

# CKY, RECOGNIZER VERSION.

**INPUT:** string of  $n$  words.

**OUTPUT:** yes/no

**DATA STRUCTURE:**  $n \times n$  table.  
rows labeled 0 to  $n-1$   
columns labeled 1 to  $n$   
cell  $(i, j)$  lists constituents  
found between  $i$  and  $j$

For each  $i$  from 1 to  $n$ :

Add to  $(i-1, i)$  all categories allowed  
for the word between  $i-1$  and  $i$

For width from 2 to  $n$ :

For start from 0 to  $n - \text{width}$ :

Define end to be  $\text{start} + \text{width}$

For mid from  $\text{start} + 1$  to  $\text{end} - 1$

For every constituent in  $(\text{start}, \text{mid})$

For every const. in  $(\text{mid}, \text{end})$

For all ways of combining them (if any):

Add the resulting const.  
 $(\text{start}, \text{end})$  if it's not  
already there.

# EARLEY'S ALGORITHM. (1970)

Nice combo of our previous ideas from today:

- incremental interpretation
- no restrictions on the form of the grammar  
( $A \rightarrow BC \text{ spoon } Dx$  is an okay rule thanks to dotted rules)
- $O(n^3)$  worst case, but faster for many grammars
- uses left context and optionally right context to constrain search

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**INPUT:** string of  $n$  words

**OUTPUT:** yes/no (i.e., recognizer, but can turn into parser)

**DATA STRUCTURE:** columns 0 thru  $n$ , corresponding to gaps between words  
column  $j$  is a list of entries like  
( $i, A \rightarrow XY \cdot ZW$ )  
meaning there could be an  $A$  starting at  $i$ , and we have found the  $XY$  part of it from  $i$  to  $j$ .

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# EARLEY'S ALG, RECOGNIZER VERSION, NO LOOKAHEAD.

Add  $ROOT \rightarrow .S$  to column 0.

For each  $j$  from 0 to  $n$ :

For each dotted rule in column  $j$ ,  
(including those we add as we go!),  
look at what's after the dot:

If it's a word  $w$ , **SCAN**:

If  $w$  matches the input word  
between  $j$  and  $j+1$ ,  
advance the dot and add the  
resulting rule to column  $j+1$ .

If it's a nonterminal  $X$ , **PREDICT**:

Add all rules for  $X$  to the  
bottom of column  $j$  with the  
dot at the start: e.g.,  $X \rightarrow .YZ$

If there's nothing after the dot,

**COMPLETE**:

We've finished some constituent  $A$   
that started in column  $i < j$ .  
So for each rule in column  $i$   
that has  $A$  after the dot:

Advance the dot and add  
the result to the bottom  
of column  $j$ .

Output "yes" just if we have  $ROOT \rightarrow S$   
in column  $n$ .

**NOTE:** Don't add an entry to a column if it's  
already there!