

SVM with Multiclass

Note Title

11/25/2006

Task: learn a classifier $h: \mathcal{X} \rightarrow \mathcal{Y}$

Hypothesis: $h_{\omega}(x) = \arg \max_y \sum_{i=1}^n \omega_i \cdot f_i(x, y) = \arg \max_y \omega^T f(x, y)$

Task: select weights $\{\omega_i\}$ using training data:
 $\{(\underline{x}^{(i)}, y^{(i)} = t(\underline{x}^{(i)})) : i = 1 \text{ to } N\}$ $t(x) = \text{true label of } x$.

Define: $\Delta f_x(y) = f(x, t(x)) - f(x, y)$

Minimize $\frac{1}{2} \|\omega\|^2 + C \sum_x \xi_x$ $\forall x, y$

s.t. $\omega \cdot \Delta f_x(y) \geq \Delta t_x(y) - \xi_x$

Condition $\Rightarrow \xi_x \geq 0$, by setting
 $y = t(x)$, $\Delta f_x(t(x)) = 0 = \Delta t_x(t(x))$.

$\Delta t_x(y) = 1$, if $y \neq t(x)$
 0 , otherwise.

Primal Problem

$L_p(\omega, \xi_x : \alpha_x)$

$= \frac{1}{2} \|\omega\|^2 + C \sum_x \xi_x$

$- \sum_{x, y} \alpha_x(y) (\omega \cdot \Delta f_x(y) - \Delta t_x(y) + \xi_x)$

Dual Problem

$L_d(\alpha_x) = \sum_{x, y} \alpha_x(y) \Delta t_x(y)$

$- \frac{1}{2} \left\| \sum_{x, y} \alpha_x(y) \Delta f_x(y) \right\|^2$

s.t. $\sum_y \alpha_x(y) = C$, $\forall x$. $\alpha_x(y) \geq 0$.

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Special Case: Markov Network.

E.G.

a | b | c | d | e | f

$h: \mathcal{X} \rightarrow \mathcal{Y}$
Dataset

$$\mathcal{Y} = \mathcal{Y}_1 \times \dots \times \mathcal{Y}_k$$

$$S = \{(\underline{x}^i, \underline{y}^i) : i = 1, \dots, n\}$$

Let $\underline{y}^i$ be multivalued.

in form: $\underline{y}^i = (y^i_1, \dots, y^i_k)$

with $y^i_a \in \{\pm 1\}$, for each a .

basis function:

$$f_j: \mathcal{X} \times \mathcal{Y} \rightarrow \mathbb{R}.$$

classification function:

$$h_{\underline{\omega}}(\underline{x}) = \arg \max_{\underline{\omega}} \sum_{i=1}^n \omega_i f_i(\underline{x}, \underline{y})$$

E.G. OCR example

Each $\mathcal{Y}_i$ is a character

$\mathcal{Y}$ is a full word.

webpage classification:

Each $\mathcal{Y}_i$ is a webpage label.

$\mathcal{Y}$ is a label for an entire website.

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Set of features.

$$f_{\underline{x}}(\underline{y}) = \sum_{(i,j) \in E} f_{\underline{x}}(\underline{y}_i, \underline{y}_j)$$

penalty:

$$\Delta t_{\underline{x}}(\underline{y}) = \sum_{i=1}^l \Delta t_{\underline{x}}(y_i)$$

 E - edges of Graph.with $\Delta t_{\underline{x}}(y_i) \equiv \mathbb{I}(y_i \neq (t(\underline{x}))_i)$.Primal Problem

$$\min \frac{1}{2} \|\underline{\omega}\|^2 + C \sum_{\underline{x}} \underline{\xi}_{\underline{x}}$$

$$\text{s.t. } \underline{\omega}^T \Delta f_{\underline{x}}(\underline{y}) > \Delta t_{\underline{x}}(\underline{y}) - \underline{\xi}_{\underline{x}}, \quad \forall \underline{x}, \underline{y}.$$

Dual Problem

$$\max \sum_{\underline{x}, \underline{y}} \alpha_{\underline{x}}(\underline{y}) \Delta t_{\underline{x}}(\underline{y}) - \frac{1}{2} \left\| \sum_{\underline{x}, \underline{y}} \alpha_{\underline{x}}(\underline{y}) \Delta f_{\underline{x}}(\underline{y}) \right\|^2$$

$$\text{s.t. } \sum_{\underline{y}} \alpha_{\underline{x}}(\underline{y}) = C, \quad \forall \underline{x}, \quad \alpha_{\underline{x}}(\underline{y}) \geq 0, \quad \forall \underline{x}, \underline{y}.$$

Now we can interpret $\alpha_{\underline{x}}(\underline{y})$ as a probability distributor $\rightarrow$ with normalization C .

The dual function is a function of the expectation of $\Delta t_{\underline{x}}(\underline{y})$ & $\Delta f_{\underline{x}}(\underline{y})$.

$$\Delta t_{\underline{x}}(\underline{y}) = \sum_i \Delta t_{\underline{x}}(y_i), \quad \Delta f_{\underline{x}}(\underline{y}) = \sum_{(i,j) \in E} \Delta f_{\underline{x}}(y_i, y_j)$$

(4) The dual problem reduces to a formulation on the marginal unary and pairwise distributions

$$\mu_x(y_{ij}) = \sum_{y \sim [y_{ij}]} \alpha_x(y), \quad \forall (ij) \in E$$

$$\forall y_{ij} : \exists x$$

$$\mu_x(y_i) = \sum_{y \sim [y_i]} \alpha_x(y), \quad \forall i, \forall y_i, \forall x.$$

$$y \sim [y_i, y_j]$$

full assignment

y consistent with
partial assigned

y_i, y_j .

Must enforce consistency
constraints between the $\mu_x(y_{ij})$
and $\mu_x(y_i)$

$$\sum_{y_i} \mu_x(y_{ij}) = \mu_x(y_i), \quad \forall y_j, \forall (ij) \in E.$$

Because:

$$\sum_y \alpha_x(y) \Delta_x(y) = \sum_y \sum_i \alpha_x(y) \Delta_{t_x}(y_i) = \sum_{i, y_i} \Delta_{t_x}(y_i) \sum_{y \sim [y_i]} \alpha_x(y)$$

$$= \sum_{i, y_i} \mu_x(y_i) \Delta_{t_x}(y_i)$$

similarly, the second term can be expressed in form

$$-\frac{1}{2} \sum_{x, x'} \sum_{(ij)} \sum_{y_i, y_j} \mu_x(y_{ij}) \mu_{x'}(y_{ij}) \Delta f_x(y_i, y_j)^T \Delta f_{x'}(y_i, y_j) + \sum_x \sum_{i, y_i} \mu_x(y_i) \Delta_{t_x}(y_i).$$

Algorithm can maximize this, by variant of
Belief Propagation

Go to 7

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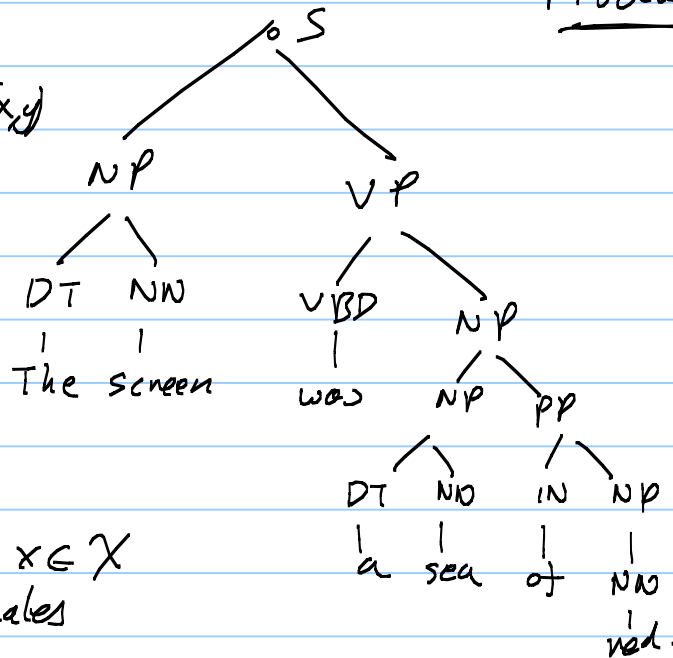
Max-Margin Parsing

Production Rules

$S \rightarrow (NP, VP)$

$$f_{w_0}(x) = \arg \max_{y \in G(x)} w_0 \cdot \Phi(x, y)$$

Input x ,
 $G(x)$ is a set of
parses.



G maps an input $x \in X$
to a set of candidate
parses $G(x) \subseteq Y$

$$L: X \times Y \times Y \rightarrow \mathbb{R}^+$$

$L(x, y, \hat{y})$ is the penalty for making parse y
on data x when tree parse is $\hat{y}$.

Want to estimate a linear discriminant form:

$$f_{w_0}(x) = \arg \max_{y \in G(x)} w_0 \cdot \Phi(x, y)$$

The arg max can be performed by
dynamic programming.

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$$y_i = \arg \max_{y \in G(x_i)} \underline{\omega} \cdot \underline{\Phi}(x_i, y)$$

The margin of the parameter $\underline{\omega}$ on example i and parse y

$$\underline{\omega} \cdot \underline{\Phi}(x_i, y_i) - \underline{\omega} \cdot \underline{\Phi}(x_i, y) = \underline{\omega} \cdot (\underline{\Phi}_{i, y_i} - \underline{\Phi}_{i, y})$$

Primal.

$$\min \frac{1}{2} \|\underline{\omega}\|^2 + C \sum_i \xi_i$$

y_i is the true parse.

$$\text{s.t.} \quad \underline{\omega} \cdot (\underline{\Phi}_{i, y_i} - \underline{\Phi}_{i, y}) \geq L_{i, y} - \xi_i \quad \forall y \in G(x_i)$$
$$L_{i, y} = L(x_i, y, y_i)$$

Dual.

$$\max C \sum_{i, y} \alpha_{i, y} L_{i, y} - \frac{1}{2} \left\| \sum_{i, y} (I_{i, y} - \alpha_{i, y}) \underline{\Phi}_{i, y} \right\|^2$$

$$\text{s.t.} \quad \sum_y \alpha_{i, y} = 1, \quad \forall i, \quad \alpha_{i, y} \geq 0, \quad \forall i, y$$

$I_{i, y} = \mathbb{I}(x_i, y, y_i)$ indicates whether y is the true parse y_i .

$$\underline{\omega}^* = C \sum_{i, y} (I_{i, y} - \alpha_{i, y}^*) \underline{\Phi}_{i, y}$$

(At 5)

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PCFG

Chomsky Normal Form (CNF)

$A \rightarrow BC$

A, B, C non-terminal symbols

$A \rightarrow a$

a terminal symbol

Two types of parts

First type

$\langle A, s, e, i \rangle$

non-terminal A

e.g.

start-point s

NP, 3, 5 "a sea"

end-point e

sentence i .

e.g.

$S \rightarrow NP, VP, 0, 2, 7$

Second type:

$A \rightarrow BC, s, m, e, i$

split point m

Countable set of parts R ,

$R(x, y)$ is the set of parts belonging to a particular parse

All rules are in binary-branching form
so $|R(x, y)|$ is constant across different derivations y
for the same input x

$$\Phi(x, y) = \sum_{r \in R(x, y)} \phi(x, r)$$

$$L(x, y, \hat{y}) = \sum_{r \in R(x, y)} l(x, y, r)$$

One possibility - define $L(x, y, v) = 0$ only if
the non-terminal A spans words $s_1 \dots s_n$
in the derivation y and 1 otherwise.

Indicator variable $I(x, y, v) = 1$, if $v \in R(x, y)$
 $= 0$, otherwise

Dual:

$$C \sum_i E_{\alpha_i} \|L_i(y) - \frac{1}{2}\|^2 \leq \sum_i \|\phi_i(y_i) - E_{\alpha_i}[L_i(y)]\|^2$$