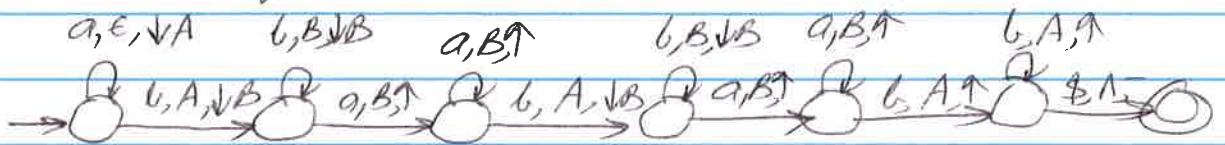


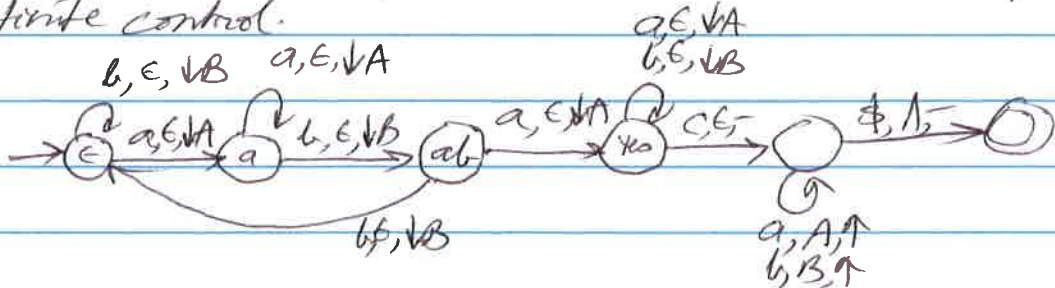
I $L = \{ a^i b^j a^j b^k a^i \mid i, j, k \geq 1 \}$

The nesting indicates the proper matchings.



$L_{10} = \{ x \in \Sigma^* \mid aba \text{ is a substring of } x \}$

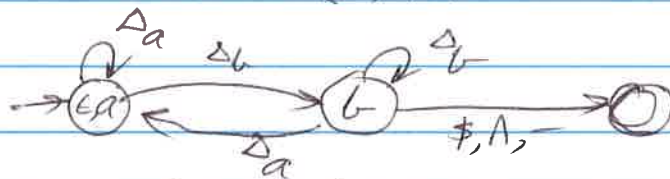
Σ, Σ^* relationship via the pushdown, substring via the finite control.



$L_{18} = \{ x \mid \#_a x = \#_b x, \text{ last symbol of } x = b \}$

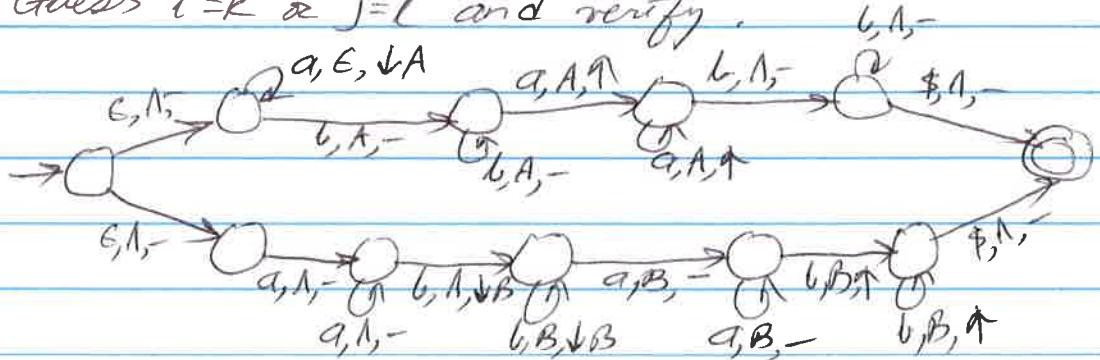
$\#_a = \#_b$ by pushdown; last symbol by the finite control.

Let $\Delta_a = \begin{cases} a, A, \downarrow A \\ a, \Lambda, \downarrow A \\ a, B, \uparrow \end{cases}$ and $\Delta_b = \begin{cases} b, A, \uparrow \\ b, \Lambda, \downarrow B \\ b, B, \downarrow B \end{cases}$



II $L_2 = \{ a^i b^j a^k b^l \mid i, j, k, l \geq 1, (i=k) \text{ or } (j=l) \}$

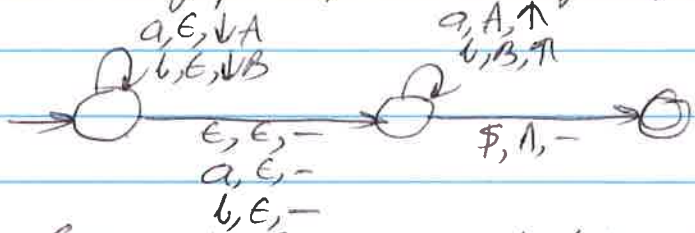
Guess $i=k$ or $j=l$ and verify.



$$L_3 = \{x \mid x = x^R\}$$

$x = x^R$ iff ~~$\exists y \in \{a,b\}^*$~~ $(\exists y \in \{a,b\}^*, c \in \{a,b\}) (x = y y^R c \text{ or } x = y c y^R)$

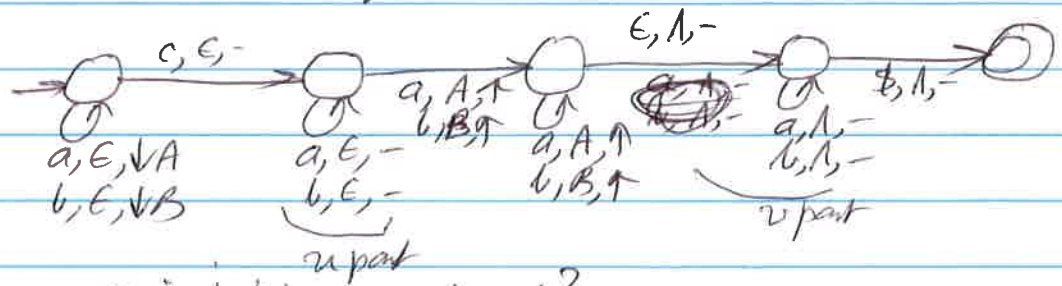
On the y part push, on y^R match & pop.



$$L_{13} = \{x c y \mid x^R \text{ is a substring of } y\}$$

$$y = u x^R v$$

Push x , guess & throw away u part, then match. After the pd becomes empty throw away the v part.



III (a) $L = \{a^i b^j c^k \mid 1 < i < j < k\}$

For any k , let $z = a^k b^{k+1} c^{k+2}$. Note that $z \in L$ & $|z| = 3(k+3) > k$.

For any way of writing

$$a^k b^{k+1} c^{k+2} = u v w x y \text{ st. } |v w x| \leq k \text{ \& } |v x| > 0$$

since $|v w x| \leq k$, $v w x$ and hence $v x$ cannot contain all b 's.

If $v x$ doesn't contain c 's then $u v^3 w^3 x^3 y \notin L$ since $i < j < k$ cannot be maintained.

If $v x$ doesn't contain a 's then $u v^0 w^3 x^3 y \notin L$ since $i < j < k$ cannot be maintained.

Hence L is not an npda lang.

(m) $L = \{x \mid \#_a x = \#_b x \geq \#_c x\}$

let $z = \overset{k_1}{a} \overset{k_2}{b} \overset{k_3}{c}$. As before vx cannot contain all b's.

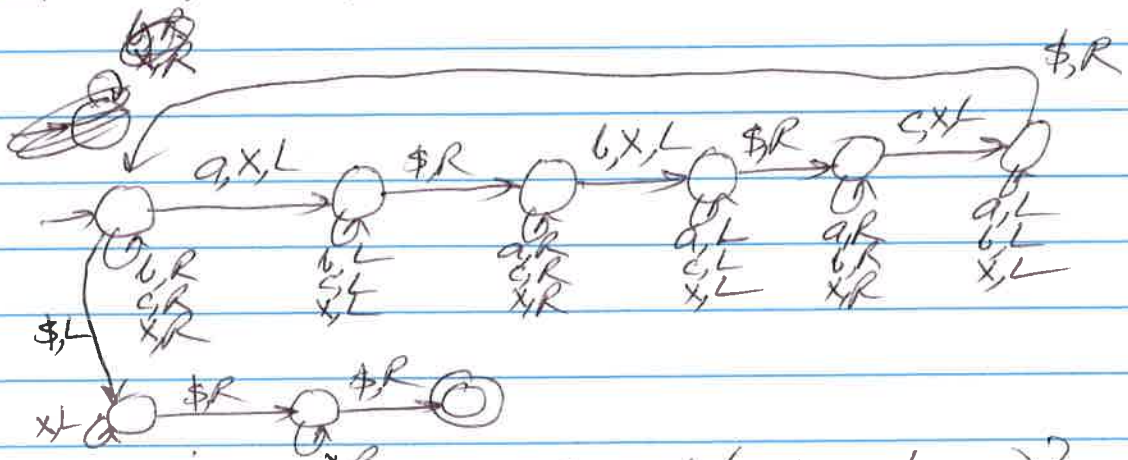
If vx doesn't contain c's, then $uv^0wx^0y \notin L$ since d's and/or b's will be reduced & $\#_a z \geq \#_b z, \#_b \geq \#_c z$ cannot hold.

If vx doesn't contain a's, then $uv^2wx^2y \notin L$ since $\#_a z = \#_b z$ and/or $\#_a z \geq \#_c z$ cannot hold.

Hence L is not an npla lang.

IV (a) $\{x \mid \#_a x = \#_b x = \#_c x\}$.

Repeatedly cross off one a, one b, one c.



(e) $\{a^i b^j c^k \mid i, j, k \geq 1, i=j \text{ and } (k=i \text{ or } k=2i)\}$

check form $\in a^i b^i c^k$, check $i=j$; then $\#_c = \#_a$ or $\#_c = 2\#_a$.

See next page.

