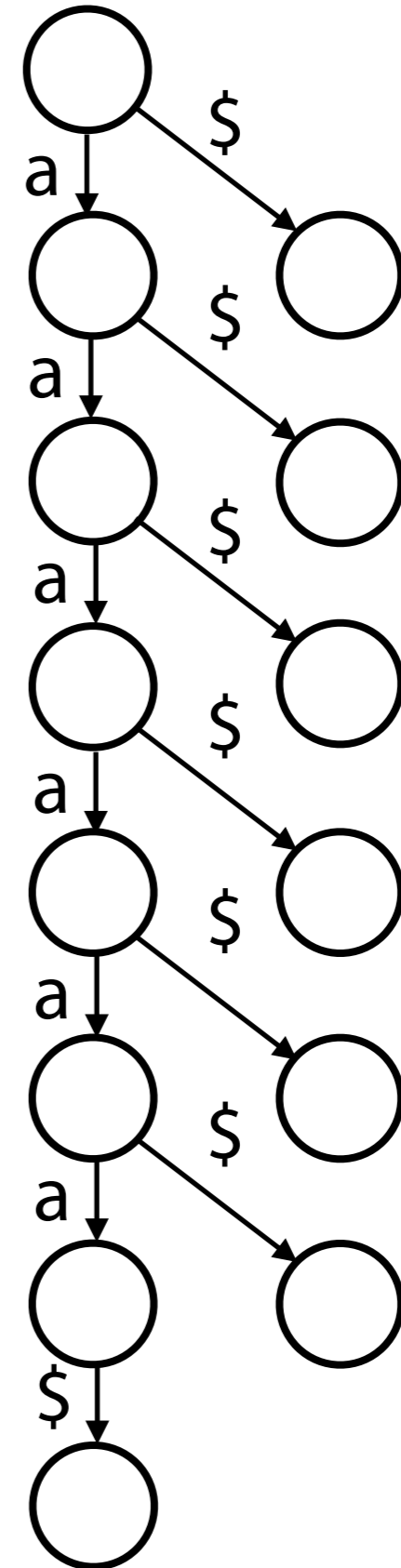
Does "prefix sharing" save us?

Suffix trie

Take repetitive strings of the form $T = a^m\$$ ($a^m\$$)

T: **a a a a a a \$**
 a a a a a \$
 a a a a \$
 a a \$
 a \$
 \$

Growth is $O(m)$, thanks to prefix sharing

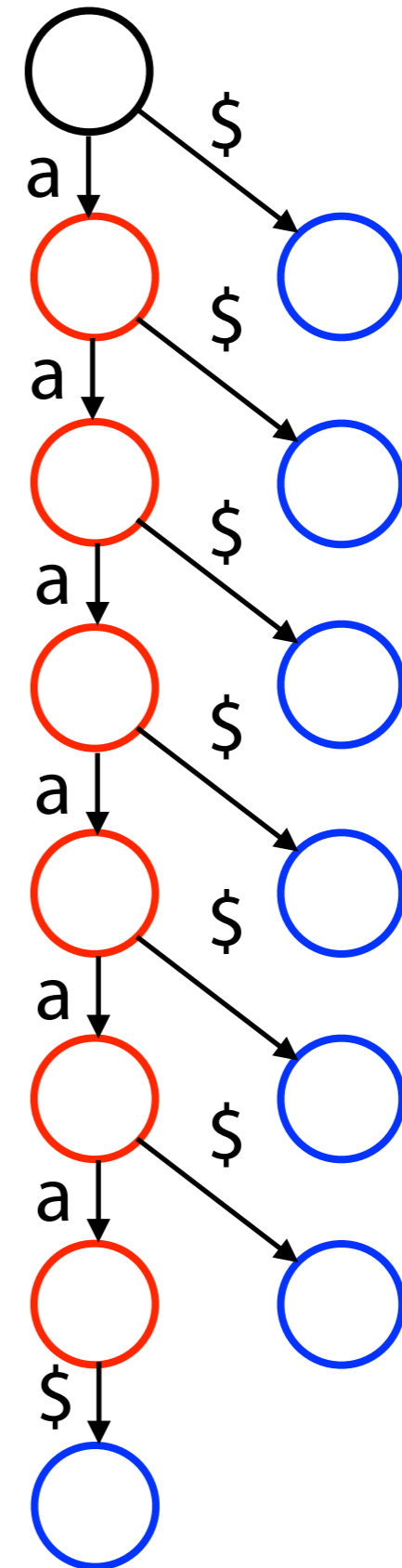


Suffix trie

Take repetitive strings of the form $T = a^m\$$ ($a^m\$$)

T : **a a a a a a \$**
 a a a a a \$
 a a a a \$
 a a \$
 a \$
 \$

Growth is $O(m)$, thanks to prefix sharing

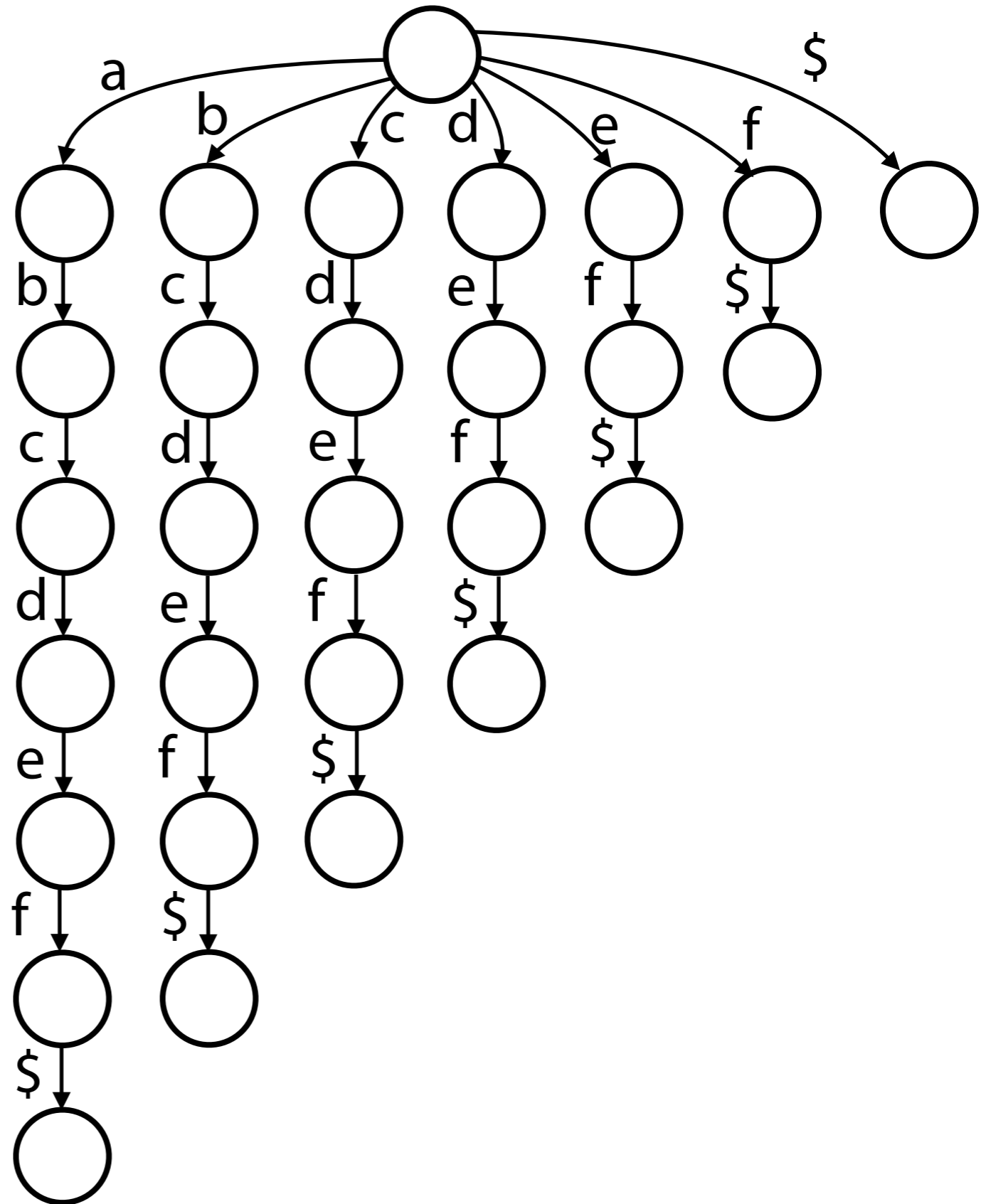


Suffix trie

Can suffixes have **no** prefix sharing?

Yes: all distinct characters

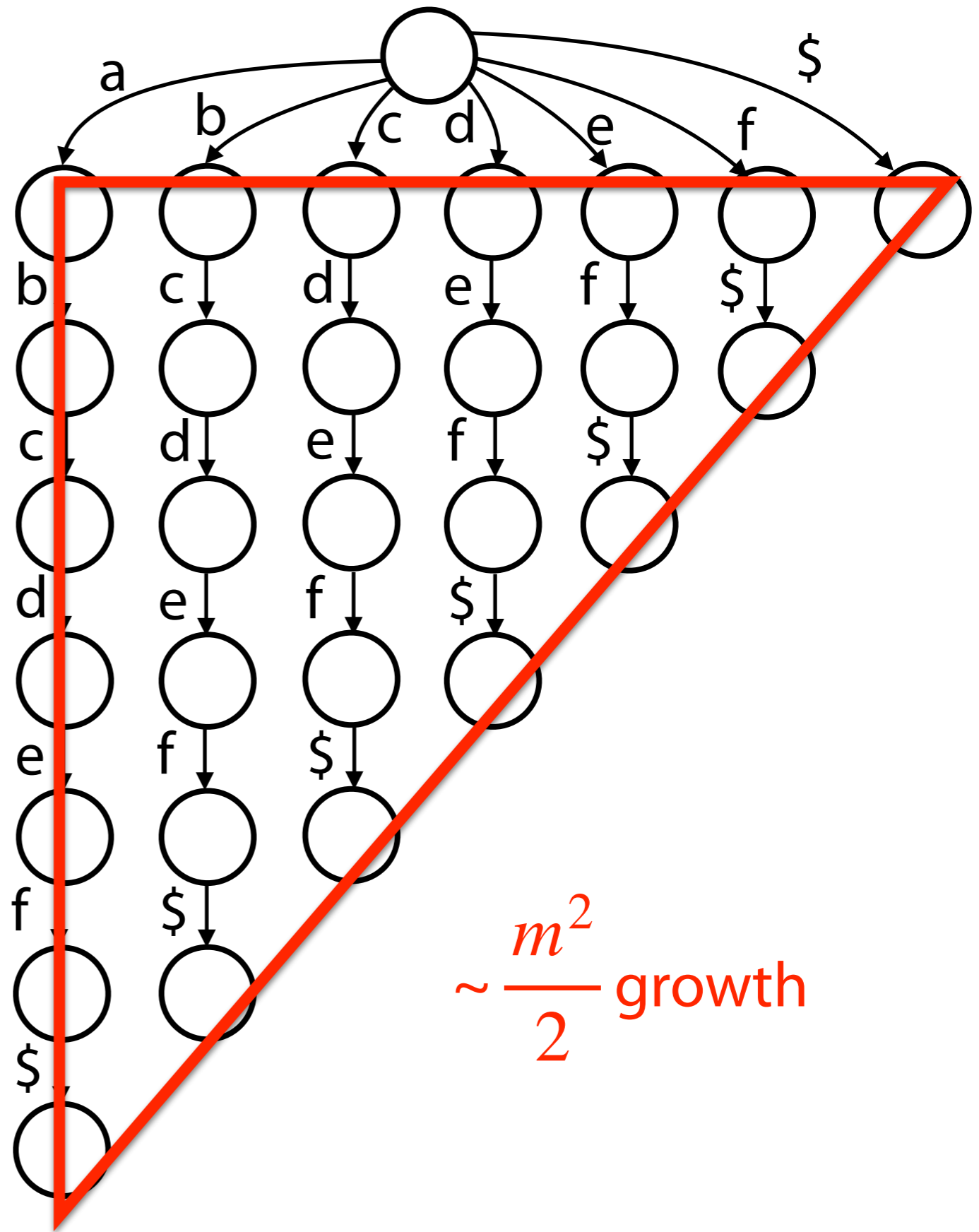
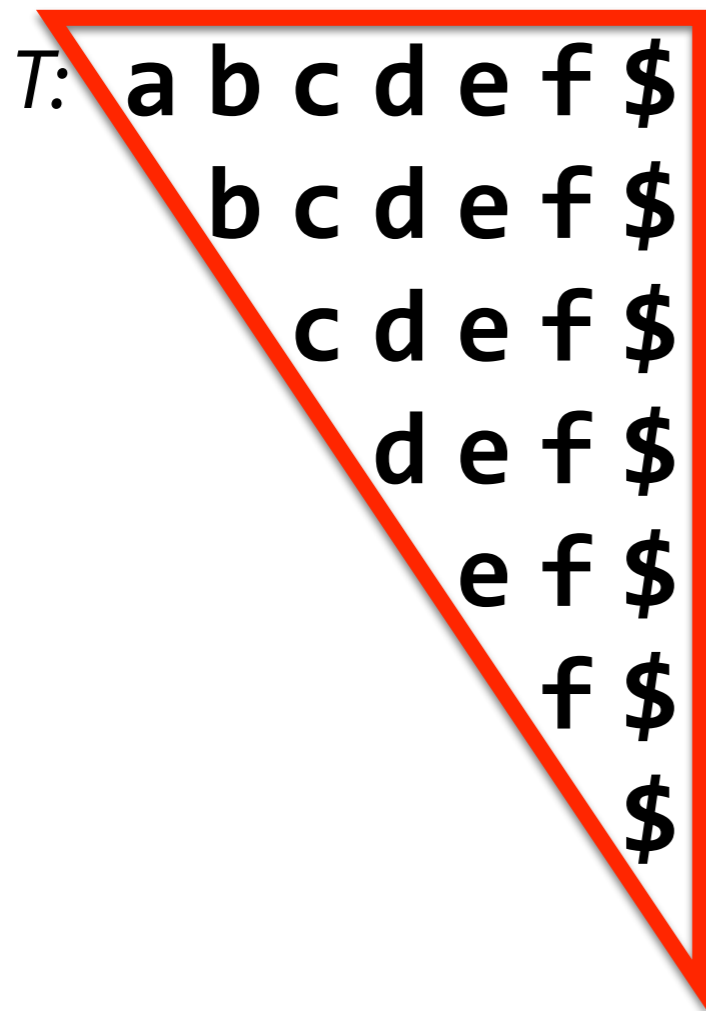
T: **a b c d e f \$**
b c d e f \$
c d e f \$
d e f \$
e f \$
f \$
\$



Suffix trie

Can suffixes have **no** prefix sharing?

Yes: all distinct characters



Suffix trie

Even when alphabet is $\{a, b\}$,
we can find strings where
suffix trie grows with $O(m^2)$

Suffix trie

Even when alphabet is {a, b},
we can find strings where
suffix trie grows with $O(m^2)$

***T*: a a a b b b \$**

Suffix trie

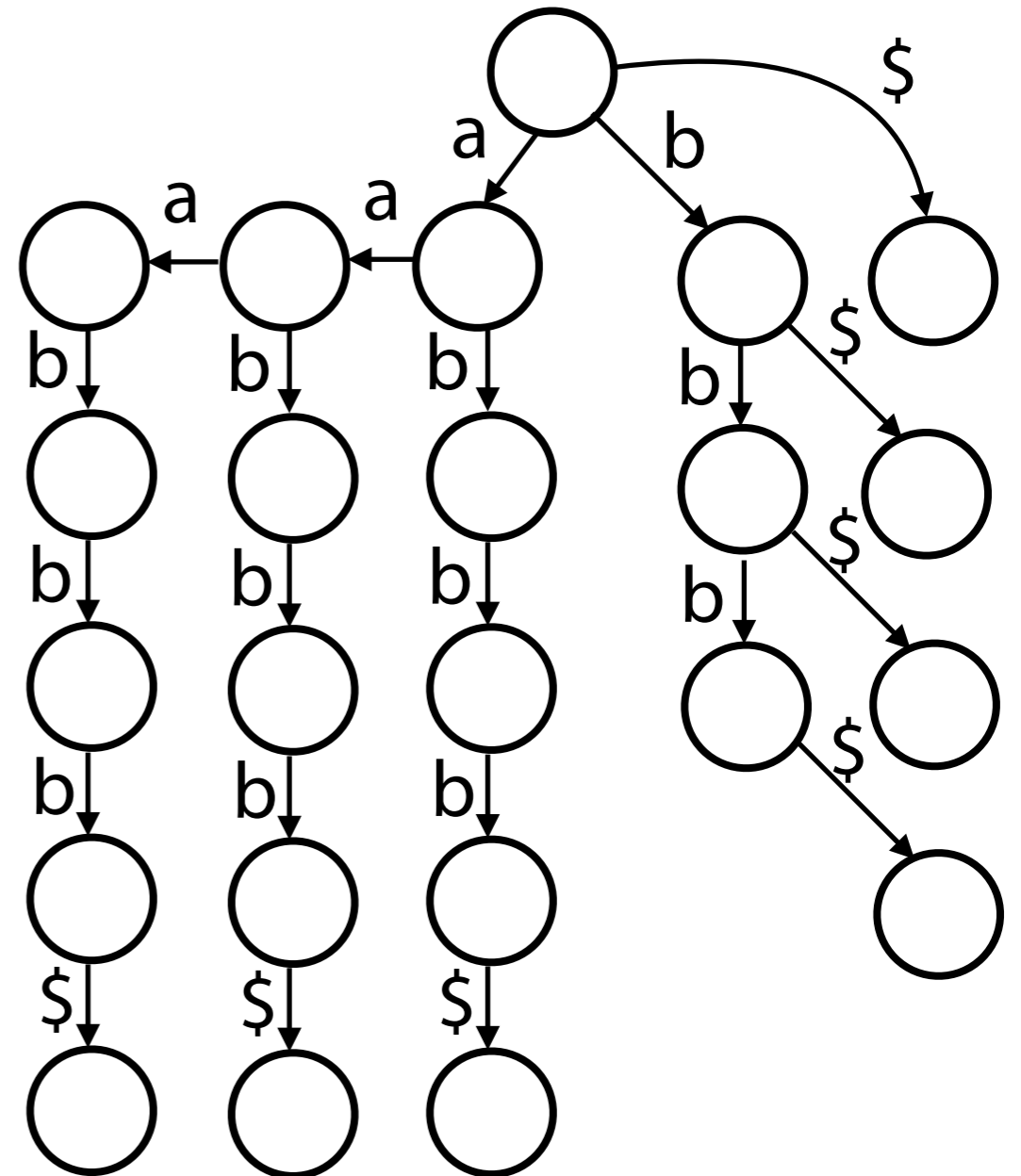
Even when alphabet is {a, b},
we can find strings where
suffix trie grows with $O(m^2)$

T: a a a b b b \$
 a a b b b \$
 a b b b \$
 b b b \$
 b b \$
 b \$
 \$

Suffix trie

Even when alphabet is {a, b},
we can find strings where
suffix trie grows with $O(m^2)$

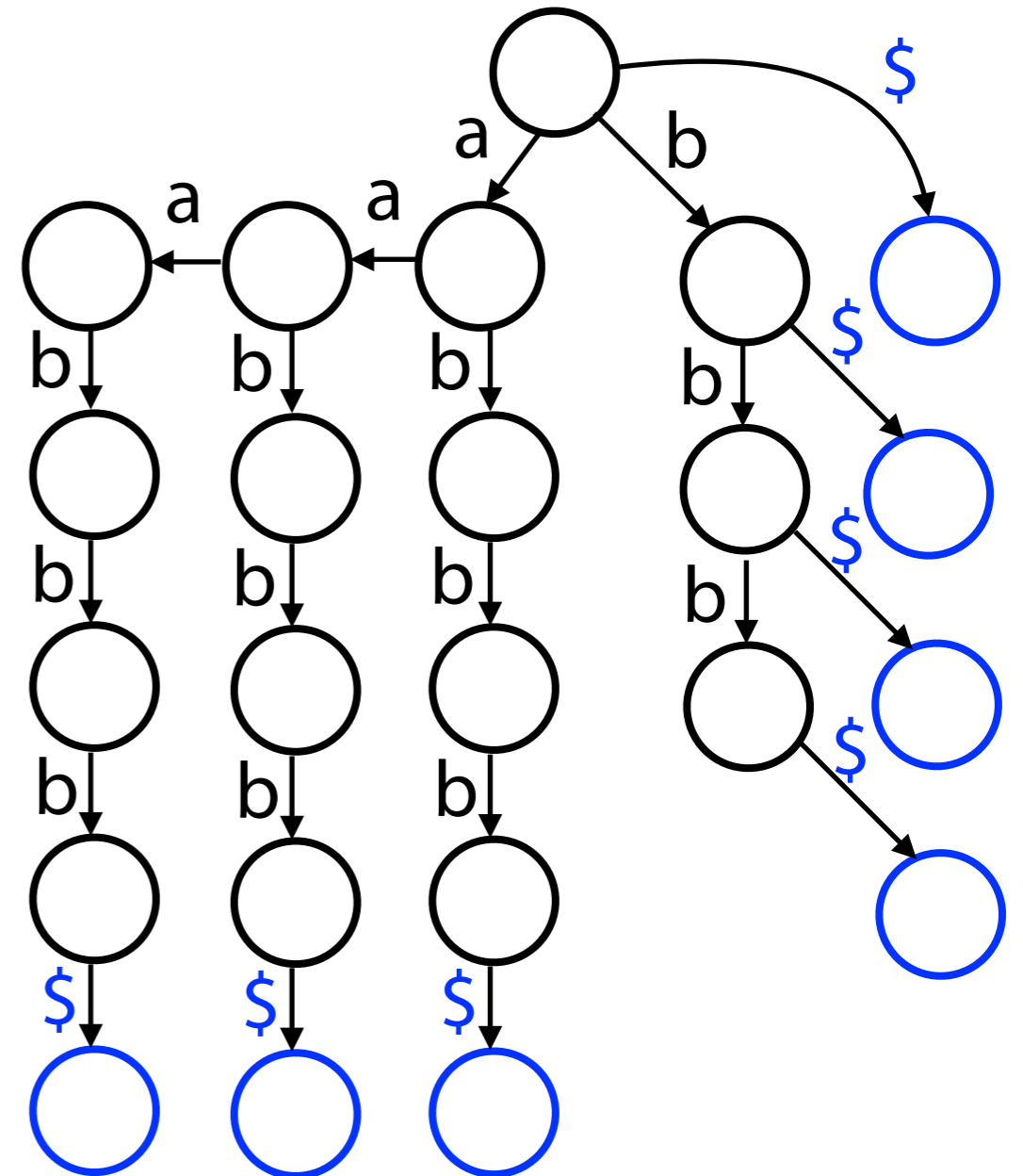
T: **a a a b b b \$**
 a a b b b \$
 a b b b \$
 b b b \$
 b b \$
 b \$
 \$



Suffix trie

Even when alphabet is {a, b},
we can find strings where
suffix trie grows with $O(m^2)$

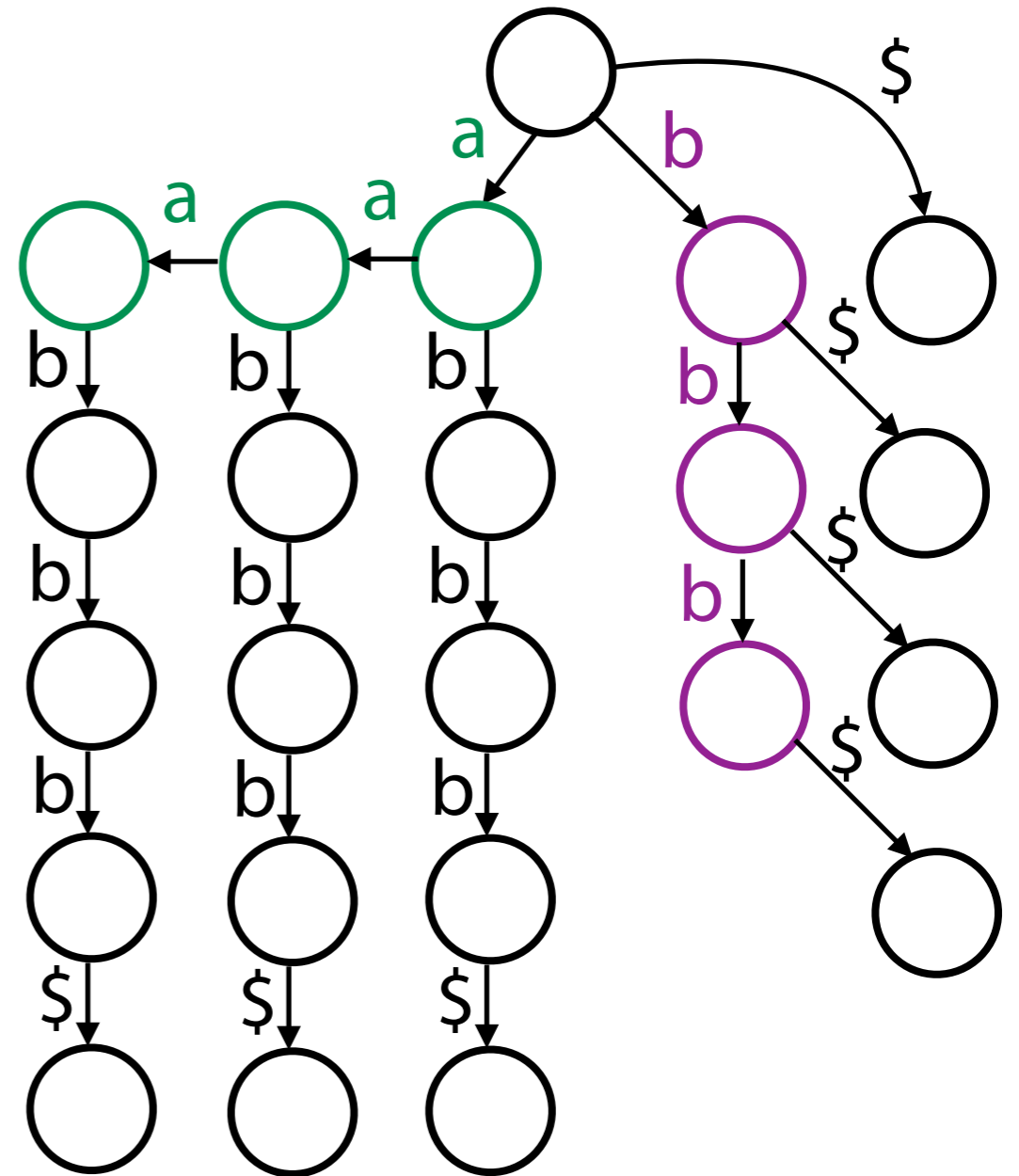
T: **a a a b b b \$**
 a a b b b \$
 a b b b \$
 b b b \$
 b b \$
 b \$
 \$



Suffix trie

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suffix trie grows with $O(m^2)$

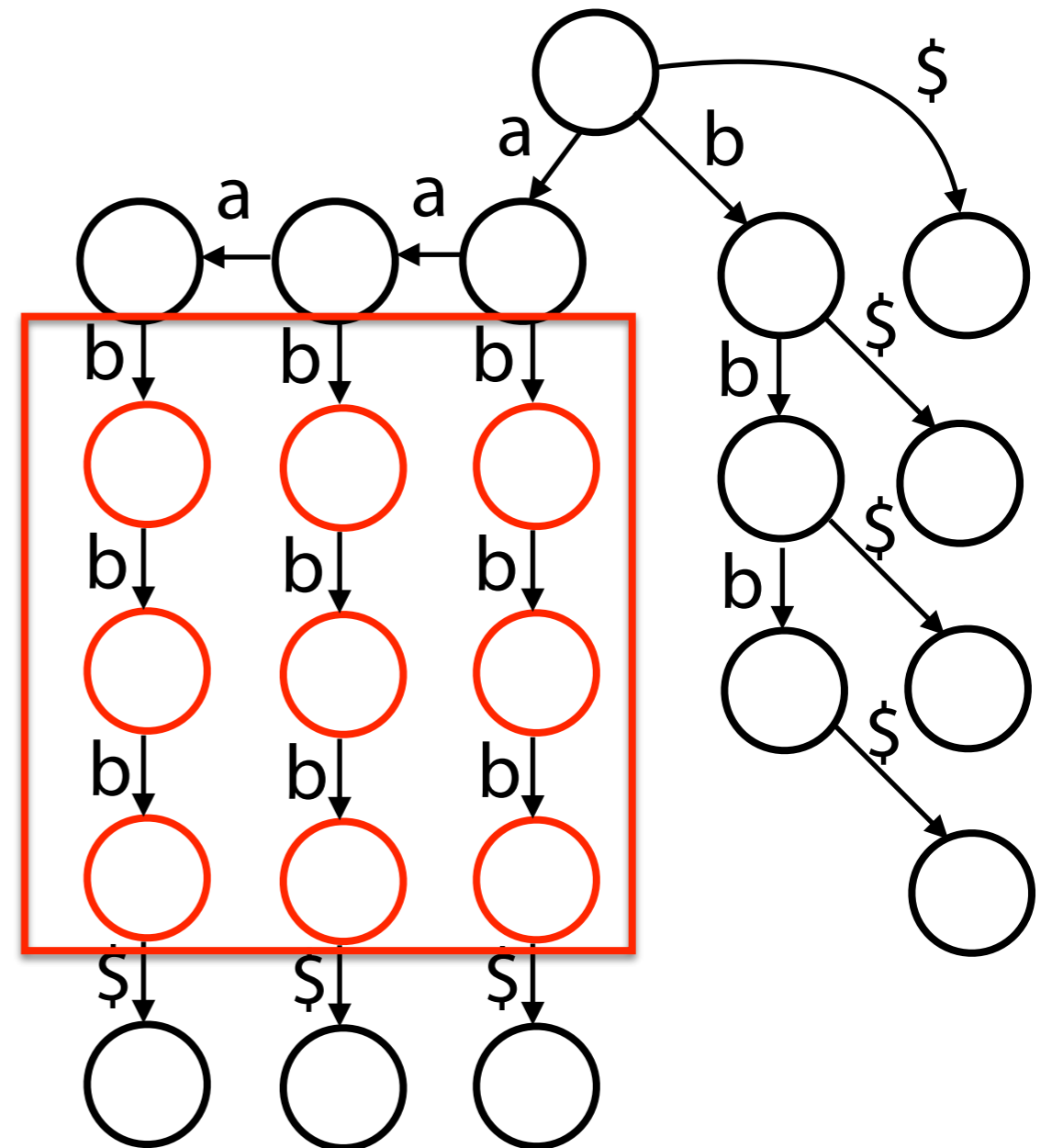
T: **a a a b b b \$**
 a a b b b \$
 a b b b \$
 b b b \$
 b b \$
 b \$
 \$



Suffix trie

Even when alphabet is {a, b},
we can find strings where
suffix trie grows with $O(m^2)$

T: a a a **b b b** \$
a a **b b b** \$
a **b b b** \$
b b b \$
b b \$
b \$
\$

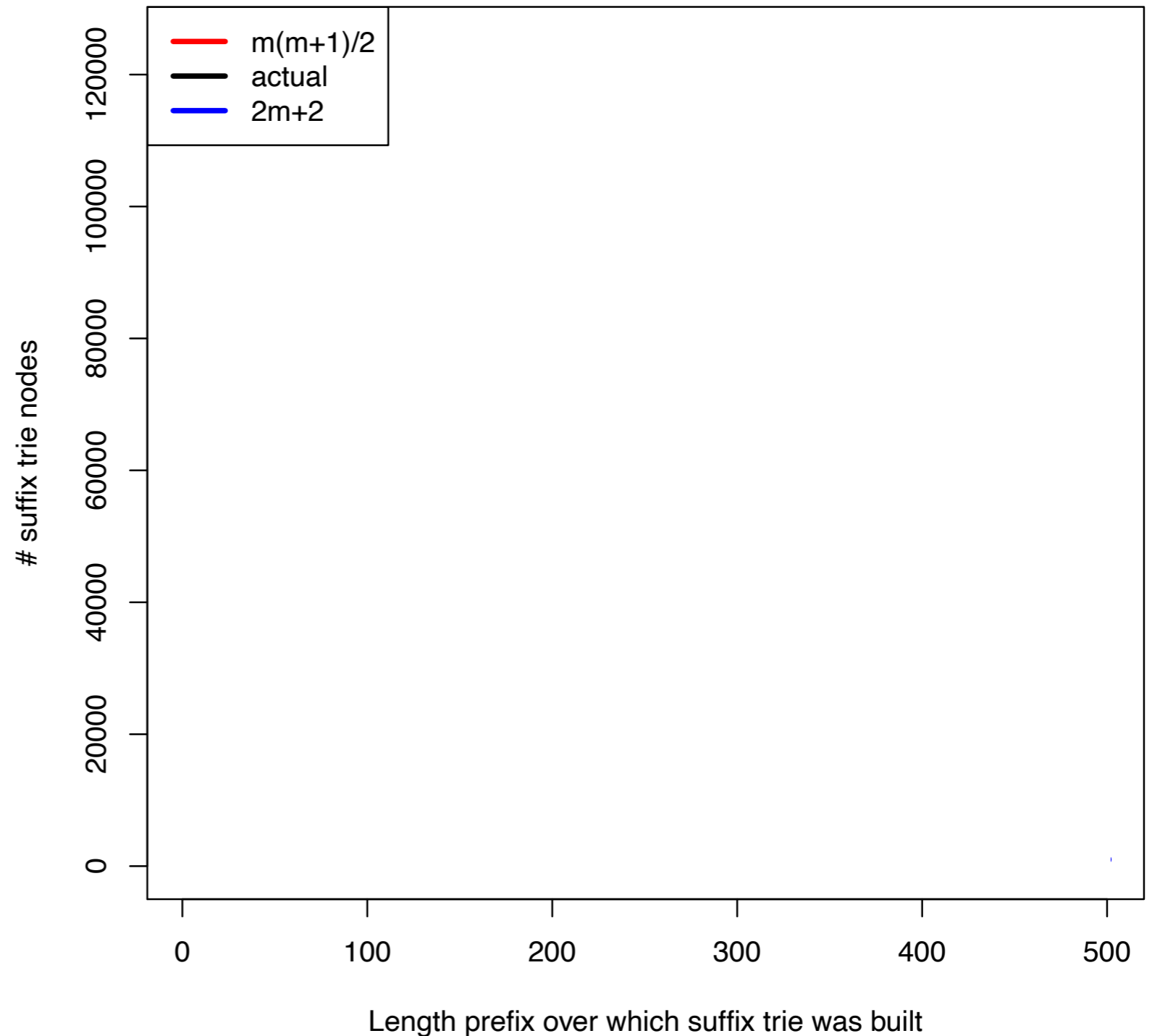


$$\sim \left(\frac{m}{2}\right)^2$$

Suffix trie: actual growth

Built suffix tries for the first 500 prefixes of the lambda phage virus genome

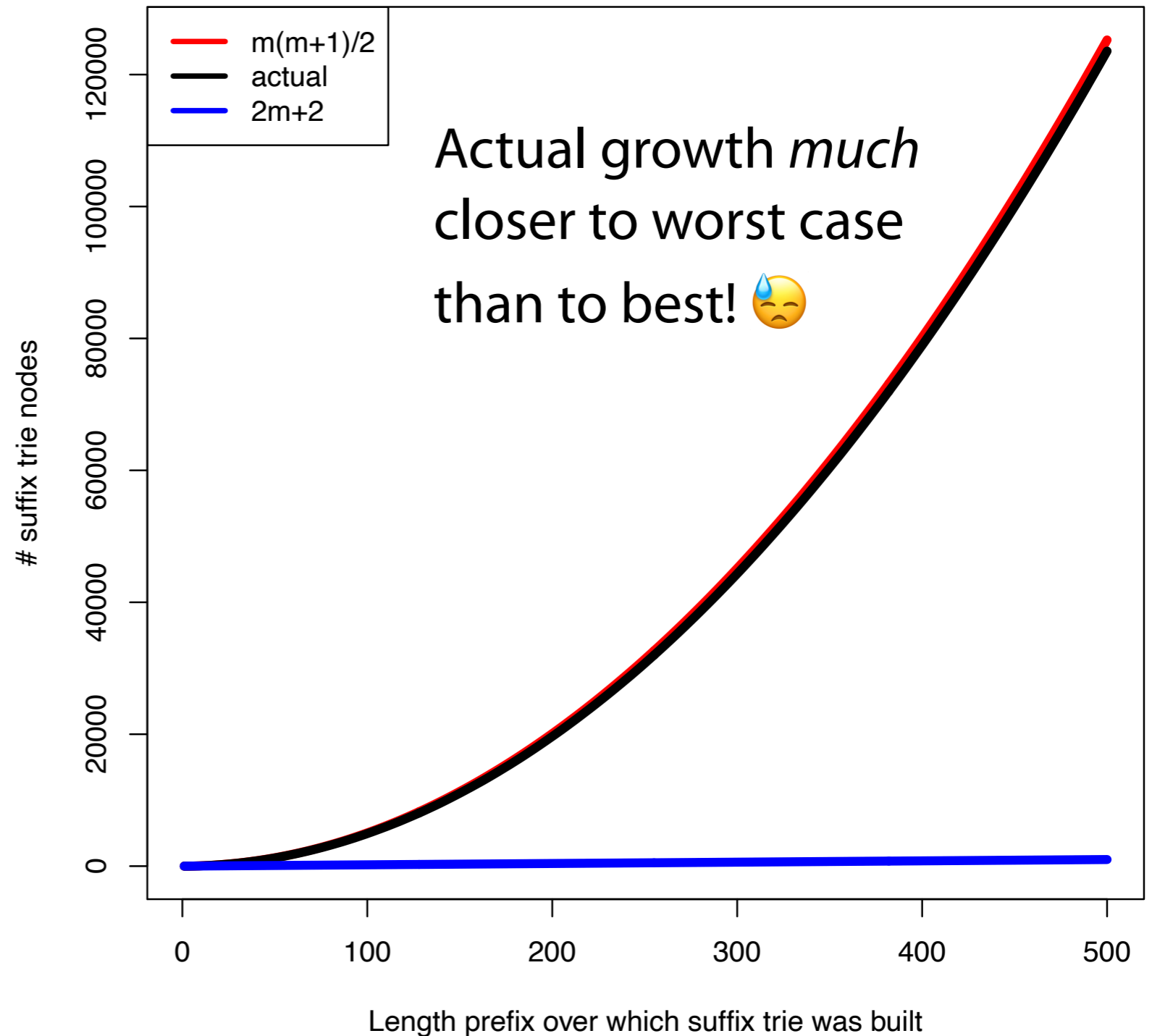
Black curve shows how # nodes increases with prefix length



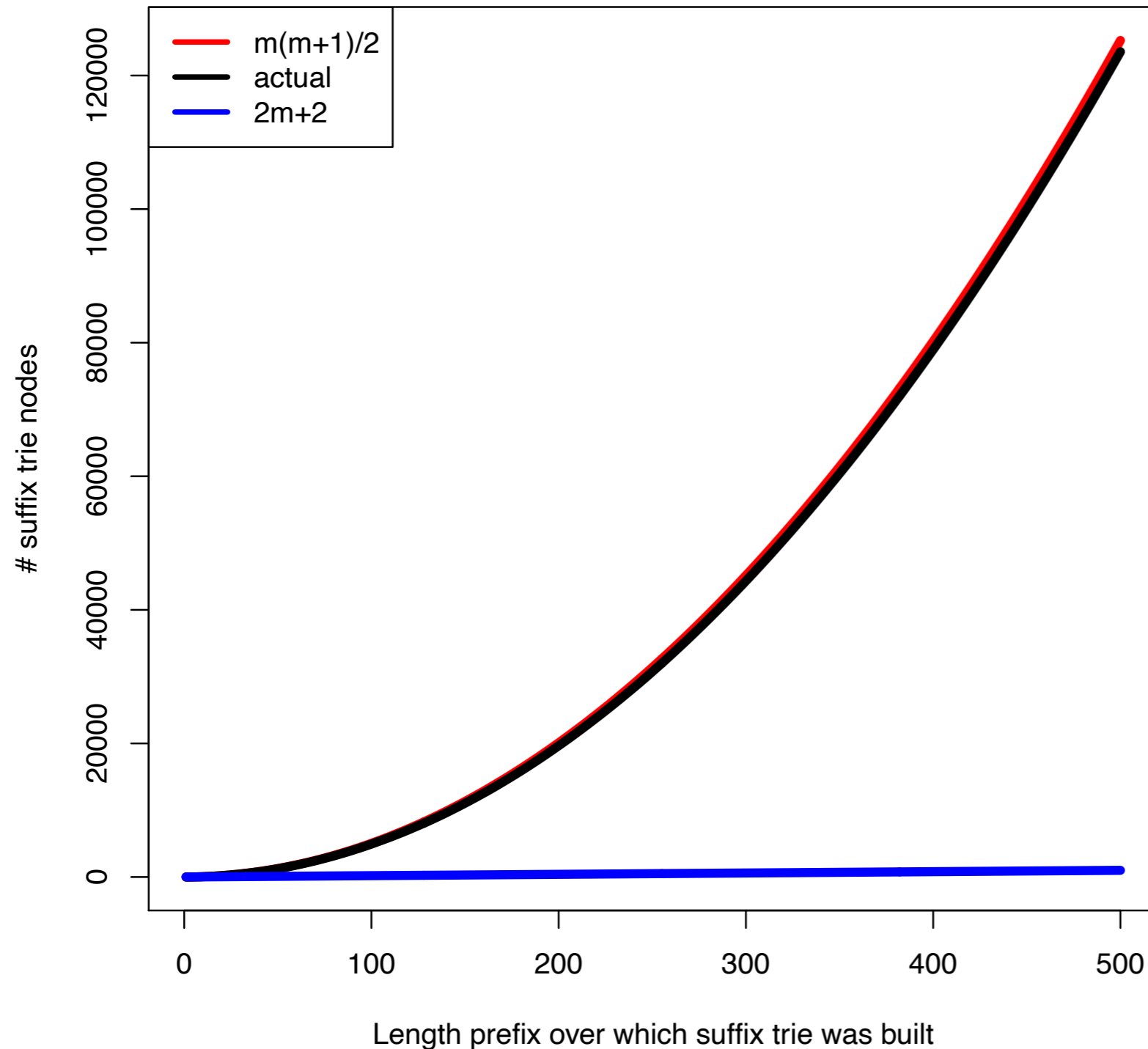
Suffix trie: actual growth

Built suffix tries for the first 500 prefixes of the lambda phage virus genome

Black curve shows how # nodes increases with prefix length



Suffix trie: actual growth



Human genome is $3 \cdot 10^9$ bases long

If $m = 3 \cdot 10^9$, m^2 is far beyond what we can store in memory

Suffix trie

How do we
shrink the trie?

In next video...

