VoteBox: a verifiable, tamper-evident electronic voting system

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Talk outline

Background

Trustworthiness of electronic voting machines Why it's worth improving them

The design of VoteBox

Durability and audit

Ballot casting assurance

Beyond

Background



DRE voting machines (Direct Recording Electronic)







US Presidential election (2000)

HAVA (2002)

DREs discredited

High-profile failures in real elections

A few examples:

2006: Sarasota, FL undervoting ~18,000 ballots blank in the congressional race (~15%) margin of victory: 369 votes

2008: video documentation of "vote flipping" touch-screen calibration? buggy input filters?

Ongoing: long lines due to complex setup, equipment problems, etc.



DREs discredited

Software bugs & design flaws identified by e-voting researchers

2003 Analysis of Diebold AccuVote TS Leaked source code analyzed [Kohno et al. 2004] Poor software engineering, incorrect cryptography, vulnerable to malicious upgrades, multiple voting



2006 "Voting-machine virus" developed Self-propagating malicious upgrades that spread from machine to machine, altering votes and leaving no trace [Feldman et al. 2006]

DREs discredited

Software bugs & design flaws identified by e-voting researchers

2007 Involvement by computer scientists in statewide voting systems audits

groundbreaking access to source code of commercial voting systems

Top-To-Bottom Review (California)

- All machines certified for use in CA found to have serious bugs & be vulnerable to attack
- Viral-style attacks found in all systems
- **EVEREST** study (Ohio)
 - All machines certified in OH found vulnerable (validating CA-TTBR)
 - Showed that hundreds of votes were lost in 2004



malfunctions

could result in changed or lost votes

design flaws

could let attackers alter the election outcome without leaving evidence



Result: undermined trust in elections





voters prefer electronic voting

S. P. Everett, K. K. Greene, M. D. Byrne, D. S. Wallach, K. Derr, **D. R. Sandler**, and T. Torous. *Electronic voting machines versus traditional methods: Improved preference, similar performance.* In CHI 2008.

legitimate benefits accessibility feedback flexibility satisfaction

Can we design a line terms of the terms of terms of the terms of te

"better" = ?

1. resistance to failure & tampering

essential vote data should survive hardware failure, poll worker mistakes, attempts to attack the system

2. tamper-evidence

if we are unable to prevent data loss, we must always be able to detect the failure

3. verifiability

two useful properties:

cast-as-intended

"Was my vote recorded faithfully?" very, very hard for DREs to satisfy

counted-as-cast

"Has my vote been tallied correctly?" can be somewhat addressed with recounts

resistance to failure & tampering

prevent or minimize data loss

tamper-evidence

if resistance is futile

verifiability

cast-as-intended; counted-as-cast

(DRE user experience)

a computer science problem

resistance to failure & tampering
replication; gossip
tamper-evidence
secure logs
verifiability
cryptography

Auditorium

Ballot challenge





A field study.





March 7, 2006:

Democratic primary election (County's first use of DREs)

An unusual situation

Voters given a choice:



OR



DRE (ES&S iVotronic)

Paper (central ES&S op-scan)

Flores v. Lopez

~50,000 votes cast
Margin of victory: ~100 votes
The loser suspected the DREs
...because he looked better on paper
Lawsuit
Bring in the experts.



initial investigation: gathering data

(April 2006)

Webb Co. data

Raw binary data from Compact Flash cards Opaque, undocumented format Text output from tabulation process IMAGELOG.TXT (cast ballots) EVENTLOG.TXT (more on that later)

What we found

Asmoking gun? Evil voting machines? HACKS??

inherently difficult to find evidence with DREs!

What we (really) found

Anomalies in the event logs

Per-machine records

Captured during machine run time

Transferred to tabulator (IMAGELOG.TXT)

A timeline of important election events

e.g. "terminal open," "ballot cast," ...

Example event log

Votronic	PEB#	Туре	Date	Time	Event
5140052	161061	SUP	03/07/2006	15:29:03	01 Terminal clear and test
	160980	SUP	03/07/2006 03/07/2006 03/07/2006	15:31:15 15:34:47 15:36:36	09 Terminal open 13 Print zero tape 13 Print zero tape
	160999	SUP	03/07/2006 03/07/2006 03/07/2006 03/07/2006 03/07/2006 03/07/2006 03/07/2006	15:56:50 16:47:12 18:07:29 18:17:03 18:37:24 18:41:18 18:46:23	20 Normal ballot cast 20 Normal ballot cast 20 Normal ballot cast 20 Normal ballot cast 22 Super ballot cancel 20 Normal ballot cast 20 Normal ballot cast
	160980	SUP	03/07/2006	19:07:14	10 Terminal close

Problem #1 Logs starting mid-day

03/07/2006 15:29:03 Terminal clear and test 03/07/2006 15:31:15 Terminal open

Polls opened around 7 AM across Webb Co. What happened between 7 and 3:30? Lost votes?

Problem #2 Election events on wrong day

A normal voting pattern:

Votronic	PEB#	Туре	Date	Time	Event	
5142523	161061 161115 161109	SUP SUP SUP	02/26/2006 03/06/2006 03/06/2006 03/06/2006 03/06/2006	19:07:05 06:57:23 07:01:47 07:03:41 10:08:26	01 Terminal clear and 09 Terminal open 13 Print zero tape 13 Print zero tape 20 Normal ballot cast	test
			[9 more	ballots cast]	
	161115	SUP	03/06/2006 03/06/2006	19:29:00 19:29:00	27 Override 10 Terminal close	

The election was held on 03/07! Ballot box stuffing the day before?

A different pattern:



26 machines with exactly two ballots cast the day before (always for the same guy)

We learned that these might be "logic and accuracy test" votes, erroneously included in the tally



initial investigation

follow-up trip: direct inspection




History for Laredo, TX

Tuesday, April 25, 2006 — View Current Conditions

Daily Summary	1			
« Previous Day	April	; 25 ; 2	2006 🗧 Go	Next Day »
Daily	Weekly	Monthly	Custom	1
	Actual:		Average :	Record :
Temperature:				
Mean Temperatu	ure 87 °F / 30	°C	-	
Max Temperatur	e 101 °F/3	8 °C	85 °F / 29 °C	101 °F / 38 °C (2006)
Min Temperature	9 73 °F/22	°C	64 °F / 17 °C	55 °F / 12 °C (2001)

source: wunderground.com

Findings

Machines containing only two votes

Hardware clock appeared normal Most likely L&A test votes

Others

Hardware clock set incorrectly ...just enough to account for anomaly This is not proof of correct behavior!

Problem #3 Insufficient audit data

We couldn't collect data from every machine Many were cleared after the election (Poll workers not supposed to do this!) Paper records missing Zero tapes Cancelled ballot logs Procedural errors by administrators, pollworkers (but the machines didn't help)

"Mistakes were made."

Mistakes were made

Violations of election procedures

Counting test votes in final results Loss of zero tapes and other paper logs Erasure of some machines

Local (mis)configuration

Hardware clocks set wrong

These things cast doubt on the results

Honest mistakes or illegitimate votes?

No way to be sure. Believable audits impossible.

Research goals

Make it easier to audit results after the election Make it harder to make mistakes on election day In particular:

Prove

every vote tallied is valid every valid vote is present

Tolerate

accidental loss/deletion of records election-day machine failure





Connect the machines together.

"The Auditorium"

Auditorium's approach

Store everything everywhere

Massive **redundancy**

Stop trusting DREs to keep their own audit data

Link all votes, events together

Create a **secure timeline** of election events Tamper-evident proof of each vote's legitimacy

D. Sandler and D. S. Wallach. **Casting Votes in the Auditorium.** In Proceedings of the 2nd USENIX/ACCURATE Electronic Voting Technology Workshop (EVT'07).

Ingredient: hash chains

"Machine turned on" (HASH = 0x1234)
"Cast a vote after event 0x1234" (HASH = 0xABCD)
"Cast a vote after event 0xABCD" (HASH = 0xBEEF)
"Turned off after event 0xBEEF" (HASH = 0x4242)

A hash-chained secure log

Every event includes the cryptographic hash (e.g. SHA1) of a previous event [Schneier & Kelsey '99]

Result: provable order

If Y includes H(X), then Y must have happened after X

Any individual change to the log

invalidates all later hashes (breaks the chain)

To alter, insert, or delete a single record

you must alter every subsequent event as well

Ingredient: timeline entanglement

Entanglement = "chain with hashes from others" Result: event ordering between participants

[Maniatis & Baker '02]



Malicious machines can't retroactively alter their own logs

it would violate commitments they have already exchanged with others

Ingredient: broadcast

All-to-all communication

All messages signed & sent to every VoteBox Each machine records each message independently \rightarrow massive replication $O(N^2)$, but *N* is small in a polling place

Mechanism for entanglement

When publishing a new message, include hashes of recent messages in the log (regardless of their origin!)

Broadcast entanglement = Auditorium

A pragmatic benefit

The supervisor console

Assistance for poll workers

Shows status of all machines Votes cast, battery running low, etc.

Helps conduct the election

Open/close polls, authorize machines to cast ballots

Less opportunity for poll-worker error

Ballots distributed over the network

Booths are **stateless**, interchangeable (Supervisor can have a spare as well)





Voting in the Auditorium



voting machines connected in a private polling-place network all election events are signed and broadcast each broadcast is logged by every machine isolated failures won't lose data secure logs provide a global timeline for meaningful audits

"Everyone hears everything in the Auditorium."

Unusual prior art





The Papal Conclave Proceedings **closed to outsiders** All ballots cast **in plain view** All ballots **secret**

How do you audit a secure log?

"Audit logs are useless unless someone reads them. Hence, we first assume that there is a software program whose job it is to scan all audit logs and look for suspicious entries."

-Schneier & Kelsey '99

Where is that program?

"suspicious" is domain-specific

QUERIFIER: an audit log analysis tool

Predicate logic for expressing rules over secure logs Key predicate: "precedes" — requires graph search Querifier runs on a complete log ("OK" / "Violation") or iteratively on a growing log ("OK so far" / "Violation")

D. Sandler, K. Derr, S. Crosby, and D. S. Wallach. **Finding the evidence in tamper-evident logs.** In Proceedings of the 3rd International Workshop on Systematic Approaches to Digital Forensic Engineering (SADFE'08).





Privacy

Secure log of votes could be a problem

When decrypted for tallying, votes are exposed in **order** An observer could match them with voters Loss of privacy \rightarrow bribery & coercion*

Anonymity through clever ballot ordering

re-encryption mixnets

lexicographic sorting

These would still require the ballots to be removed from the ordered audit logs

Ballots in VoteBox

logically, a cast ballot is a vector of counters one per *candidate*

e.g., for one race with three candidates: ballot = (a, b, c) $a, b, c \in \{0, 1\}$

ballots may therefore be summed

tally = \sum ballot_i = ($\sum a_i, \sum b_i, \sum c_i$)



Ballots should be sealed

protected from prying eyes

once cast, they should be readable only by the parties trusted to count them

But how do we count them?

Remember, we don't want to decrypt them in order

Homomorphic encryption

An encryption scheme with a special property

mathematical operations can be performed on ciphertexts, the result of which is a valid ciphertext

We can use this to tally without decrypting

e.g.,

$$\mathrm{E}(x)\odot\mathrm{E}(y)=\mathrm{E}(x+y)$$

for some homomorphic operation " \odot "

Homomorphic ElGamal does this nicely

Other research voting systems use this cryptosystem **Adder** [Kiayias et al. '06]; **Civitas** [Clarkson et al. '08]; **Helios** [Adida '08]

ElGamal encrypted ballots

Encryption & decryption:

$$E(c, r, g^{a}) = \langle g^{r}, (g^{a})^{r} f^{c} \rangle$$
$$D(\langle g^{r}, g^{ar} f^{c} \rangle, a) = \frac{g^{ar} f^{c}}{(g^{r})^{a}} = f^{c}$$

Homomorphic property using multiplication:

 $\langle g^{\mathbf{r}}, g^{a\mathbf{r}} f^{\mathbf{c}} \rangle \cdot \langle g^{\mathbf{r}'}, g^{a'\mathbf{r}'} f^{\mathbf{c}'} \rangle = \langle g^{\mathbf{r}+\mathbf{r}'}, g^{a\mathbf{r}+a'\mathbf{r}'} f^{\mathbf{c}+\mathbf{c}'} \rangle$

- f, g group generators
- *c* plaintext (counter)
- *r* random (chosen at encryption time)
- *a* (private) decryption key
- g^a (public) encryption key

cast-as-intended

How can I be sure my vote is faithfully captured by the voting machine?

"software independence"

an *undetected* system problem cannot create an *undetectable* change in the results

or,

equipment failures can't affect the result

paper — directly inspect the ballot before casting
electronic — ?

current DREs fail this test miserably



this doesn't work:

"logic & accuracy testing"

this is helpful: trusted hardware

VoteBox's approach: ballot challenge

ballot challenge

a technique due to [Benaloh '07]

at the end of the voting session:

- force the machine to commit to the ballot it is about to cast
- 2. the voter chooses to **cast** the ballot or **challenge** the machine to reveal its commitment

the voting machine cannot distinguish this from a real vote

no artificial L&A testing conditions
ballot challenge



ballot commitment

What is the commitment?

How do we force the machine to produce proof of what it's about to cast on the voter's behalf?

Benaloh's proposal

print the encrypted ballot behind an opaque shield You can't see the contents, but you can see the page the computer cannot "un-print" the ballot

How do you test the commitment?

Decrypt it.

But decryption requires the private key for tabulating the whole election!

ElGamal encrypted ballots

More than one way to decrypt a counter:

$$E(c, r, g^{a}) = \langle g^{r}, (g^{a})^{r} f^{c} \rangle$$

$$D(\langle g^{r}, g^{ar} f^{c} \rangle, a) = \frac{g^{ar} f^{c}}{(g^{r})^{a}} = f^{c}$$

$$D(\langle g^{r}, g^{ar} f^{c} \rangle, r) = \frac{g^{ar} f^{c}}{(g^{a})^{r}}$$

- f, g group generators
- *c* plaintext (counter)
- *r* random (chosen at encryption time)
- *a* (private) decryption key
- g^a (public) encryption key

challenging the machine

When challenged, the machine must reveal r

We can then decrypt this ballot (only) and see if it's what we expected to see

In Benaloh, the encrypted ballot is on paper

An irrevocable output medium

decrypting requires additional equipment

VoteBox happens to have its own irrevocable publishing system

One that doesn't run out of ink or paper

Auditorium.

Challenges in Auditorium

When challenged,

a VoteBox must **announce r on the network** Irrevocable thanks to the properties of Auditorium We still need help decrypting the commitment, even given **r**

If we are careful, we can send challenges offsite

Allow a third party to assist in verifying the challenge Trusted by the challenger!



Ballot challenges: cast-as-intended verification preserving privacy without artificial test conditions.

3. Conclusion







lots of research on individual pieces of the e-voting problem

VoteBox integrates these techniques in a single system.

Auditorium (Sandler et al.)

robustness, tamper-evidence

Ballot challenge (new adaptation of Benaloh) verifiability

Other techniques

Smaller TCB through pre-rendered UI [Yee '06]

D. R. Sandler, K. Derr, D. S. Wallach. **VoteBox: A tamper-evident, verifiable electronic voting system.** In USENIX Security 2008.

platform



VoteBox is open-source

votebox.cs.rice.edu & code.google.com/p/votebox suitable for further research, HCI experiments, class projects, security analysis



HCI research

Platform for human factors research & experimentation

VoteBox's ballot designed jointly with CHIL

VoteBox-HF includes extensive instrumentation for HCI work

Questions answered include:

"Do DREs improve performance?" "Do voters notice if DREs malfunction?"

Research output

workshop papers, journal articles, conferences (CHI), two theses

Collaboration ongoing





thanks



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