Clark's Select

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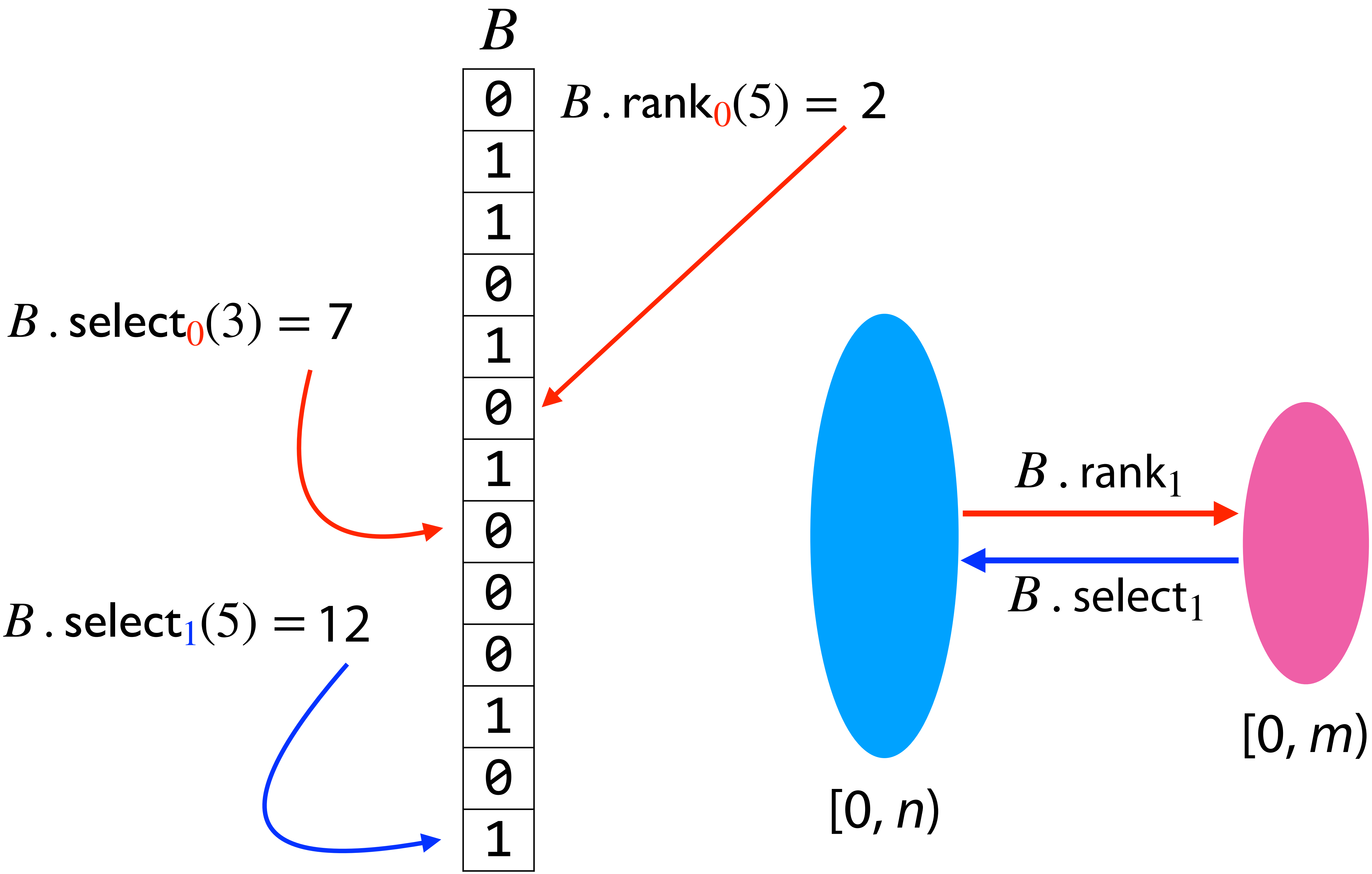
Clark's select

Unlike rank:

- Chunks are defined by # 1s, not # bits
- Two layers of “sparsity”
- Answer is an offset into bitvector

Bitvectors



B
0
1
1
0
1
0
1
0
0
0
1
0
1
$B\,.\,\mathrm{rank}_0(5) = 2$
$B\,.\,\mathrm{select}_0(3) = 7$
$B\,.\,\mathrm{select}_1(5) = 12$
$B\,.\,\mathrm{rank}_1$
$B\,.\,\mathrm{select}_1$
$[0, n)$
$[0, m)$

Clark's select

Split into $\log^2 n$-weight chunks

...

Sparse ($\geq \log^4 n$-length)

Lookup table for each 1-bit ✅

Dense ($< \log^4 n$-length)

Split into $\sqrt{\log n}$-weight sub-chunks

Sparse ($\geq 1/2 \log n$-length)

Lookup table for each 1-bit ✅

Dense ($< 1/2 \log n$-length)

Lookup table for *all possible* sub-chunks ✅

Clark's select

Split the string into chunks each containing $\log^2 n$ 1-bits

Larger chunks are ***sparse***; 1's spread out

Shorter chunks are ***dense***; 1's packed together

Clark's select

Each chunk contains $\log^2 n$ 1-bits

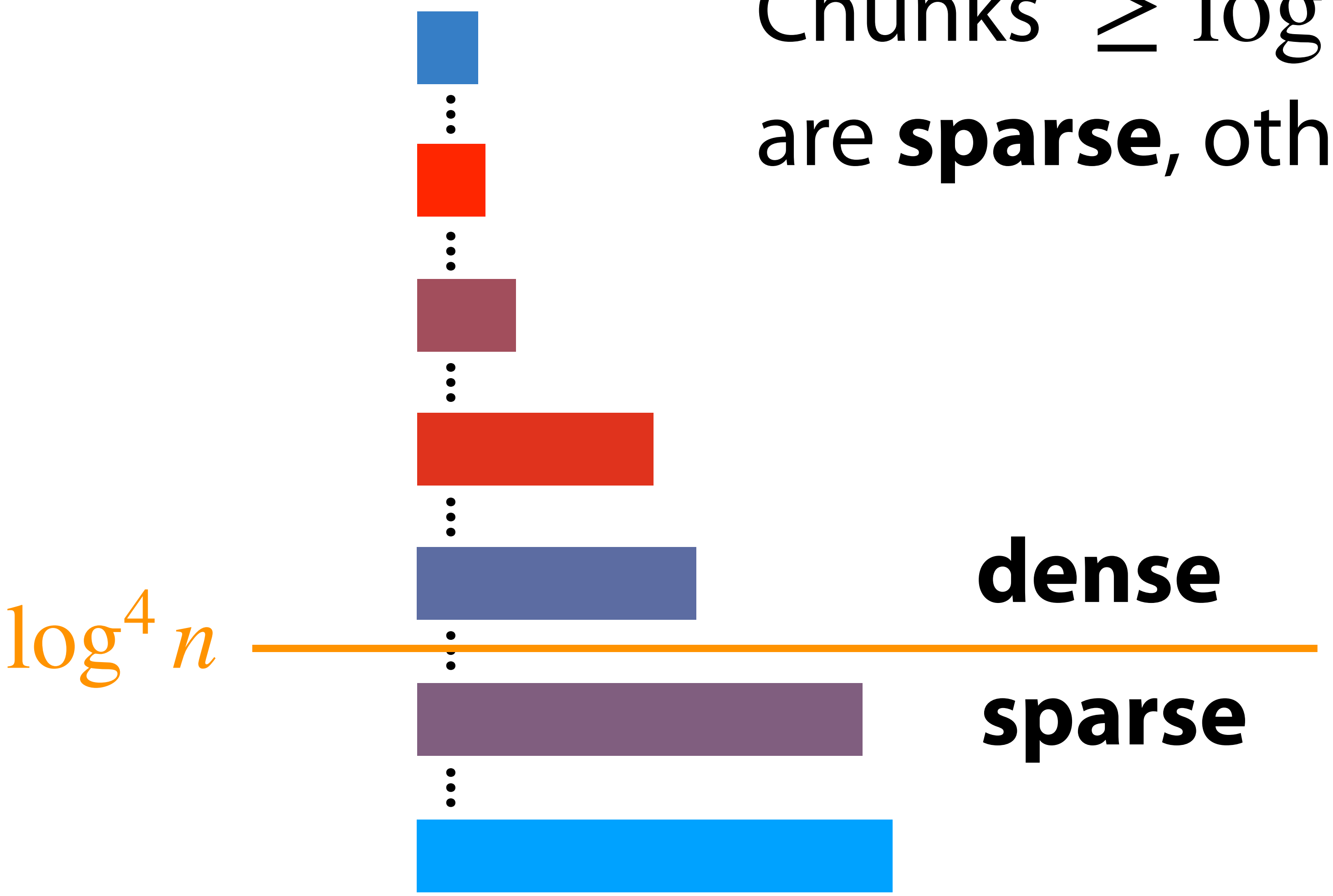
We store offset of each chunk start

This takes:

$$O\left(\frac{n}{\underbrace{\log^2 n}_{\text{Worst case: every bit set}}}\log n\right) = O\left(\frac{n}{\log n}\right) = \check{o}(n) \text{ bits}$$

Clark's select

Chunks $\geq \log^4 n$ bits in length are **sparse**, others **dense**



$\log^4 n$ is square of the # of set bits per chunk, $\log^2 n$

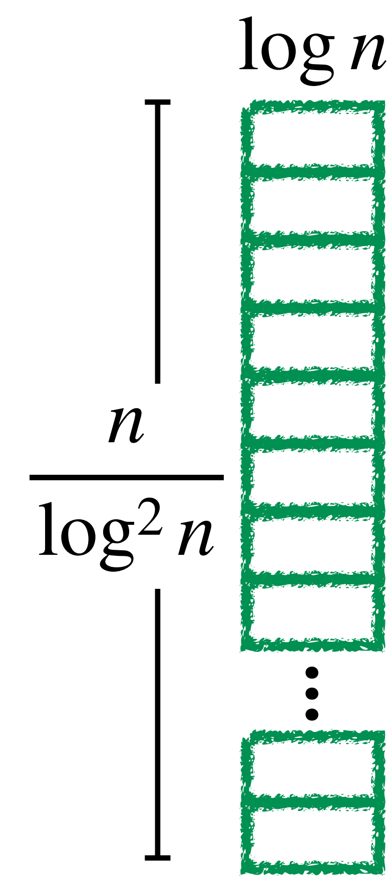
("sparse" roughly means "less than $\sqrt{\text{\# bits}}$ are set")

Clark's select: sparse case

Pre-calculate B . select$_1$ for 1-bits in sparse chunks

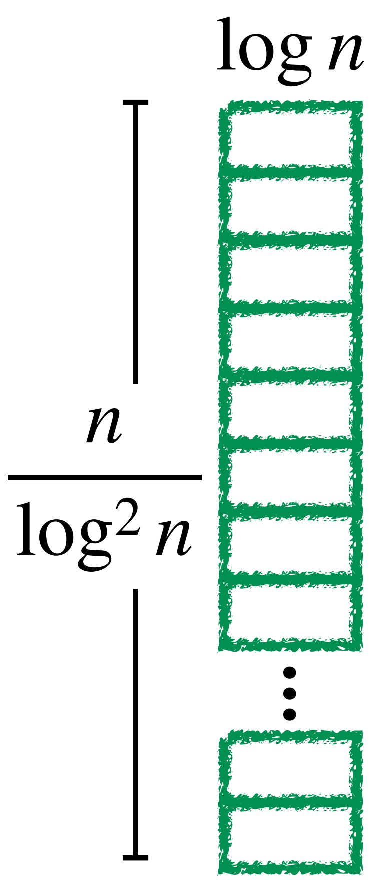
$$O\left(\underbrace{\frac{n}{\log^4 n}}_{\text{Max \# sparse chunks}} \cdot \underbrace{\log n}_{\text{\# bits to store 1 answer}} \cdot \underbrace{\log^2 n}_{\text{\# answers per chunk}}\right)$$

$$= O\left(\frac{n}{\log n}\right) = \check{o}(n)$$



Offsets for 1-bits in sparse chunks

Clark's select



Offsets for 1-bits
in sparse chunks

So far, strategy for select is:

(a) find what chunk it's in (division)

(b) if chunk is **sparse** ($\geq \log^4 n$ bits)

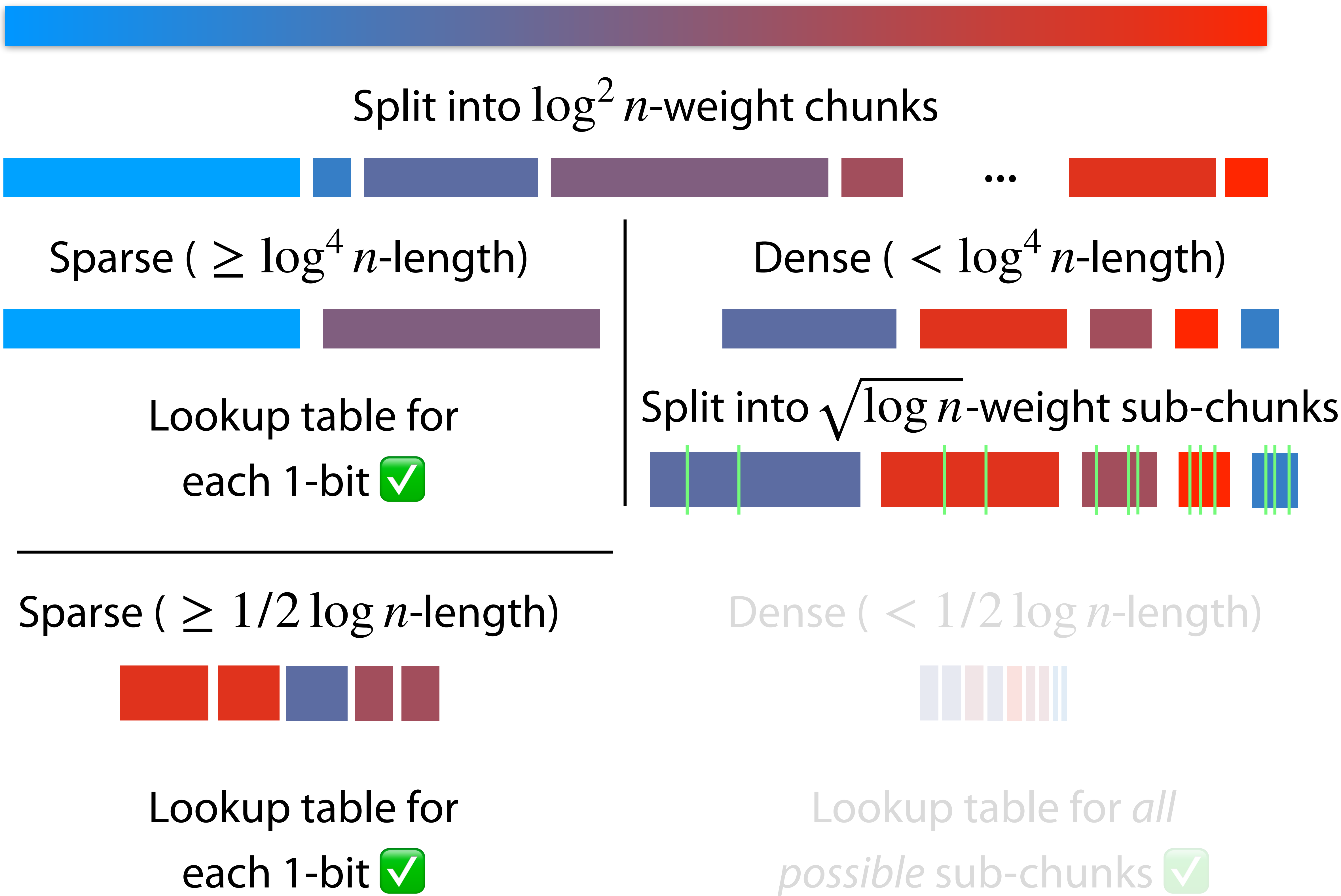
(b.i) look up in sparse offset table

(c) if chunk is **dense** ($< \log^4 n$ bits)

TODO

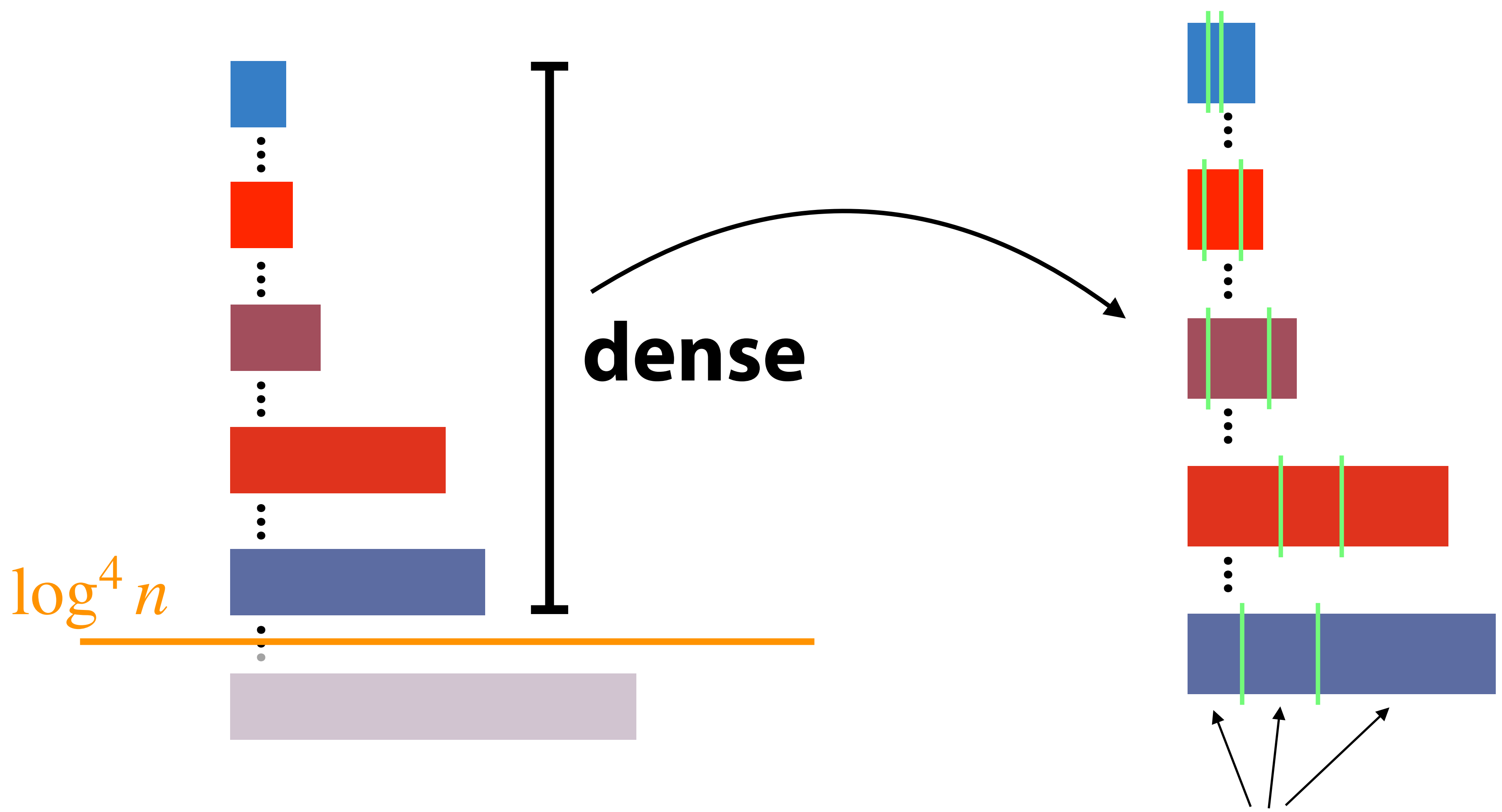
So far, space is $\check{o}(n)$

Clark's select



Clark's select: dense case

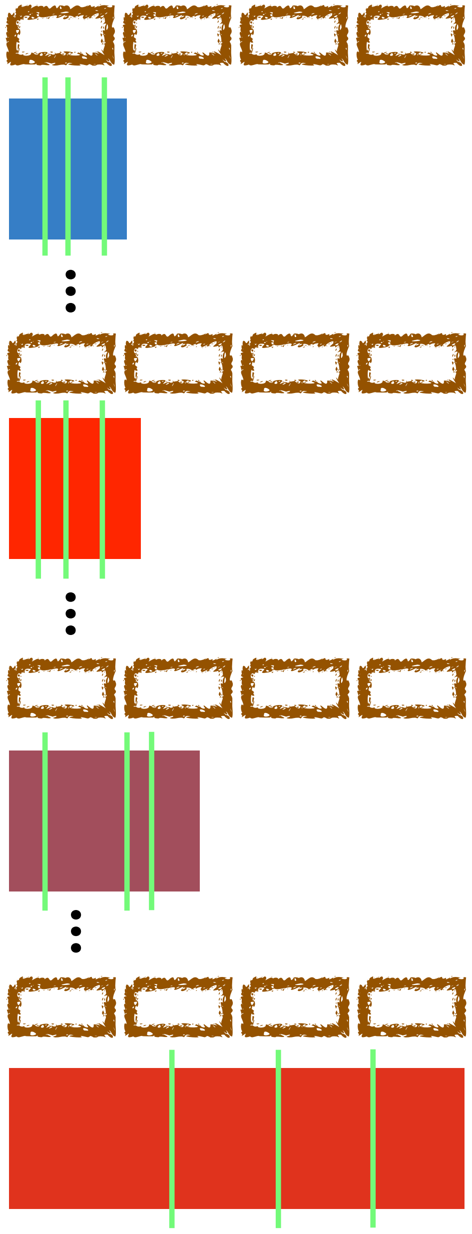
Dense chunks are shorter than $\log^4 n$ bits; further subdivide these to sub-chunks of $\sqrt{\log n}$ 1-bits each



Chunks with $\log^2 n$ 1s

Sub-chunks with $\sqrt{\log n}$ 1s

Clark's select: dense case



Store relative offset per sub-chunk

There are $\leq n/\sqrt{\log n}$ sub-chunks

Containing chunk has $< \log^4 n$ bits, so relative offset fits in $O(\log \log^4 n) = O(\log \log n)$ bits

Overall: $O\left(\dfrac{n \log \log n}{\sqrt{\log n}}\right) = \check{o}(n)$

Clark's select

So far, strategy for select is:

(a) find what chunk it's in (division)

(b) if chunk is **sparse**

(b.i) look up in sparse offset table

(c) if chunk is **dense**

(c.i) look up chunk’s offset

(c.ii) find what sub-chunk it’s in (division by $\sqrt{\log n}$)

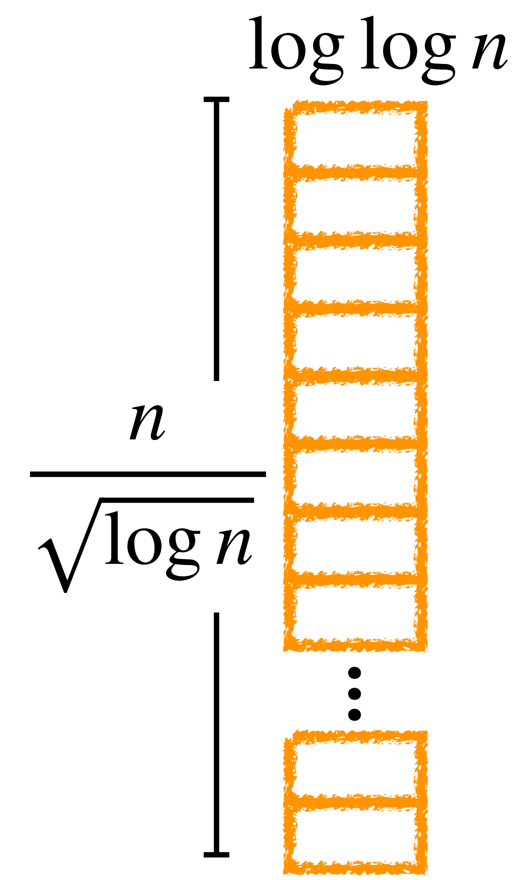
(c.iii) look up sub-chunk’s relative offset

TODO: need to look *within* sub-chunks

Clark's select: dense/sparse case

Sub-chunks with $\geq 1/2 \log n$ bits are **sparse**;
we simply store relative offsets for every 1-bit

Overall: $$O\left(\underbrace{\frac{n}{1/2\log n}}_{\text{Max \# sparse sub-chunks}}\ \underbrace{\log\log n}_{\text{\# bits to store 1 answer (rel. to chunk)}}\ \underbrace{\sqrt{\log n}}_{\text{\# 1-bits per chunk}}\right)$$



$$= O\left(\frac{n\sqrt{\log n}\log\log n}{\log n}\right) = O\left(\frac{n\log\log n}{\sqrt{\log n}}\right) = \check{o}(n)$$

Clark's select

Split into $\log^2 n$-weight chunks

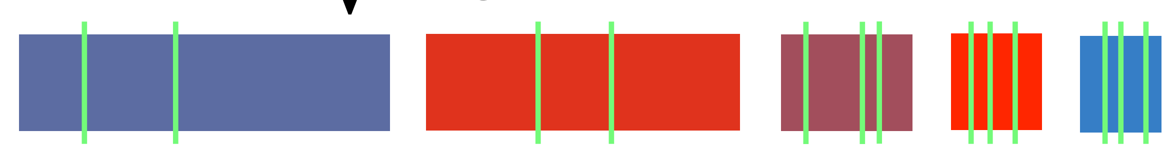
Sparse ($\geq \log^4 n$-length)

Lookup table for

each 1-bit ✅

Dense ($< \log^4 n$-length)

Split into $\sqrt{\log n}$-weight sub-chunks

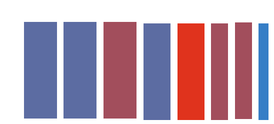


Sparse ($\geq 1/2 \log n$-length)

Lookup table for

each 1-bit ✅

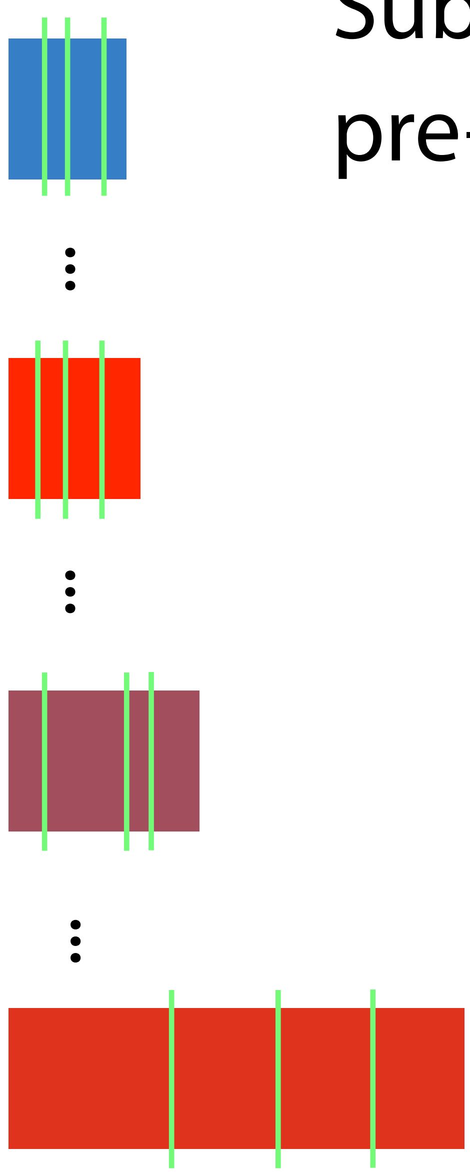
Dense ($< 1/2 \log n$-length)



Lookup table for *all*

possible sub-chunks ✅

Clark's select: dense/dense case



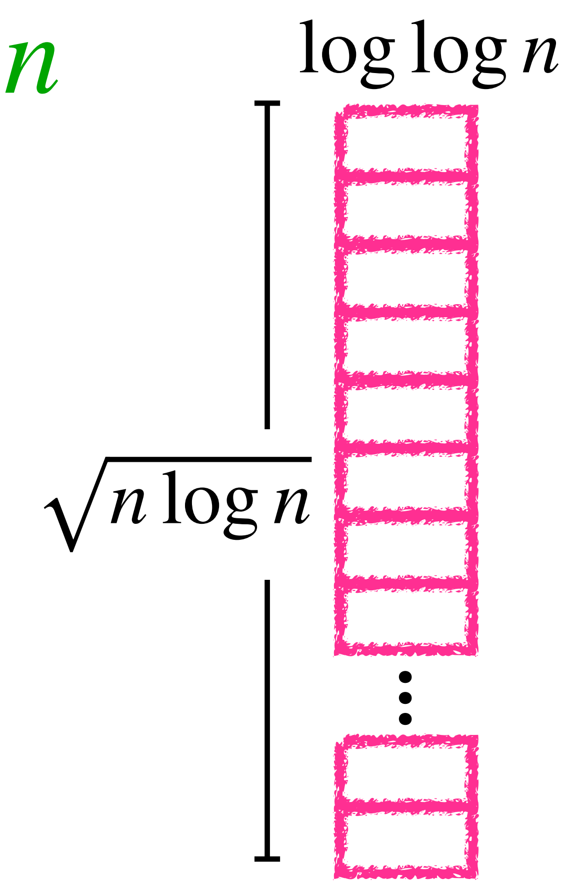
Sub-chunks $< 1/2 \log n$ bits are **dense**;
pre-calculate answers for all such chunks, like rank:

$$2^{1/2 \log n} \cdot \sqrt{\log n} \cdot \log \log n$$

possible bitvectors | possible 1-bits | answer

$$= O\left(\sqrt{n \log n} \log \log n\right)$$

$$= \check{o}(n)$$



Clark's select

(a) find what chunk it's in (division by $\log^2 n$)

(b) if chunk is **sparse** ($\geq \log^4 n$ bits)

(b.i) look up answer in sparse offset table

(c) if chunk is **dense** ($< \log^4 n$ bits)

(c.i) look up chunk's offset

(c.ii) find what sub-chunk it's in (divide by $\sqrt{\log n}$)

(c.iii) look up sub-chunk's relative offset

(c.iv) if sub-chunk is **sparse** ($\geq 1/2 \log n$ bits)

(c.iv.A) look up answer in sparse 1-bit table

(c.iv.B) return **(c.i)** + **(c.iii) + (c.iv.A)**

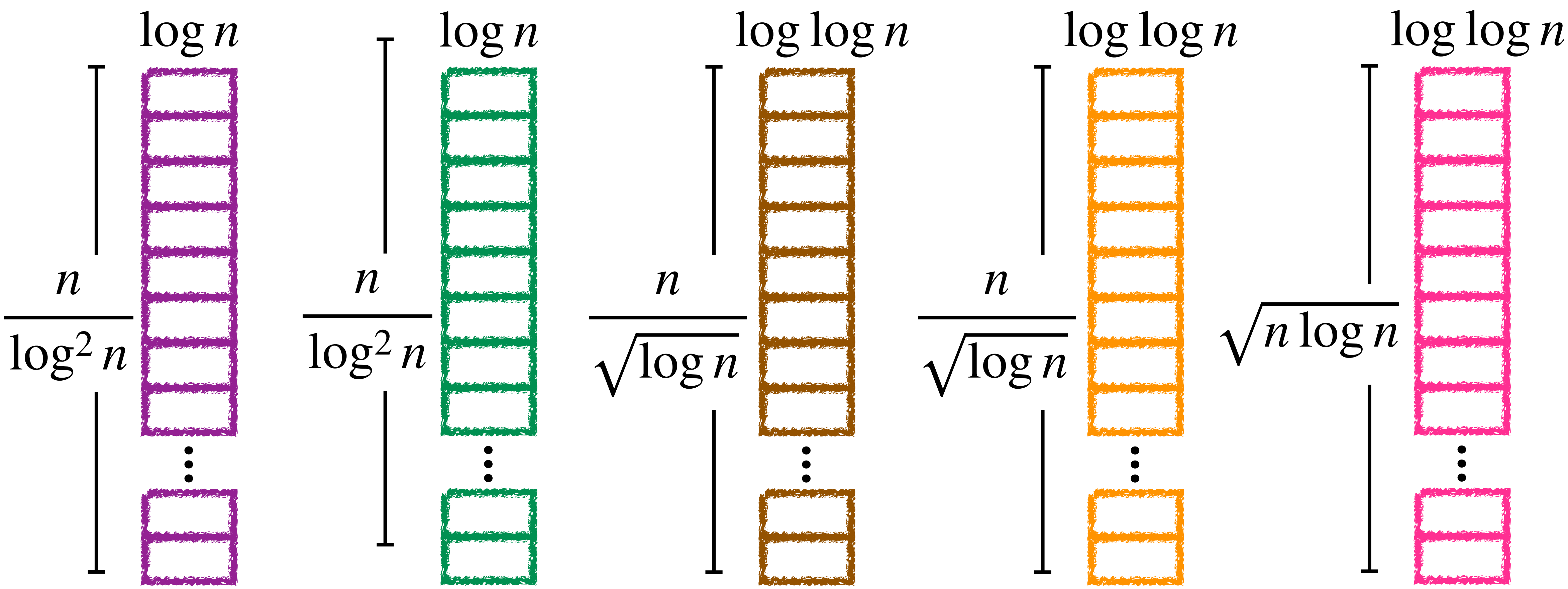
(c.v) if sub-chunk is **dense**

(c.v.A) look up answer in all possible dense/dense table

(c.v.B) return **(c.i)** + **(c.iii) + (c.v.A)**

Clark's select

Overall, space is $\check{o}(n)$



Sparse chunk offsets

Offsets for 1-bits in sparse chunks

Dense sub-chunk relative offsets

Answers for 1-bits in dense/sparse sub-chunks

Answers for all possible dense/dense sub-chunks

Bitvectors

	Time	Space (bits)	Note
B . access	$O(1)$	n	Lookup
B . select$_1$	$O(1)$	$\check{o}(n)$	✅ Clark
B . rank$_1$	$O(1)$	$\check{o}(n)$	✅ Jacobson