Zero-Knoweldge Proofs

Abhishek Jain



Zero Knowledge: History

Invented by Shafi Goldwasser, Silvio Micali, Charlie Rackoff in 1980s







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- ▶ Paper rejected three times! Accepted the 4th time in 1985.
- ► Shafi and Silvio won the 2012 Turing Award for work on encryption and proof systems.

Scenario: Where's Waldo?







A "Hey Bob, I found Waldo!"



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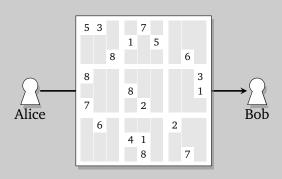






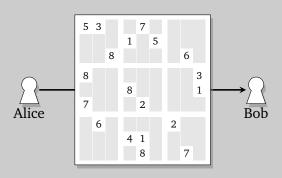
- A "Hey Bob, I found Waldo!"
- B "That was way too fast, I don't believe you."

Scenario: Sudoku



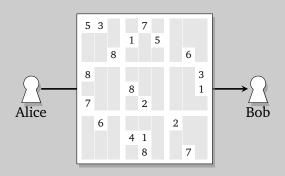
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- A "This one has a solution, trust me."





A "Can I have access to the database? It's me, Alice."

Scenario: Authentication



- A "Can I have access to the database? It's me, Alice."
- B "OK, send me your password so I know it's you."

A Problem of Trust and Information

Alice wants to convince Bob of something

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What might a possible solution look like?



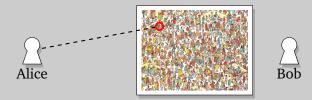






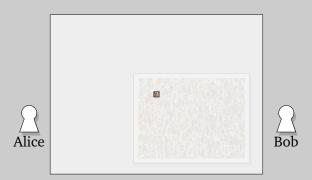
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1. Alice places opaque cardboard with hole over picture, revealing Waldo



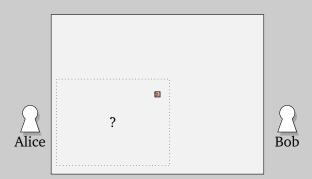


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Bob gets no information about Waldo's location within picture!



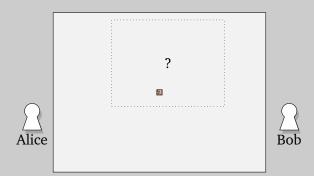


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Philosophy

Fuzzy Definition

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Fuzzy Definition

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What does it mean to

- prove something?
- give out information?

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Today: Interactive Proofs











A true story [R. Fisher, Mathematics of a Lady Testing Tea, 1956]:

M "Tea poured into milk tastes different than milk poured into tea."





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Intro Zero-Knowledge Defining Constructing Generalizing Applications

Fisher's Smart Idea: Interactive Proof







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Random challenge In private, flip a coin to decide which to pour first (tea or milk).

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Response Muriel guesses.

- If Muriel can really tell, she gets it right.
- ▶ If no difference in two kinds of teas, she has 1/2 chance of guessing correctly.



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- ▶ If Muriel can really tell, she gets it right.
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Repeat Repeat *n* times.

▶ If no difference in two kinds of teas, she has $(1/2)^n$ chance of guessing *all* correctly.

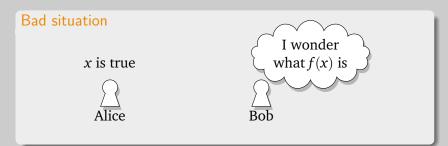


Intro Zero-Knowledge Defining Constructing Generalizing Applications

Epistemology: What is Knowledge?

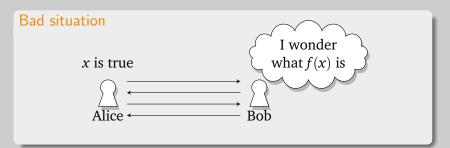


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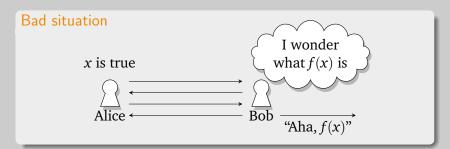




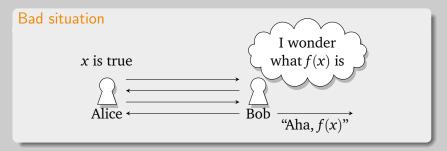
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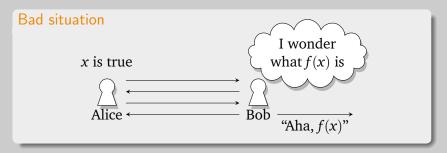






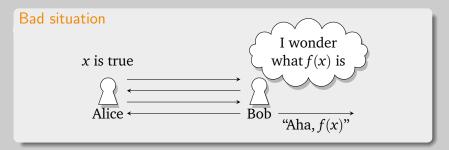


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Zero Knowledge (Want to say):

Everything Bob can compute *after* seeing the transcript, he could have computed *before* seeing the transcript.



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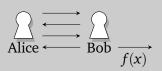
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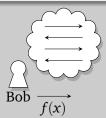
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Clever Definition

Interaction is zero-knowledge if Bob could generate transcripts without interacting with Alice:



Whatever Bob could compute after seeing the transcript ...



... there is a way to compute without interaction!

Paradox?

- ▶ Transcript should *convince* Bob of something new
- ▶ Bob could have generated transcript himself







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Bob already knew the correct responses to challenges

► Convinced by *how* the transcript was generated (in response to his challenges)



Formal Definition

Definition [GMR 1985]

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- ► The prover can always convince the verifier of any true statement
- ► The verifier can't be convinced of a false statement (even by a cheating prover), except with very low probability
- ► There is an efficient procedure to output "same-looking" protocol transcripts



Zero-knowledge protocol:











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1. Alice randomly relabels $\{1, \dots, 9\}$





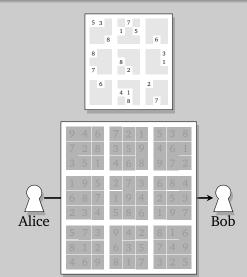






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1

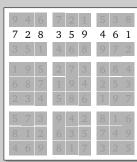
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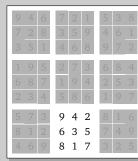


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4. Repeat *n* times









Observation

If Alice can answer all challenges successfully, her scratch card satisfies:

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What if Alice is cheating (there really is no solution)?

⇒ No scratch card can correctly answer all challenges.



Suppose Alice tries to prove an incorrect statement. Let c be a challenge that is bad for Alice's scratch card.

- Bob picks random challenge (28 choices)
- ▶ With probability 1/28, Bob chooses c and Alice is caught!
- With probability $\leq 27/28$, Alice's cheating undetected



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Key Idea

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Key Idea

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When n = 2500, Alice caught with 99% probability.



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Each of these Bob could have generated himself (without the solution)!



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Theorem

Every NP statement can be proven in zero-knowledge.





Lots of things!

Authentication without passwords



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- Stronger encryption
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Disclaimer: ZK proofs very bad for teaching courses:

- Students convinced that professor knows a lot
- Students gained no additional knowledge



Take 600.442!



Thanks for your attention!

fin.

I hope this was a "talk about zero knowledge," not a "zero-knowledge talk."