



Stacks and Queues

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



What is a Stack?

Stores a set of elements in a particular order
Accessed in Last-In-First-Out (LIFO) fashion

Real life examples:

- Pile of books
- PEZ dispenser
- Cup trays in cafeteria

CS examples: program execution stack,
parsing/evaluating expressions

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Stack Abstract Data Type

push(*o*): insert *o* on top of stack

pop(): remove object from top of stack

top(): look at object on top of stack (but don't remove)

size(): number of objects on the stack

isEmpty(): does (size == 0)?

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Java Interface for Stack ADT

```
public interface Stack {
    public int size();
    public boolean isEmpty();
    public Object top() throws
        StackEmptyException;
    public void push(Object element);
    public Object pop() throws
        StackEmptyException;
}
```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Array-based Stack Implementation

Allocate array of some size

- Maximum # of elements in stack

Bottom stack element stored at index 0

first index tells which element is the top

increment *first* when element pushed,
decrement when pop'd

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Array-based Implementation

```
public class ArrayStack implements Stack {
    private Object[] S;
    private int topIndex = -1;
    public void push(Object obj) throws StackFull {
        if (size() == S.length)
            throw new StackFull("full");
        S[++topIndex] = obj; }
    public Object pop() throws StackEmpty {
        if (isEmpty())
            throw new StackEmpty("empty");
        Object elem = S[topIndex];
        S[topIndex--] = null;
        return elem; }
}
```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Analysis

Each operation is $O(1)$ running time

- Independent of number of items in stack
- push, pop, top, size, isEmpty

Space can be $O(n)$ or may be much more

- depends if n is known at initialization time



Linked List Stack Implementation

Benefits

- Avoids maximum stack size restriction
- Only allocates memory for stack elements actually used

How

- Allocate a node for each stack element
- Nodes are chained together by reference to next node in the chain



Linked List Node

```
public class Node {
    private Object element;
    private Node next;
    public Node(Object e, Node n) {
        element = e; next = n; }
    ...
}
```



Linked List Stack Implementation

```
public class LinkedStack implements Stack {
    private Node top = null;
    private int size = 0;
    public void push(Object elem) {
        Node v = new Node();
        v.setElement(elem);
        v.setNext(top);
        top = v;
        size++; }
    public Object pop() throws StackEmpty {
        if (isEmpty()) throw new StackEmpty("empty");
        Object temp = top.getElement();
        top = top.getNext();
        size--;
        return temp; }
}
```



Analysis

All stack functions still $O(1)$

- push, pop, top, size, isEmpty



What is a Queue

Stores a set of elements in a particular order

Accessed in First-In-First-Out (FIFO) fashion

Real life examples:

- Waiting in line at cafeteria
- Waiting on hold for technical support

CS Example: Buffered I/O



Queue ADT

enqueue(*o*): Insert object *o* at *rear* of queue

dequeue(): remove object from *front* of queue

size(): number of elements

isEmpty(): size == 0?

front(): look at object at front of queue

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Queue Interface in Java

```
public interface Queue {
    public int size();
    public boolean isEmpty();
    public Object front() throws QueueEmpty;
    public void enqueue (Object element);
    public Object dequeue() throws QueueEmpty;
}
```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Array-based Queue Implementation

Array of fixed size

Index array element for front and rear of queue

Indices “wrap around” when they cross end of array

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Array Queue Implementation

```
public class ArrayQueue implements Queue {
    private Object[] Q;
    private int size=0;
    private int front=0, rear = 0;
    public void enqueue(Object o) {
        if (size() == Q.length) throw
            new QueueFull("full");
        Q[rear] = o;
        rear = (rear + 1) % Q.length;
        size++;
    }
}
```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



List Queue Implementation

Head and tail node references for front and rear of queue

Insert at tail, remove from head

- Remove from tail too slow for singly linked list
 - Updating tail reference with new tail takes full traversal
- So use tail of list for rear of queue

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



List Queue Implementation

```
public class ListQueue
    implements Queue {
    private Node head = null;
    private Node tail = null;
    private int size = 0;
    . . .
}
```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



List Queue

```

public void enqueue(Object obj) {
    Node node = new Node(obj, null);
    if (size == 0)
        head = node;
    else
        tail.setNext(node);
    tail = node;
    size++;
}
public Object dequeue() {
    Object obj = head.getElement();
    head = head.getNext();
    size--;
    if (size == 0)
        tail = null;
    return obj;
}

```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Analysis

All queue operations are $O(1)$

- `size()`, `isEmpty()`
- `enqueue()`, `dequeue()`, `front()`

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Double-ended Queue

Sometimes called “deque” (děk)

Similar to stack and queue

- Allows insertion and removal at both ends of the queue
- Stack or queue is easily implemented using a deque

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Deque ADT

`insertFirst(e)` : insert element at front

`insertLast(e)` : insert element at rear

`removeFirst()` : remove first element

`removeLast()` : remove element at rear

`first()` : examine first element

`last()` : examine last element

`size(e)` : number of elements in deque

`isEmpty()` : `size == 0?`

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Doubly Linked List

Singly linked list inefficient for removing from tail

Has prev reference as well as next reference

Can use *sentinel* nodes to reduce the special cases

- node has no element, just next or prev reference

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Deque Implementation

```

public class MyDeque implements Deque {
    DLNode header, trailer;
    int size;
    public MyDeque() {
        header = new DLNode();
        trailer = new DLNode();
        header.setNext(trailer);
        trailer.setPrev(header);
        size = 0;
    }
}

```

Johns Hopkins Department of Computer Science
Course 600.226: Data Structures, Professor: Jonathan Cohen



Deque Implementation

```
public void insertFirst(Object o) {
    DLNode second = header.getNext();
    DLNode first =
        new DLNode(o, header, second);
    second.setPrev(first);
    header.setNext(first);
    size++;
}
```

Could be null if no sentinels
(trailer could also be null)



Deque (no sentinels)

```
public void insertFirst(Object o) {
    DLNode second = header;
    DLNode first = new DLNode(o, null, second);
    header = first;
    if (second == null)
        trailer = first;
    else
        second.setPrev(first);
    size++;
}
```



Deque Analysis

All operations still $O(1)$

Doubly linked list

- nodes slightly larger
- more references to keep up to date